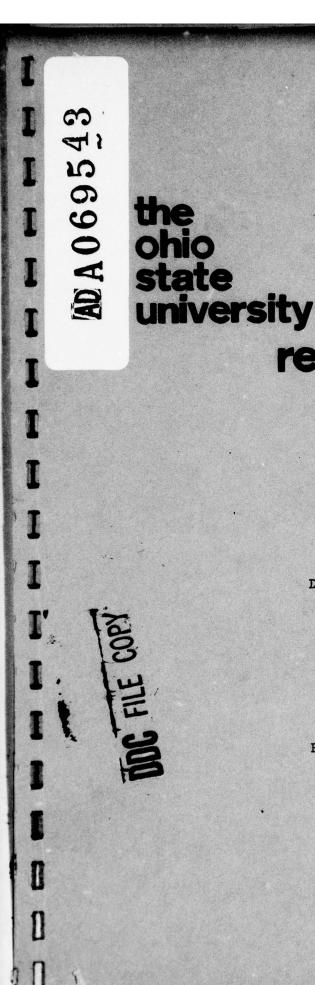
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A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS-ETC(U)

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A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYSTEMS FROM A DBMS AND USER PERSPECTIVE: EXTENSIONS TO THE INFORMATION PROCESSING SYSTEM SIMULATOR

Thomas G. DeLutis
Department of Computer and Information Sciences

For the Period August 1, 1977 - January 31, 1978

U.S. ARMY RESEARCH OFFICE P.O. Box 12211 Research Triangle Park, North Carolina 27709

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The process of analyzing and designing contemporary information systems is a complex one. Models of proposed system designs can be extremely helpful in determining their behavior prior to the expenditure of manpower, dollars and time on unsatisfactorily performing systems.

The goal of this research was to investigate simulation methodologies and language capabilities which could ease model synthesis and analysis efforts in support of system analysis and design. This research focused specifically on data base subsystems and on satisfaction of user performance requirements. (over

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An integral part of the research was the modeling of a complex information system architecture to test the appropriateness of the information processing system simulator design facility (IPSS) as a design tool.

This document describes the investigative and modeling efforts and reports

on the following conclusions:

1. IPSS is a useful tool for modeling in a topdown, modular manner complex information system architectures, and

2. The proposed new capabilities would significantly reduce modeling effort.

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A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYSTEMS FROM A DBMS AND USER PERSPECTIVE: EXTENSIONS TO THE INFORMATION PROCESSING SYSTEM SIMULATOR

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1.0 SNYOPSIS OF THE RESEARCH

1.1 INTRODUCTION TO THE PROBLEM

In the past quarter century, computer technology has had a profound impact upon the manner in which data is processed. With each advance in computer hardware technology has come associated advances in software capability both in respect to the processing of data into more meaningful packets and in the management of the computer hardware. A seemingly never ending cyclic phenomenon has developed with regard to computer based information processing systems. Each expansion in the processing capabilities only increases the appetite of the user community and consequently new demands are placed on the system designers and analysts to increase system throughput, reduce response time and provide more services while maintaining or reducing costs. This has mandated the better utilization of expensive computer hardware and software resources and associated personnel and thus, the development of complex system architectures incorporating sophisticated supervisors to handle the job scheduling, task management and resource allocation functions, multiprocessors to increase throughput and integrated distributed data bases to meet the user's information needs.

Contemporary computer based information processing system architectures are reflective of the information systems they support. As the information needs become more sophisticated, so does the computer system. In part, this phenomenon is the result of new technologies which permit increased volumes of data to be incorporated into information systems at reduced operating cost and faster retrieval time per quantum of data stored. Unfortunately, the complexities of modern society have not only increased the enterprise's need for information but also the transiency of data's value. Furthermore, needed information is more integrative in nature in order to reflect the interdependencies of modern organizations. Ill-designed systems which do not meet their operational requirements are expensive in terms of the time, manpower and dollars lost in their design and implementation.

The design and analysis problem is a complex one, and its complexity is augmented by the costs associated with such systems. Furthermore, the analyst cannot, in most cases, rely on past experiences since each system must be tailored to the particular environment in which it will operate. The need exists, therefore, for modeling technologies which provide the analyst/designer with insights into the behavior of complex on-line, real-time information processing system. Additionally, such technologies must be attuned to the descriptive and evaluative needs of the information processing system so as to minimize the cost and effort for their use both with regard to synthesizing models and evaluating statistics resulting from their execution.

Ideally, models provide the designer and analyst with the ability to identify and characterize an information system's user, data base, software and hardware components and their logical, control and data interfaces. During the lifecycle of an information system, modeling activities provide a feedback capability to the designer and analyst which enhances their ability to attain and maintain performance objectives. By being representations of systems or subsystems, models serve three purposes -- as an aid to design by providing insight into the behavior of selected system components; as a predictor of performance changes that would result for alterations to existing systems; and as a guide to monitoring activities by identifying components to be monitored and statistics to be collected.

Discrete event simulation models can provide a number of benefits to system designers and analysts. They can also be expensive endeavors which remain unused. Among the important benefits to be achieved from this form of performance assessment are:

- Insight gained into the behavior of the constituents comprising a system,
- 2. Identification of performance sensitive components,
- Identification of system and component performance in response to user activity (e.g., volume and mix), and

4. Estimates of performance changes which would result from new hardware, software (1), and data base architectures and system procedures.

Simulation models serve as predictors of system activity; they are cost effective only when the costs associated with model synthesis, validation and evaluation are significantly lower than the cost of building benchmarks or prototype systems.

The modeling of contemporary systems requires the inclusion of submodels for data base, data base access software and support facility components. Without a systematic and modular approach to the total modeling effort, costs can become prohibitive. Therefore, the first step in achieving a modeling capability is the identification of a set of functional components common to the information processing system architectures to be investigated. A second requirement for a cost effective modeling capability is the availability of models which permit the analyses of alternative architectures and/or system activity with a minimum of modification. Third is the ability to incorporate new evaluative technologies and procedures into existing models. Fourth, simulation models must produce meaningful statistics from model evaluation; in particular, it is important that the statistics they generate are useful to the analyst and designer and that they can be validated. Finally, the performance assessment system should allow the reuse of component characterizations in order to effect technology transfer and thus reduce overall cost.

To be effective, simulation facilities must be compatible with the designer and analyst's evaluation goals. This implies the ability to perform a spectrum of modeling tasks from macroscopic analyses of preliminary designs to detailed assessments of instruction execution characteristics. Model synthesis must be consistent with the information known about the system and output data commensurate with the evaluation needs. Required is an integrated view of information system

⁽¹⁾ Software includes operating system modules, application programs and general data base access software.

activities and interfaces which permits a top down, modular approach to model synthesis and is capable of accommodating present and anticipated hardware and software architectures.

1.2 RESEARCH OBJECTIVES

The Information Processing System Simulator (IPSS) has been $developed^{(2)}$ with four general objectives in mind:

- 1. Ease of model snythesis,
- 2. Modularity in model construction,
- 3. System self-containment, and
- 4. Portability.

Additionally, it is recognized that the system must maintain executional efficiency and must be usable by a wide spectrum of users, from the novice to the sophisticate. A continuing goal is to further the development of IPSS within the framework of the above objectives.

The purpose of this research was to investigate methods for extending the IPSS language and statistical facilities, focusing upon those facilities which could increase the definitional and evaluative capabilities of IPSS. The research has led to the incorporation of new language and statistics gathering features into the IPSS. These include (a) DBMS -- associated definitional statements and performance measures, and (b) user and task-oriented evaluative capabilities which provide:

- Definitional statements to ascertain the performance of a model relative to user goals;
- Definitional statements specific to network and hierarchical type DBMS architectures;
- Enhanced standard statistical outputs to provide additional insights into model behavior; and
- 4. Interfaces to internal model facilities thereby permitting additional model control and statistics generation.

⁽²⁾IPSS's development was sponsored in part through research grants GN-36622 and SIS-75-21648 from the National Science Foundation. (Office of Science Information Services) to The Ohio State University.

These new features permit modeling activities to be more closely identified with actual analysis and design functions. Additionally, they reduce model antificiality and provide more meaningful and identifiable statistics. The benefit to the modeler is a reduction in time and effort required to synthesize and validate models of complex information processing systems.

1.3 RESULTS OF THE RESEARCH

The research proposal identified modeling constructs and statistics as candiates for inclusion into the IPSS design facility. The vehicle for assessing the usefulness of these extensions was the modeling of a complex system architecture to be described in Section 3. Details of the evaluative efforts and the results established are found in Appendices A, B and C.

For evaluative purposes, the proposed IPSS extensions were placed in the following categories:

- A. MODEL STRUCTURE RELATED
 - Data Base Subsystem Submodels
 - a. IPSS Data Base Access Component
 - . IPSS Data Base Structure Component
 - 2. Newwork of Submodels as a Model
- B. LANGUAGE EXTENSION RELATED
 - 1. Data Base Subsystem Definition
 - 2. System Effectiveness Evaluation Definition
 - Network Definition
- C. STATISTICS RELATED
 - 1. System Effectiveness Evaluation
 - 2. Data Base Subsystem Behavior
 - General

In Appendix C the results of the evaluation are discussed. In all 37 statements and 123 statistics identified in the proposal were evaluated. The results of the evaluation are summarized in the following table.

Table 1.3-1. Summary of the Evaluation of Proposed IPSS Language Extensions

ASSESSMENT CONCLUSION	DISPOSITION	TOTAL	
	CODE	Statements	Statistics
Not Examined	NE	6	20
Examined and Rejected	ER	0	0
Examined and Found to be Useful —			
Has been implemented as proposed	EU-I	16	0
Has been implemented but merged into another facility	EU-IM	6	7
Has been defined and not implemented	EU-D	0	93
<pre>Has been defined, used in paper models, but not implemented</pre>	EU -DP	9	0

The modeling activities also demonstrated the feasibility and suitability of applying IPSS and its underlying methodology to the synthesis of models of complex information systems. The system modeled to demonstrate these capabilities was a host/back-end processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system function of task management, job scheduling and resource allocation.

The IPSS methodology proved capable of providing the modeler with a modular, top-down approach to system analysis and model synthesis thus facilitating a solution to this complex problem. Several models were constructed, each of which reflected system details of a particular aspect of the system. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their

integration into an overall complete model. The SIDPERS/IDMS workload model represents the software processes of the applications and IDMS in detail while characterizing the operating systems and computer hardware as black boxes. From this basic model, two subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood synopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicating both potential bottlenecks and under-utilized components.

The models were all internally verified. That is, through experimentation, the internal processing consistency was verified to be accurately reflective of the real world processes. Due to the lack of operational data, however, it was not possible to validate them through comparison to actual system performance. Thus, although some confidence can be placed in the relative values of the statistics reported in Appendix B no validation has been made with respect to the magnitude of the data values.

The experiment demonstrated the suitability of IPSS to model complex systems while requiring minimal time for model synthesis. The modeling effort required approximately four man months. The resulting models have proven to be versatile as well as adaptable to changing requirements indicating that significant portions can be used for other application areas. For example, the DMBS and operating system

processes as well as hardware facilities need not be modified when another major application area is modeled, thus resulting in a substantial reduction in future related modeling efforts.

In addition, this experiment has served to focus future development work as well as to validate current goals. In particular, the experiment demonstrated the need for simulation facilities specifically attuned to the characterization of integrated data structures as represented in IPSS/DBS. These initial models did not use these features since they were not fully operational, and thus the logical structure of the data base was characterized at a more abstract level than desired. Furthermore, experimentation with IDMS record mapping policies was not attempted because of the complexities involved in relation to the time available.

A number of experiments on the basic models became readily apparent. These were:

- Identification of the maximum number of terminals which could be supported as the system approaches saturation;
- 2. Response time behavior for system processes and resources as loading increases and/or mix of transaction types is varied;
- 3. Effect on throughput as the transaction error rate is varied;
- Response time behavior as a function of transaction arrival rate;
- 5. Effect on system throughput and response time with a m-way, multi-threaded DBMS;
- 6. Impact of presorting transactions on system throughput. The initial model did not contain operating system or computer hardware characteristics. However, later models did have these characterizations. Thus the effect of changes in their characterizations on throughput and response time could also be examined in future analyses.

The modeling was done using the original IPSS, now designated IPSS/BASIC. The IPSS/DBS features were not employed. As a result detailed models of DBMS behavior could not be easily modeled. However, paper models using the proposed syntax were written comparatively quickly

(one man-month which included review of the IDMS literature). This exercise demonstrated the value of the IPSS/DBS capabilities to reduce the modeling effort needed to characterize DBMS software and data base scheme.

Finally, the incorporation of the IPSS/DBS into the current IPSS has extended the concept of information systems modeling into the area of networks of simulation models. Since the IPSS/DBS was designed to execute either as a stand alone or asynchronously with IPSS/BASIC models, it became necessary to devise a mechanism to accommodate this concept. This mechanism has been extended so that an arbitrary collection of nodes (3) can be incorporated into an IPSS model. The implementation requirements for this concept are discussed in Appendix C.

⁽³⁾ An IPSS model node is defined to be an uniquely identified submodel comprised of IPSS/BASIC, IPSS/DBS or IPSS/BASIC-IPSS/DBS components.

2.0 THE INFORMATION PROCESSING SYSTEM SIMULATOR (IPSS)

The Information Processing System Simulator (IPSS) is a special purpose digital simulation facility designed to facilitate the performance assessment of complex computer based information processing systems. It represents the realization of a generalized methodology for the performance assessment of information systems and has been designed to be applicable to every phase of the information processing system life cycle. Additionally, models developed for one phase can be easily and inexpensively modified in both scope and depth of detail. This enables the model to evolve in parallel with the system's progression through its lifecycle.

The methodology employed by IPSS divides the elements of an information processing system into three categories according to the function they serve: data bases, services which access the data bases, and support facilities. As illustrated in Figure 2-1, a system is viewed as a hierarchy of services and data bases with each level supporting the service and data needs of the next higher level and requesting support from the adjacent lower level. Performance measures identify the behavior of services relative to their data base access activity, acquisition and use of support facilities, and use by other services. The methodology emphasizes data base access activity and has identified those system elements at each data base level which contribute to this activity. Within IPSS, the characterization of system elements has been formulated in a manner such that their interaction during model evaluation automatically produces the desired performance measures. An illustration of the modeling process is shown in Figure 2-2.

IPSS is more than a language in the sense of a GPSS or a SIMSCRIPT; it is a complete system. In addition to statements defining resources and tasks for the system under investigation, IPSS contains special statements directing the IPSS system in a number of auxiliary functions including the use and management of user defined library facilities, model compilation, and model execution sequencing.

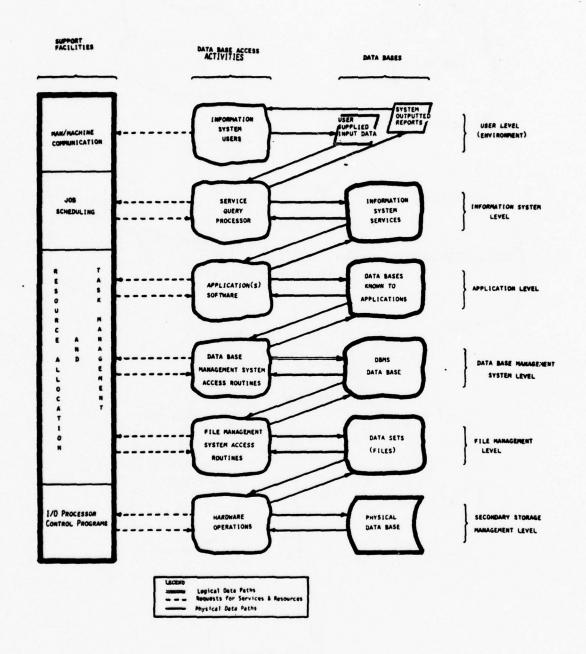


Figure 2-1 Hierarchical View of Information Processing Systems

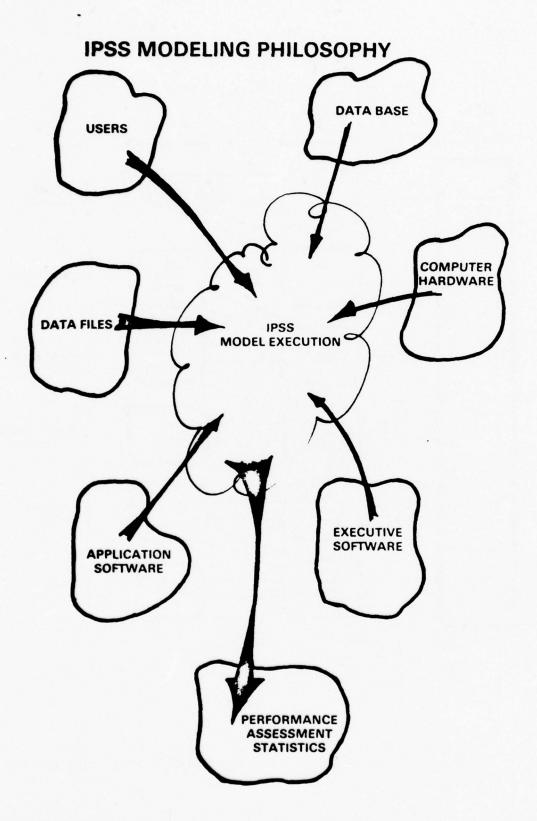


Figure 2-2 IPSS Modeling Philosophy

The IPSS language facilities have been carefully designed to focus the modeler's synthesis activities on the IPSS methodological underpinnings. In this manner the development of models of complex systems can be done in a systematic, top down manner. This affords the modeler the same advantages that these techniques provide in information system analysis, design and implementation. Figure 2-3 illustrates the three phases in the use of IPSS for IPSS model synthesis. At the begining all that is available to the modeler (4) is general knowledge of the information system (at some stage in its life cycle). Phase I is an implicit one during which the modeler employs the IPSS methodology to separate overall system knowledge into specific knowledge of the following four functional components:

- A. System users and their request characteristics;
- B. Services to be provided to users and to other services, and their inter service connections;
- C. Data base resources and configurations;
- D. Hardware resources and configurations.

The IPSS language has been specifically formulated to characterize these functional components,

During Phase II the above functional components are characterized via the IPSS language into independently defined model components. Presently five distinct model components are possible to describe an information system as shown in Figure 2-4. (The sixth component, the model director, is used to control simulation activities during Phase III, the model execution phase.) Not all are required to model a particular computer facility. Additionally, more than one of each component (except model director) can appear in an IPSS model. A synopsis of the role of each component follows:

 SYSTEM RESOURCES -- contains definitions for all hardware, software and data resources comprising an information system model. Included in the SYSTEM RESOURCES component is the IPSS supplied clockwork mechanism to schedule and control simulated events and to determine

⁽⁴⁾ A modeler is assumed to be a systems/analyst or a systems designer.

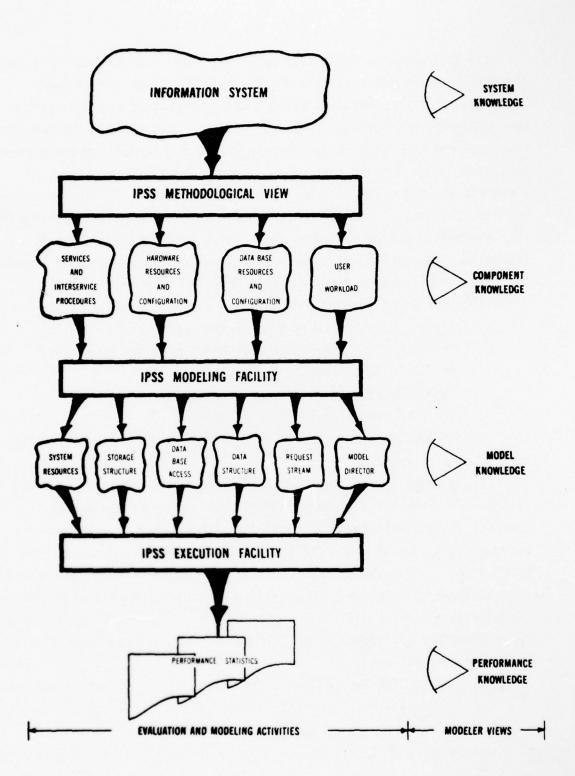


Figure 2-3 Relation Between IPSS Methodology, Language and Execution Facility

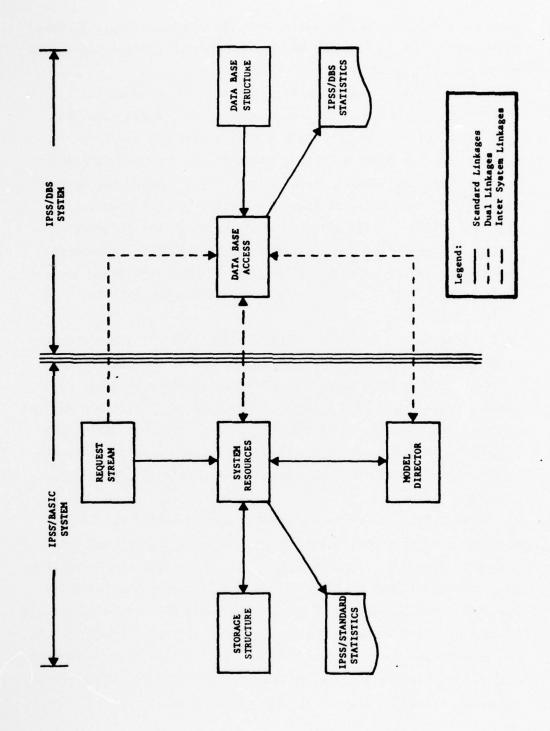


Figure 2-4 Components of an IPSS Model (with DbS Extension)

when the simulation is to terminate. The clockwork logic is based on the next most immediate event philosophy for controlling discrete event digital simulations;

- 2. STORAGE STRUCTURE -- supplements the definition of data set facilities. It provides the model with general definitions for data set files in terms of their formats, size and location relative to each other and their location in secondary storage. Additionally, the STORAGE STRUCTURE component is used to define secondary storage space management policies for the model;
- 3. <u>REQUEST STREAM</u> -- defines the external requests for information system services; that is, the users of the system being modeled. The modeler is required to define types of users and their interarrival times. IPSS converts these times into a composite arrival time stream;
- 4. DATA BASE ACCESS -- contains the definitions of all the resources required by the DBMS. These include the hardware resources of buffers and user work areas as well as application programs and DBMS software. All DBMS-related entity type facilities are defined within the component. DATA BASE ACCESS is similar to the SYSTEM RESOURCES component in that it contains its own simulation clockwork mechanism similar in purpose to that of the REQUEST STREAM component.
- bata structure -- provides the modeler with a set of facilities which allows the definition of logical data structures and the characterization of relationships among them. This can be applied to a variety of DBMS architectures and application environments. The DATA STRUCTURE component permits the modeler to investigate the effects on system behavior caused by alternate set, record type, and access path definitions. The definitional facilities provided enable the modeler to investigate a wide spectrum of logical data structure organizations and allocation policies.

Also illustrated in Figure 2-4 is the division of the IPSS components into two subsystems -- IPSS/BASIC and IPSS/DBS $^{(5)}$. This has been done to permit a modeler to concentrate his efforts on a particular aspect of the problem. Thus, the capability to model only the data base or the system aspects of a total system is provided the modeler.

Associated with each component are IPSS language facilities which permit the modeler to define and characterize information system entities and attributes, gather statistics and control model behavior. Over 100 such statements are available to the modeler.

The function of Phase III is to assemble the model, initiate its execution and display the generated performance assessment statistics. IPSS provides a modeler with a number of statistics concerning the behavior of modeler defined entities and IPSS supplied built-in information system services. (6) These statistics fall into eight general categories:

- 1. Operational Statistics,
- 2. Request Stream Statistics,
- 3. I/O Activity,
- 4. Queueing Statistics,
- 5. Utilization Statistics,
- 6. Wait Statistics,
- Service Statistics, and
- 8. Task/Activity Statistics.

Additionally, the modeler can employ the complete facilities of the Fortran language to develop his own statistics. Statistics are printed automatically at the conclusion of each model simulation unless explicitly inhibited.

(6) These services are characterizations of the basic I/O operations for tape. DASD, unit record, and terminal devices.

This configuration of IPSS components is currently being implemented. The CSC model reported here employed only IPSS/BASIC components. However, the IDMS was also written using IPSS/DBS facilities but has yet to be validated.

Finally, IPSS contains a library facility which permits the modeler to retain previously defined models, model components and code segments for later referral. This feature can greatly reduce modeling time.

An implicit goal of the IPSS is to achieve the maximum applicability of the methodology and the IPSS facility. Therefore, the following criteria have served as guidelines throughout its development.

- Breadth of Applicability -- the ability to model the behavior of contemporary and foreseeable system architectures and operating environment;
- 2. <u>Functional View of Systems</u> -- the ability to identify and characterize system components and activities based on their function independently of a particular architecture or environment;
- 3. Top Down, Modular Model Synthesis -- the ability to model to a level of detail commensurate with the desired analysis of research objectives;
- Extensible Structure -- the capability to incorporate new, higher level descriptive facilities and performance measures into the methodology and simulation system; and
- 5. <u>Flexibility of Use</u> -- the ability to be used by a wide spectrum of modelers from the experienced system analyst/designer and researcher to the practitioner and student.

Examples illustrating the broad flexibility and applicability of the extended IPSS system follow.

Modeling Activities

Modeling activities are those activities concerned with the development of evaluative capabilities for examining the behavior of existing and proposed software, hardware and data base architectures.

 Models of specific data base management systems for the purpose of examining their behavior under a variety of different criteria such as loading, data structure, and data mapping.

- Models for evaluating the influence of data structures on system performance over a broad spectrum of DBMS loadings and support architectures.
- Models of systems where the central site computing facilities are comprised of computer networks in order to examine the effects of alternative scheduling and tasking philosophies.
- 4. Models for evaluating the interaction of front-end and back-end computing capabilities.
- 5. Models for investigating the effects of queue scheduling policies for basic system activities such as job scheduling, task management, resource allocation, and paging strategies for virtual memory systems.
- 6. Models to evaluate data base security procedures; in particular, the identification of authorized users to the system and the impact of such activities on overall performance.
- 7. Development of a simulation capability and support library which permits the modeling of a wide range of contemporary architectures for the purpose of evaluating new technologies as they become available.

Performance Measurement Research

Performance measures are those measures which assist designers and analysts in their performance evaluation activities.

- 1. Measures of system performance based on user oriented goals.
- 2. Measures of the impact of new applications on current systems.
- Development of higher level measures which permit the identification of critical system activities based on applications performance objectives.
- 4. Development of measures and procedures which aid in identifying (from a set of alternatives) the "best" alternative to achieve system performance objectives.

5. Development of measures to assist the data base administrator. (It is assumed that data base administration is an on-going activity and its major function is to maintain and assure the integrity of the data base and to improve its performance via changes in its data structure.)

These activities were possible due to the unique design of IPSS which permits a top down modular view of information processing systems and the special language constructs and facilities developed specifically for the above purposes.

The IPSS system is designed to analyze information processing systems in their totality. Initially, IPSS mechanism for the definition and evaluation of hardware, computer system software and storage level data structures were developed and have been extensively tested. The next phase in the continued development of IPSS was to incorporate features for the description, manipulation and evaluation of logical data structures of the kind normally encountered in modern data base management systems. While the current IPSS capabilities substantially reduce the effort required to model system and data base level software, hardware and dataware, modelers can model all aspects of an information processing system. A high level query processing capability is currently being implemented and development is planned for a telecommunication specific extension.

3.0 BACKGROUND OF THE RESEARCH NEEDS

The main goals of the research are the extension of the current IPSS system to include language constructs and statistics attured to data base management system modeling activities, and to the measurement of system performance from a user perspective. The following discussion provides a background to evaluative methodologies associated with the research needs for these areas.

3.1 METHODOLOGIES FOR THE EVALUATION OF USER GOALS

As the complexity of current and anticipated systems increases, the analyst's need for more sophisticated evaluative techniques becomes more apparent. The types and number of services being provided by an information system are growing rapidly. Where once was a homogenous set of similar user classes which could be evaluated by relatively simple measures, now operates a multitude of user classes and services demanding various conflicting performance requirements and lacking a single common measure of performance. Thus, from an external performance point of view, the analyst must evaluate information systems against a multiple set of user criteria.

From an internal performance view, current and anticipated demands force the analyst to also evaluate multiple criteria. As the number of services provided increases, the size of the delivery system must also increase. This implies that the number of resources required to build the system will increase. Information system resources -- hardware, software and dataware are expensive to buy, install and operate. Thus, the analyst must also be concerned with the cost effective use of these resources.

Current evaluative technologies focus on only one criteria in the system design equation, either the use or the system's resource performance. The ability to simultaneously ascertain the impact of resource performance on user goals or vice versa is not readily achievable through these methodologies. One of the purposes of this evaluative

expansion is to establish a formal liaison between the evaluation of user goals as a function of system behavior and the analysis of resource performance as a function of user activity.

User oriented analyses with objective functions based on response time, throughput, and cost have been (are continuing to be) reported in the literature. Most frequently, analytic approaches use queueing models as their basis (GAV76, BUZ73, KLE76 are representative of this type of analysis). Due to the necessity to maintain tractable models, many simplifications are required for a model's analytical solution.

Simulation models have also been applied to user oriented analysis (CON65, ROE70). Unfortunately, these models yield only average and/or aggregate measures of system response. As a result of these simplifications, the analyses produced by both of the approaches fails in many cases to identify the relationship between users and resources. Therefore, they are suspect when used to predict the impact on system performances of modifying the current environment.

Alternatively, performance analyses can be made from the system's standpoint, treating the user and his goals in the aggregate. The most common approach is a subsystem study, where a particular part of the information system complex is isolated, with the subsystem user(s) represented by a stochastic generator, both analytic and simulative. The most emphasized areas of research has been the I/O subsystem (ABA68, NAH73, BOU72) and CPU utilization (KLE72, HEL70, MLS76). The problem with this level of evaluation is that, although providing valuable local intuitive insight, these models rarely relate to the ultimate information system user, and, therefore, do not provide realistic insight into global performance.

Examinations of complete systems have also been made. Exhaustive hardware/software measurements have been analyzed by Lindsay and Cole (LIN76, COL72) while simulation models, including an aggregate user component, have been built by Reeves and Pooch, Norland, and Lum (REE75, NOR71, LUM70). Alghough results of the evaluations include resource utilization statistics and user oriented measures such as response time, there is little attempt in these models to relate particular resource

usage to the effort on user goal attainment. But from practical experience it is evident that there is indeed a relationship between user goals and resource usage. In fact, Buzen (BUZ76) has recently proposed some fundamental laws for computer performance which relate resource activity to global system/user measures such as response time and throughput.

Thus, current evaluative efforts have not been able to provide an acceptable approach to handling the multiple criteria environment. Because of the inherent conflicting nature of the criteria an optimal design is not possible, and therefore, good system design is assumed to satisfice both performance related criteria. The approach embodied in this evaluative expansion was to first establish a casual relationship (liaison) between user goal attainment and system resource activity. Then after measuring the resultant behavior, a multi-objective technique was employed to produce a satisficed evaluation. The liaison will be specifically maintained in IPSS via the Task and Activity facilities. A mathematical programming technique, multiple goal programming (MGP) has been adapted to evaluate the multiple performance criteria (CHN77).

The importance of providing a task modeling capability is evident by examining current multi-programming system design philosophies. There has been a continuing trend toward hierarachic design of operating systems (DIJ68, ZUR68, MAD76, SDB77), both top-down and bottom-up. Such designs imply that particular functions are assigned to particular levels within the hierarchy, e.g., job scheduling at one level, hardware I/C control at another, etc. Each of the functions (levels) also in a sense, manage a class of resources via allocation, usage and de-allocation. In order to model and evaluate the effect of user demand on each of these functional levels, and thus, on their subordinate resources, one must be able to model the tasking function.

The overall modeling philosophy of IPSS, coupled with the addition of the TASK and ACTIVITY facilities allow the modeler to accomplish this. The purpose of the TASK facility is to maintain control over the resultant resource activity with respect to its causal user demand. This resource activity is further broken down by activities which characterize the individual system functions.

The control that the TASK facility exerts over its subordinate Activities and resources is essentially its maintenance of statistics. In IPSS the modeling resource usage automatically produces an extensive set of queueing and utilization statistics. The characterization of software processing also produces, automatically, service by service statistics on service time and hierarchy interactions. The TASK facility is associated with a particular user class to complete the user-resource liaison. It is the function of the TASK facility to maintain a set of statistics measuring the behavior of its subordinate resources. These statistics, collected dynamically, essentially partition of the aggregate queueing, utilization and service time statistics to each causal TASK (user class). Furthermore, the collection and maintenance of these statistics is performed automatically by IPSS. Thus, the first part of this approach to information systems establishes the link from user demand to resource activity and automatically produces statistics.

The second part of the approach to be implemented in this expansion is an evaluation of the multiple performance criteria. To be most effective, this evaluation should relate resource activity to user oriented performance goals. This entails the production of higher level performance measures, that may not be solely time related such as cost. This completes the two-way liaison between user demand and resource activity by measuring the contributions of resources to the satisfaction of user goals. This is, in part, accomplished via the TASK/ACTIVITY liaison, but the evaluation, however, is performed by a multiple goal programming based procedure.

The evaluative view of an information system formulated by the MGP procedure is identical to the modeling view characterized by the TASK and ACTIVITY facilities whereas using the TASK and ACTIVITY aids in the creation of a top-down modular performance evaluation (PE) model. In fact, the variables of the model are the TASK and ACTIVITIES. This analogy was designed to allow the automatic collection of statistics by the TASK facility during normal modeling be directly applicable to the PE model for evaluation.

The result of evaluating this PE model is two-fold. First, the desired evaluation of the user criteria is produced, on a global goal basis and on an individual goal basis (Task basis). Second, each activity is also evaluated relative to its service time performance, across all tasks (user classes). This evaluation can determine if a given activity is being under-utilized, saturated or used to its proper capacity, where proper is with respect to all tasks. These activity measures can be used as guidelines when deciding which designs/alternatives to follow. The most probable performance inhibiting activities can be identified in this way. An overview of the underlying methodology of this multiple criteria evaluation can be found in the paper by Chandler and DeLutis (CHN77).

3.2 METHODOLOGIES FOR THE EVALUATION OF DATA BASE MANAGEMENT SYSTEMS

Sable (SAB71) outlines a specific evaluation procedure for data base systems which is based on establishing the cost effectiveness of each candidate system with respect to a specific problem. He proposes a multistage method for "scoring" or grading DBMS features. The scores are normalized according to an assigned nominal requirement for each feature and multipled by a weighting factor which represents relative importance. The overall system score is the sum of the weighted results. This methodology relies heavily upon the intuition of the analyst to select the important factors and to assign correct weightings. In addition, it presupposes that the evaluation data are available.

While methodologies based upon this approach are useful for establishing guidelines and controlling the overall evaluation process, they are static in nature and therefore are not useful in measuring dynamic DBMS behavior. However, the current DBMS literature contains several methodologies specifically designed for the measurement of this type of behavior. These methodologies all provide a set of descriptive facilities for the definition of both DBMS software and computer hardware. In addition, they all provide a special purpose, discrete event simulation tool for the analyst/designer.

Phase II (OWE70) is one such tool for modeling hierarchical base management systems. Its objective is to allow the modeler to study the interaction of such DBMS variables as storage structure characteristics, data distribution, device interaction, and access methods. The modeler specifies the characteristics of the data, the data accessing strategies, and the underlying hardware configuration through a special purpose user-oriented language. Data field contents are characterized by distribution functions and data fields are lelated to each other to form what is called segment (the logical structure of the data). The logical data structure is then mapped to a physical storage structure by partitioning the records into files which exist on hardware devices. The accessing strategy allows the modeler to specify the lists of records which are to be accessed, the accessing method (i.e., sequential, indexed, or direct), and the order in which the accesses are to occur. Phase II is limited to the I/O analysis of single level hierarchical data base management systems. Its design objectives are intentionally modest and thus it is not suitable for modelling more complex software interactions.

A methodology for evaluating data base scheme, physical data organization, and data base loading is described by Hulten and Soderlund (HUL76). Their simulator is called ARTS/DB (Analysis of Real Time Systems/Data Base oriented systems). It is implemented in Simula and executes on the DEC10 system. The descriptive facilities which they provide allow the characterization of data base systems in terms of CODASYL-DBTG features. An ARTS/DB model consists of four levels, i.e., definitions of the application, the DBMS, the system software, and the hardware. This provides the analyst with a useful mechanism for performing successive experiments with only slight modifications to the model.

A data base simulator based on the Data-Independent Accessing Model I (DIAM I) (SEN72) has been developed with two ojectives in mind: the comparison and selection of existing data base systems and the design and evaluation of new data base management techniques

(SCH76). This simulator consists of four subsystems similar to the ARTS/DB approach. In this case the subsystems correspond to the levels of the DIAM and are: a query generator, a query translator which uses defined access paths, and I/O generator, and the host computer system.

A three stage DBMS model building approach as well as a simulator system utilizing this approach has been described by Reiter and Finkel (REI75, REI76). First, high level user-supplied logical data view descriptions are translated into an intermediate standard form. Second, an instruction stream is generated which refers to data by block address. In this stage no assumptions are made concerning the location of data in physical storage. Third, the dynamic elements of the system are employed to produce performance measures. The particular DBMS under study is parameterized with respect to data representation, resulting in a specific model of that DBMS. The model is constructed of static and dynamic parts which results in a similarity to operational DBMS and also in executional efficiencies.

DeLutis (DEL72) presents a methodology for the analysis of auxiliary storage systems (AUXSIM). In this system the modeler describes hardware devices and software processes using language statements which are highly descriptive of these devices and processes. The AUXSIM translator produces Fortran source code which executes as a discrete event simulation of the system. This research was the predecessor to the Information Processing System Simulator (IPSS). The research proposed here was, in part, a further extension of the IPSS methodology to data base management systems.

In addition to these generalized evaluation techniques, several specific DBMS models have been developed using a variety of simulation languages and techniques. The primary evaluation criteria used by these simulations is response time and throughput.

Ghosh and Tuel (GH076) designed an experiment to model a large IMS data base to determine the effect of varying such DBMS features as the retrieval sequence and logical data structures. They collected operational data by executing IMS in a stand-alone environment. Their measurement criteria was the elapsed time required to execute a

sequence of calls. Griffith (GRI75) constructed a GPSS simulation model of Univac's DMS-1100 also using response time as the evaluation criteria. The simulation results reported by Nakamura (NAK75) include response time, resource utilization (cpu, channel, disk), as well as statistics related to transactions and queues. Hall (HAL74) describes a DBMS simulation model of a hierarchical data base whose objective was to minimize response time and costs to the user.

4.0 EXPERIMENTS

A series of modeling experiments were performed between mid June and mid September 1977 for the purpose of; (a) assessing the ease of employing the IPSS methodology and language in modeling complex systems, and (b) evaluating the usefulness of the proposed IPSS language and statistics extensions. The system modeled was an on-line real-time system comprised of; (a) data entry terminals; (b) a host main frame for the application software and user interface; and (c) a CODASYL-like data base residing on a backend mini computer. Section 4.1 summarizes the modeling activities. (The reader may refer to Appendix A for a complete discussion of each of these models.) The following subsection provides background as to why such a complex system was chosen.

The purpose of the experiment, initiated by the U.S. Army Computer Systems Command, was to evaluate a complex information system architecture via the Information Processing System Simulator (IPSS). The primary objectives of the experiments were three-fold:

- To demonstrate the capability of IPSS and its underlying methodology to accurately model a complex system,
- 2. To ascertain its ease of use, and
- 3. To develop a generalized basic model of the following systems:
 - a. an interactive version of the SIDPERS Personnel system resident in an idealized Host computer, and
 - b. The IDMS data management system resident in a PDP-11/70 serving as the back-end computer;

and through the model judge the appropriateness of the proposed IPSS language and statistics extensions.

Results indicate that IPSS is well suited for the task. Through the use of a top down modeling philosophy a complete model of the application and IDMS processing was written and tested in less than two man months and model statistics tracked well to the available performance measures. The following subsection provides an overview of the salient features of the modeling effort.

4.1 THE MODELING EFFORT

The purpose of the modeling activities was to demonstrate the feasibility and suitability of the IPSS to model large, complex information processing systems. Reflection on the modeling activities associated with the SIDPERS/IDMS experimentation provides positive evidence that the IPSS is equal to the task. One important reason that IPSS achieves its goal is that its underlying methodology provides a top down, modular approach to system analysis and model synthesis thus facilitating answers to complex analysis and design questions.

Several models were constructed, each of which reflected system details suited to the multi-stage development process which was adopted. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their integration into an overall complete model. Appendix A summarized their major features. The SIDPERS/IDMS workload model represents the software processes of the applications and the IDMS in detail while characterizing the operating systems and computer hardware as a block box. After this basic model was tested, subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. These results are summarized in Section B.2. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood snyopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicated.

The modeling activities were performed with approximately six man months of effort. All the models to be described in the next section were completed. However, three conditions occurred which prohibited the creation of the sought-after final models. These conditions were:

- The IPSS/DBS parsers were not completed in time for use. However, models of the data base reflecting the IDMS scheme have been written in the IPSS/DBS syntax and semantics specifications. Furthermore, these submodels have now been executed under the IPSS/DBS.
- 2. Combination of all the submodels (excluding the IDMS schema) could not be successfully compiled. Investigations into the cause of the problem has lead to the detection of an error in the IPSS nucleus. The nucleus did not detect an IPSS library manager message signaling the lack of library space. Thus model code was lost and was undetected until final model assembly. This condition has been recently fixed and the models have been successfully executed.
- 3. The interfacing mechanism discussed in Appendix C3 has not been completely implemented. This is due to the unexpected quantity of code which had to be modified. Therefore, the IPSS and IPSS/DBS cannot currently execute together. This problem is currently being aleviated.

Section 4.2 discusses the complete modeling program envisioned for the modeling experiment. Appendix B discusses the statistics for the submodels actually completed during the experiment.

4.2 SIDPERS/IDMS MODELS

The purpose of this section is to provide an overview of the system analysis and model design activities related to the IPSS modeling of the interactive SIDPERS/IDMS system. (7) For analysis purposes the system was divided in the following logical model components:

⁽⁷⁾ The IPSS methodology provides a modular, top-down approach to system analysis and model synthesis thus permitting the solution discussed in this section.

Hardware

Host machine configuration
Back-end machine configuration

Software

Operating system for the host machine
Operating system for the back-end machine
Terminal interaction modules
SIDPERS transactions
Interface modules
IDMS

These logical model components were used to define five inter-connected system activities each representing specific SIDPERS/IDMS functions. This organization is shown in Figure 4.2-1. A number of models were written which incorporated a one or more of the modeling components and represented one or more system activities. Figure 4.2-2 identifies the submodels and incorporated model components.

Because of the magnitude and complexity of the system, a phased approach was taken in model development. The model synthesis and validation is implicit in IPSS, thus, the approach posed no difficulties. In Phase I a model of the SIDPERS transactions as processed by IDMS (system activities 1 and 5) was constructed. However, it is important to note that the entire structure of the final model was represented in this preliminary model. This was accomplished through IPSS services for system activities 2, 3, and 4 which were essentially null routines. That is, they were invoked by other services but their internal processing function remained unspecified. In particular, terminal interaction and secondary storage accesses were characterized merely by advancing the simulator clock. This simplified approach allowed the SIDPERS transactions and the IDMS processing functions to be modeled and validated quikly. Once this workload flow model was executing successfully and generating statistics, the operating system and computer hardware model components were added in succeeding modeling phases.

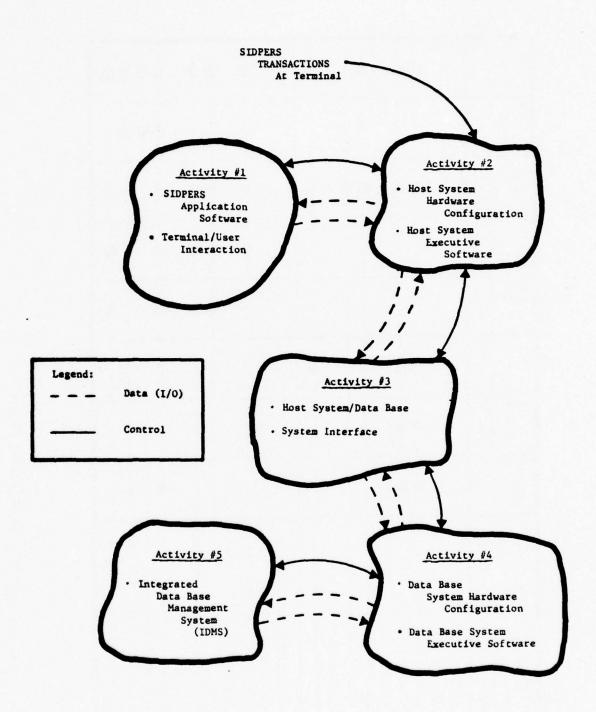


Figure 4.2-1 Structure of Final SIDPERS/IPSS Model

Models System Factiviae	SIUPERS/ IDMS Workload	Host Submodel	lt odel	Back-end Submodel	back-end Submodels	host/ Back-end Nodels	IDMS	Total
Represented	Model	VI	V2	V1	V2			
Application Environment								
Activity #1								
IDPEKS Application Software								•
Activity #2)				•)
• Cost System Hardware		•	•			•		•
• Cost System Executive Software Activity #3		•	•			•		
• Cost to PDP 11/70 Interface			•			•		•
Data base Environment (PDP 11/70)								
• Pup 11/70 to Host Interface					•	•		•
ACETUILY #4						(
PDP 11/70 Hardware PDP 11/70 Executive Software				•	• •	• •		• •
Activity #5) ()) (
DMS General Software DMS DM. Statements	•			•		•	• •	•
Data Base - Logical Structure							•	•
• Data base - Physical Structure				•	•	•		•

Figure 4.2-2 IPSS SIDPERS/IPMS Models

Phase II focused on the host environment. The base model was system Activity #1 as represented in the Phase I model with the backend treated as a black box. To this base model was added successively more detailed characterizations of the host system. The final model in this phase included the characterization for the backend interface hardware and software. A synopsis of Phase II activities is given below. Appendix A2 presents detail consideration of the modeling activities in this phase.

Phase III was concerned with the characterization of the backend system. Again the modeling proceded in an evolutionary manner. A sequence of models was planned which would add the PDP 11/70 hardware and operating system software to the IDMS workload model developed in Phase I. For these activities the front end was treated as a white box, i.e., a external source of back end requests. Specifically omitted was any attempt to model IDMS DML's or traversal activities for the SIDPERS logical data base (as defined via the IDMS DDL's). These latter functions were modeled using the IPSS/DBS facilities. A synposis of this phase is given below. Details are found in Appendix A3.

The goal of Phase IV was to combine the Phase II and III models into a single model. As mentioned at the begining of this appendix, a minor IPSS programming error, since corrected prevented their phase to be completed within the scheduled experimentation time line. Results of this modeling phase will be reported at a latter date.

Phases V was planned whose purpose was to model the IDMS DML's and the DDL described SIDPERS data base. A paper model was completed. However, execution was delayed until completion of the IPSS/DBS parsers. As with Phase IV the activity was completed after the termination of the modeling activities. This model is being reported in Mr. Brownsmith's disseration, titled "A Methodology for the Performance Assessment of Data Base Systems: An Extension to the Information Processing System Simulator Methodology", which is to be available Spring, 1979 from Ohio State University.

The goal of Phase VI was to integrate the Phase V model with the final Phase IV model. Again this activity has to await modifications to the current IPSS and IPSS/DBS.

A Synopis of the Phase II Activities (SIDPERS/HOST Models)

The SIDPERS application model was built assuming a computer free environment. The purpose of the Host submodel was to provide both the hardware configuration and the operating system of a host computer. Since no particular computer system had been identified as the host computer for the SIDPERS/IDMS project, a generalized host model was constructed. The elements of this model were characterizations of those operating system functions that support the SIDPERS application and the Host/Back-end interface. These include the job scheduling, task management and resource allocation, assumed to be available in the Phase I model, were now characterized. The hardware configuration that was modeled consisted of only these hardware components required to support the modeled host software.

The approach to construction of the host model followed the premise that the host operated in a reactive mode. The existing SIDPERS model was analyzed and when an operating system function was required, a module that characterized its behavior was added to the Phase II models. This procedure was followed for all levels of applications represented in the SIDPERS 1, and for the Host/Back-end interface needs. As a result a model was created of a host computer environment that satisfied the SIDPERS operating needs. The operating system functions included in the host model were characterized in a general, modular parameterized manner, such that at a later time, the mechanisms of the functions of a particular host computer could be easily incorporated into the existing model. Thus, although system activities 1 and 2 were communicating continuously, modifications to the latter could be made without affecting the existing SIDPERS application model. The SIDPERS/HOST submodels were tested by "black boxing" the back-end machine.

Synopsis of Phase III Activities (BACK-END/IDMS Models)

The IDMS model was built like the SIDPERS application model, assuming no particular computer environment. The purpose of the Back-end system

model was to provide the hardware configuration and the operating system of the back-end environment. Previous Army studies had designated the PDP-11/70 as the specific computer to be employed as the back-end computer. The functions of this computer system that needed to be modeled were the operating system functions to support IDMS and Host/Back-end interface routines. As with the host submodel, the back-end submodel was assumed to operate in a reactive mode, responding to requests for its services. The hardware configuration that was modeled was the complete USACSC PDR-11/70 configuration, recently moved with AIRMICS to the Georgia Institute of Technology in Atlanta.

The back-end models were built modularly, adding chracteristics of operating system or interfacing functions as they were needed. Although these characterizations were made specific to the AIRMICS PDP-11/70 computing environment, they were modeled in a parmeterized, modular manner so that if different PDP-11/70 strategies were to be substituted, the IDMS and interface submodels would not be affected. The IDMS/Back-end models were tested using the SIDPERS application model as the work load generator.

4.3 INSIGHTS GAINED FROM THE MODELING

The purpose of the modeling effort reported in this document was to demonstrate the feasibility and suitability of the Information Processing System Simulator (IPSS) with respect to modeling large complex information systems. The system modeled to demonstrate these capabilities was a host/back-end processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system functions of task management, job scheduling and resource allocation.

The IPSS methodology provided a modular, top-down approach to system analysis and model synthesis thus facilitating a solution to this complex problem. Several models were constructed, each of which

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reflected system details suited to the multi-stage development process which was adopted. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their integration into an overall complete model. The SIDPERS/IDMS workload model represents the software processes of the applications and IDMS in detail while characterizing the operating systems and computer hardware as a black box. From this basis model, two subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. These results are summarized in Appendix B. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood synopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicating both potential bottlenecks and under-utilized components.

These models were all internally verified. That is, through experimentation, the internal processing consistency was verified to be accurately reflective of the real world processes. Due to the lack of operational data, however, it was not possible to validate them through comparison to actual system performance. Thus, although some confidence can be placed in the relative values of the statistics reported in the Appendix B, no validation has been made with respect to the magnitude of the data values.

This project has demonstrated the suitability of IPSS to model complex systems while requiring a minimal time for model synthesis. The modeling effort as reported here required approximately four man months. The resulting models have proven to be versatile as well as

adaptable to changing requirements indicating that significant portions can be used for other application areas. For example, the DBMS and operating system processes as well as hardware facilities need not be modified when another major application area is modeled, thus resulting in a substantial reduction in future related modeling efforts.

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In addition, this project has served to focus future development work as well as to validate current goals. In particular, this project has demonstrated the need for simulation facilities specifically attuned to the characterization of integrated data structures as represented in IPSS/DBS. These models did not use these features since they were not fully operational, and thus the logical structure of the data base was charcterized at a more abstract level then desired. Furthermore, experimentation with record mapping policies was not attempted because of the complexities involved. Finally, several other research endeavors were identified. In particular, large complex models of computer networks could easily be constructed from models of individual nodes if a multi-simulator approach were adopted. Work is currently underway to allow multiple node simulations through extension of IPSS facilities and concepts. Extensions for modeling teleprocessing activities are also planned.

5.0 FUTURE EXTENSIONS TO THE IPSS

This research activity is just another step in a continuing IPSS research, development and engineering program whose goal is to create a comprehensive computer aided design facility appropriate to all stages in an information system's lifecycle. Three steps in the RD&E plan have now been completed. These include

- 1. 1967 1972: Initial Methodological and Simulator Development this activity resulted in the Auxiliary Storage System Simulator (AUXSYM). Details are described in "A Methodology for the Analysis of Auxiliary Storage Systems" available from University Microfilms.
- 2. 1972 1975: Extension to the methodology and complete redesign of the simulator. These activities resulted in the IPSS; details are reported in "A Methodology for the Performance Evaluation of Information Processing Systems" (DEL77) and "The Information Processing System Simulator (IPSS): Language Syntax and Semantics".
- 1975 1977: Continued methodological development to include the following
 - a. DBS specific features;
 - Assessment of system effectiveness with respect to competing user goals;
 - c. Addition of mechanisms to determine costing policies.
 - Assessment of system loading created by multi key query processing
 - e. Development of comparative evaluative techniques for simulation languages vis-a-vis performance assessment needs.
- 4. 1976 1979: Verification and Validation The purpose of these activities has been to demonstrate the effectiveness of the IPSS as a modeling facility. Validation has included paper models of different vendor hardware

software and file structures, models of idealized systems, models of existing complex systems (this report describes one such system.) A proposal by Chandler is designed to extend this activity into more specific validation activities where extrapolation can be validated.

As a result of the current and past research experiences, an extensive number of continuing research and experimental activities have been identified. They are presented here to illustrate the potential for the IPSS and its underlying methodology and to identify a means of activities which must be accomplished in order to achieve the ultimate research objective - that being to create a computer assisted design facility.

A. Extensions to the IPSS Methodology

The goal of these research activities is to continue the extension of the methodology's scope of applicability to (a) higher views of information system processes (b) other stages in the information process system lifecycle and (c) analysis and design processes. Five research activities have been identified for this area. They are:

- Telecommunications extension of the IPSS the permit characterization of data communication subsystems in a manner familiar to their structure and behavior.
- Requirement Analysis extension of the IPSS to interface with requirements analysis methodologies such as Tiechrow's PSL/PSA and TRW Corp's SCREM. (8)
- 3. <u>Sensitivity Analysis</u> development of a methodology which permits the assessment of models to identify performance sensitive components.
- 4. Estimating System Lifecycle Costs development of a methodology to examine IPSS models for the purpose of (a) estimating the number and complexity of components comprising the modeled system, and (b) using these data as input to lifecycle estimating techniques.

⁽⁸⁾ PSL/PSA is acrynon for Problem Statement Language/Problem Statement Analyzer SREM stands for System Requirement Engineering Methodology.

5. Automated Component Selection - formulates a methodology which is based on a predefined view of system components, and on one or more design criteria will automatically alter IPSS models thus identifying feasible designs for additional analyses by the modeler.

B. Extensions to the IPSS Language

The goal of this development activity is to expand the model synthesis and statistical display facilities of the current IPSS in order to reduce modeler effort. Five development activities have been identified. They are:

- Statistics Generation development of language extensions to cause the automatic gathering of standard statistics for sequential ("serially reusable") processes.
- Verification and Validation extends existing and develop new statements to support model verification and validation activities.
- 3. Operating System Characterizations extend existing and develop new facilities attuned to; (a) the modeling of information system job scheduling, resource allocation and task management functions and (b) characterizing main memory management, CPU interrupt activity and CPU instruction times.

C. Extensions to IPSS Simulation Facility

The purpose of these activities is to increase the ease of use of the IPSS simulator facility by the modeler. The following four items are proposed as extensions in this area:

1. On-Line Modeling - extend the processing environment for the IPSS to include; (a) on-line entry of IPSS language statements comprising an IPSS input stream, and (b) an on-line, real-time prompter facility to guide the modeler in the use of IPSS language syntax. The modeler will be able to switch between model stream definitions and stream processing as desired.

- 2. MACRO Library the IPSS macro and library facilities currently permits the modeler to retain models and/or model components for future reference. This activity would develop a set of generalized structures presenting the salient features of the following software:
 - a. Operating System: task management, job scheduling, resource
 allocation and memory management;
 - b. File Utilities: sort/merge, file copy, and opening and closing of files;
 - c. Data Base Management System: buffer management, high level DML's, schema structures.
- 3. <u>Graphical Display</u> the IPSS would be enhanced to permit graphical interaction between the system and the modeler. Initially this will involve only output of items such as:
 - a. Model statistics including frequency listographics.
 - b. Model structure showing the network of services described for the model and the data and file structures.
 - c. Dynamic model structure showing the sequencing of service invocation and their hierarchy during a task's execution.

Ultimately, the goal of graphics use is to prove a modeler/system communication interface which will permit interactive support of model synthesis and the display of model structure, model verification data, and model behavior statistics.

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APPENDICES

APPENDIX A

BACKGROUND TO THE MODELING EXPERIMENTS

The objective of this experiment was to provide a realistic problem environment against which the IPSS design facility (simulator and methodology) could be applied. The system modeled was chosen by Mr. Allan Curry of AIRMICS* specifically to test IPSS's design capabilities, ease of implementation, and adequacy and accuracy of the resultant statistics. The chosen system architecturally is of great interest to the U.S.A. Computer Systems Command as a replacement for exists BASOP systems. Section Al of this appendix discusses the problem environment including a full breakdown of the hardware and software characteristics of the system. Section A2 defines the IPSS models generated to characterize the system.

A1. THE PROBLEM ENVIRONMENT

The purpose of the following discussion is to provide a background to the system modeled. The system modeled evolved from a series of R&D efforts performed by the U.S. Army Computer Systems Command (USACSC) at the request of the Department of the Army, Director of Management Information Systems (DA DMIS) in direct support of their "Vertical Installation Automation Baseline" (VIABLE) Project. The VIABLE project was established to upgrade both hardware and software at Army data processing installations. Included in its scope is the competitive procurement of new ADP equipment, and the extension of functional software capabilities to include management information system requirements as well as interactive processing.

The initial USACSC research plan was developed in September 1975. Its purpose was to explore and demonstrate those technologies which would conceivably be utilized in Army data processing systems developed

^{*}AIRMICS stands for Army Institute for Research in Management Information and Computer Sciences.

over the next five to eight years. These areas included on-line processing, formal data base management systems, mini and micro computers and distributed data processing systems.

One aspect of the plan was to select a single functional area and build a baseline system which could be modified repetitively to add new features and migrate from one environment to another. The Standard Installation/Division Personnel System (SIDPERS), which processes Military Personnel actions was selected.

The computer system modeled was composed of a Front End Module (host machine) which processes SIDPERS transactions in a real-time interactive mode, attached to a Back End Module utilizing Cullinane Corporation's Integrated Data Base Management System (IDMS) to interface with the personnel data base. The backend machine modeled was the DEC PDP-11/70 System.

A1.1 SIDPERS

SIDPERS (Standard Installation/Division Personnel System) is the batch oriented multi-command personnel data processing system. It is a large system, consisting of approximately 800 types of transactions which perform update and retrieval operations on a personnel data base. However, since six of these transactions account for almost eighty percent of the total transaction activity, only these transactions were selected for incorporation into the baseline system. These transactions perform the following functions: adding a soldier to the data base, establishing duty status when a soldier arrives at an installation, changing the duty status and grade of a soldier, deleting a soldier from the data base on departure from an installation and acquiring data on a soldier and/or unit. These commands are denoted ADD, ARRIVAL, DUTY CHANGE, GRADE CHANGE, DEPARTURE, and INQUIRY respectively.

The SIDPERS test system required the creation of a small data base of sample records. Approximately 4,500 records representing 12 different

record types were selected and organized into an IDMS processable data base. The details of this data base and its organization are given in the next section.

The SIDPERS test system facilities for interactive transaction processing have been described by Schaaf (SCH77) (1). Basically, special terminal interface processing modules were designed which allowed a terminal user to enter the above transactions in either a batch or tutorial mode of operation. In batch mode, all the data for a complete transaction is entered at once. In the tutorial mode, the terminal user is prompted for every data item of the transaction. This requires an access to a VALIDITY record which contains the text for prompting. Items that fail to pass the validity check are returned to the terminal operator for re-submittal. For either mode of operation, once all the data items of a transaction pass this test, the corresponding application program is invoked. Essentially, there is one application program for each of the transactions listed above. As part of application processing, compatibility checks are performed. If a data item doesn't pass this check, several accesses are required to correct the error; namely an access to a COMPATIBLE record, and several accesses to VALIDITY records. Such errors cannot always be resolved and in these cases the transaction is terminated. Otherwise the application program issues a series of DML commands to the IDMS Modules in order to store, modify, and retrieve records in the data base. During the IDMS Processing of each DML, application processing is suspended. When the application is completed, control is returned for further interaction.

Al.2 Prior Analysis Activities of SIDPERS

The data required to characterize these activities was provided from a previous study of VIABLE transactions (SHA77), (See Tables A1.2-1 and A1.2-2 for further specifics). This document provided the following data for each of the six types of transactions:

⁽¹⁾ NOTE. All references are found at the end of the main body of this report,

Table Al.2-1 The Results of the SIDPERS Interaction Test

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	No. of entered data items	To , on bata transactions	No. of Cancelled transactions	lo. oV ni sansactions in Direct mode	No, of transactions in tutorial mode	Vo. of validity errors	% of data items with val. errors	No. of compatibility errors	% of data items with comp, errors
ARRIVAL	6	158	25 / 16%	93 / 26%	65 / 41%	59		31	
DEPARTURE	12	193	34 / 18%	111 / 57%	82 / 43%	09		101	
DUTY CHANGE	80	331	23 / 7%	208 / 63%	123 / 37%	16		170	
GRADE CHANGE	10	52	11 / 21%	31 / 60%	21 / 40%	. 15		34	
	70	9	%0 / 0	%0 / 0	6 / 100%	17		9	
INQUIRY	3	38	%0 / 0	31 / 82%	7 / 18%	2		12	
TOTALS		778	93 / 12%	474 / 61%	304 / 39%	214	2.9%	354	4.8%

Table Al.2-2 Average Entering Time

Transaction Type	Enterin Per Tra	Entering Time Per Transaction	Trans	Transaction Characteristics	Entering Data	Entering Time Per Data Item
	Tutorial Min:Sec	Direct Min:Sec	No. of Entered Data Items	Errors (about 8%)	Tutorial (sec)	Direct (sec)
ARRIVAL	5:23	2:49	6	-	32.3	16.9
DEPARTURE	6:03	2 :55	12	1	27.9	13.7
DUTY CHANGE	3:58	1:40		1	26.4	ו.וו
GRADE CHANGE	5:55	1:48	10	1	32.3	3.8
ADD	25:57		70	9	20.5	
INQUIRY	1:23	0:54	က	•	17.72	18
AVERAGE OVER A	AVERAGE OVER ALL TRANSACTIONS	NS.			27.9	13.9

- 1. the relative frequency of occurrence,
- 2. the percent direct mode,
- 3. the number of data items per transaction entry,
- 4. the percent data item validity error,
- 5. the percent data item compatibility error.
- 6. the time for direct mode data items.
- 7. the time for turorial mode data items, and
- 8. the percent of transactions canceled.

This data was incorporated into the IPSS model of SIDPERS/IDMS through Fortran DATA statements and was used to "drive" the model through representative terminal sessions. In addition, this document contains for each transaction, a list of the DMLs that are invoked and their sequence within the application program. This data was used to invoke lower level DBMS processes as described in the next section.

Al.3 IPSS Model of SIDPERS

The IPSS SIDPERS/IDMS transaction model was designed to reflect the arrival validation and DML activities of each of the six previously identified SIDPERS transactions. Figure Al.3-1 identifies the major functional areas of this model. Each functional area is represented by an IPSS service or services within the model. The function of each of these services are defined as follows:

System Loading

Within the IPSS model, the arrival of a transaction at a work station is represented by an exogenous event (defined within the Exogenous Event Stream Component). This event was scheduled to occur every 5.89 minutes plus or minus 5 minutes. This is based upon the following assumptions and parameters. First, the data base represents 203 soldiers; second, that there are an average of 8 transactions per month per soldier; third, that data base updating activities through

terminal interaction is possible 8 hours per day, 20 days per month. These figures were used to derive the transaction arrival event frequency of 5.89 minutes. The deviation of 5 minutes was selected arbitrarily to represent variation in the arrivals. In order to model a larger data base or different arrival patterns, a single statement need be changed in the procedure associated with the event. This fact is documented within the procedure and the statement identified for ease of modeler experimentation.

Transaction Arrival and Sequencing

This service represents the arrival of a transaction at a work station. The model assumes a single terminal and hence the transaction may be queued if the previous one has not completed. The model also assumes a FIFO (first-in/first-out) queueing discipline. The time that a transaction waits in this queue is a statistic automatically provided by IPSS. When a transaction is selected for processing, it is removed from the queue, the transaction Selection Service is invoked, and the service waits its completion. The total elapsed time from transaction selection to completion is also maintained by IPSS. Upon completion, the next transaction is selected from the queue and the service becomes idle. Idle time and busy period statistics are also kept. All the statistics mentioned above are maintained for each of the services and will not be further documented in this section.

This service can easily be modified to reflect an arbitrary number of terminals simply by establishing the maximum number of terminals permitted to be in operation at once, and keeping track of the current number. Only those transactions which arrive when the current number is equal to the maximum number need by queued. An infinite number of terminals can be modeled simply by not comparing these two variables and hence never queueing a transaction.

Transaction Selection

This service selects the type of transaction that has arrived and associates with it a set of parameters. The selection is made by a probability distribution function based on relative frequency data provided by the Schaaff (SCH77) study. Once this selection is made, this service invokes the terminal interaction and validation service and waits for its completion. When the processing associated with this lower level service has completed, the Transaction Selection Service, determines if the transaction has been canceled (via a return code) and, if not, invokes the proper Application Service and waits for its completion.

Terminal Interaction and Validation

This service represents all the activities associated with the entry and validation of a transaction by a terminal operator; and can be conceptually divided into two parts, namely those activities that occur in obtaining valid transaction through terminal interaction in direct mode and those activities required to process and complete the transaction in tutorial mode. These processes are both characterized through tabular data and probability distribution functions. Input from the terminal, for example, is represented by advancing the simulator clock. In separate models, facilities representing the hardware and software processes are characterized and thus these activities are more accurately measured. The tutorial input mode requires several more parameters than the direct mode to characterize the events. The final activity of this service is to determine whether the transaction is to be canceled. This is communicated to the invoking service by a parameter value.

⁽²⁾ These parameters represent the transaction characteristics and are derived from the data in Tables Al.2-1 and Al.2-2.

Application Services

The model contains a set of six functionally similar services, each of which represents the processing performed by an application program for one of the transactions. These services are all invoked by the Transaction Selection Service and they all at some time invoke the IDMS Interface Service. The processing performed by these transactions is represented by a series of DML invocations to this lower level service. The internal data manipulation activities are not represented in the model since their impact on performance is assumed to be minimal. The DMLs are identifical in type and sequencing to those reported by Schaaf (SCH77); their format is further discussed in the next section.

While not all performance evaluations characterize application processing to the detail represented in this model, this approach was deemed best for this model for the following reasons: (1) the data was readily available, (2) the mechanics of model construction was facilitated by IPSS language modeling features and (3) the model architecture emulates the processing performed by the real system. Thus, when the model is used in the future it will be evident what processes are being modeled.

Within IPSS, a service is dynamic facility and as such statistics are automatically collected for its use. For each of the services identified in Figure Al.3-1, utilization statistics are produced. These include: busy period, idle period, concurrency, seizure, and transaction transmit time statistics. In addition, queueing statistics are also produced. However, because of the nature of the model, only the Transaction Sequencing Service has queueing associated with it. These queueing statistics include: busy period, idle period, queue length, queue entry, and transaction transmit time statistics.

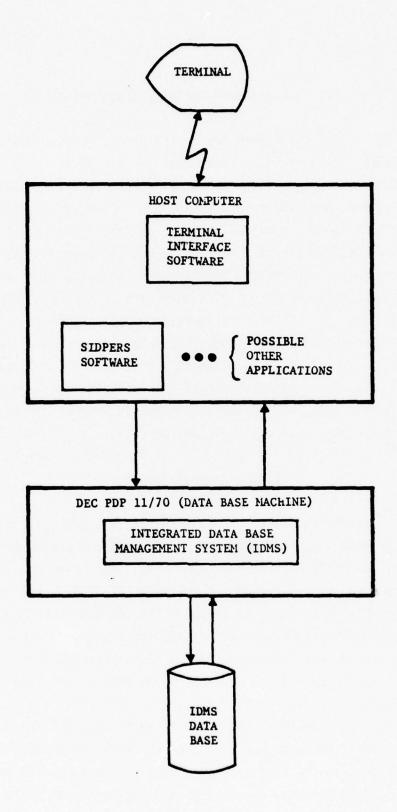


Figure A2-1 Experimental Version of VIABLE SIDPERS Processing Environment

A2. THE HARDWARE/SOFTWARE ENVIRONMENT

The configuration established as the protype for the VIABLE SIDPERS project employs the host/backend concept (Figure A2-1). The IDMS SIDPERS data base would reside on and be controlled by the backend machine. The purpose for this structure is to reduce the load on the host machine in order to extend the life of the mainframe. Previous studies have not designated a particular host machine for use by USACSC. Currently, the VIABLE SIDPERS test data base is maintained on the USACSC HCC IBM 370 system, with both host and backend machines residing there. One of the most immediate applications of the models developed during this project is to evaluate different host environments, in particular, the IBM 370 environment versus the PDP-11/70 environment. Thus, the goal of the modeling effort with respect to the host machine is to develop a model of a machine independent generalized host machine, allowing for each substitution of environments.

A2.1 The Host Machine

To accurately model the IBM 370/165 host was deemed to be too time consuming a task with regard to the objectives of the project. Therefore, an idealized host environment was identified which paralleled the hardware and software needs of the VIABLE/SIDPERS. Host hardware is shown in Figure A2.1-1. It included the terminal, CPU, secondary storage equipment (I/O Processor, disk control unit and eight IBM 3330 type disk units), and a data channel to the backend machine. Also shown in this figure are the hardware components assumed for the backend computer.

Host software functions in addition to the SIDPERS are shown in Figure A2.1-2. The major components of the application/host environment are the teleprocessing (TP) interface with the user, the TP interface with the backend machine, the host's disk I/O software, and the SIDPERS

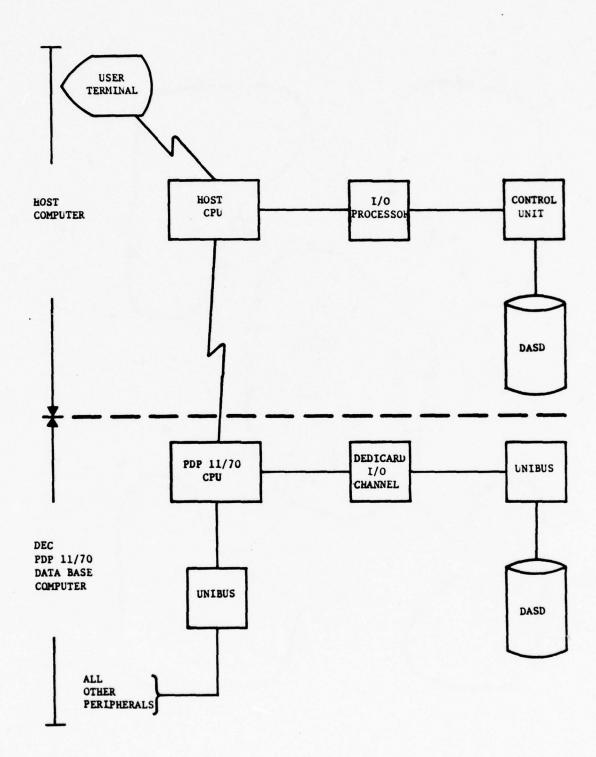
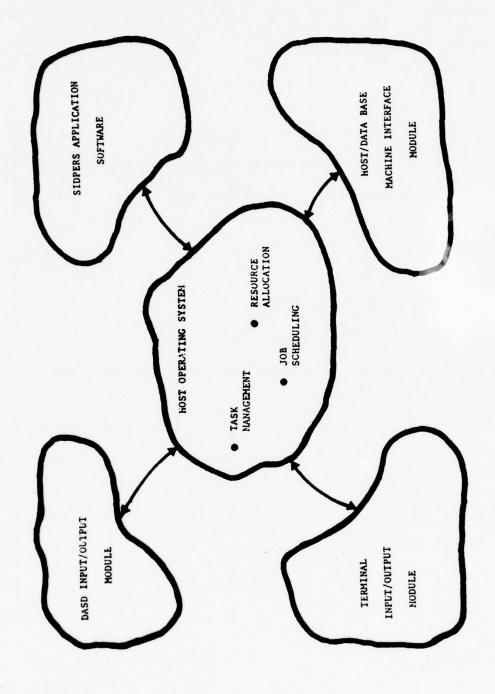


Figure A2.1-1 Modeled Hardware Configuration



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Figure A2.1-2 Modeled Host System Resident Software

application itself. The generalized host operating system (OS) is assumed to be reactive in operation, i.e., executing only to respond to the needs of the other software components.

Since VIABLE SIDPERS is interactive, communication with the user is maintained throughout each step in the process. In this capacity, the host performs the function of scheduling the communication, managing the TP line and sending/receiving the messages. After the initial log-on procedure the host operating system schedules the initiation of SIDPERS processing. The input and validation of SIDPERS data involves, in addition to the host's TP functions, disk I/O for the purpose of retrieving error messages for transmission to the terminal.

The crucial element of this model is the liaison between the application in the host machine and the DBMS in the backend machine. The DML request must be transmitted to the backend for processing and the retrieved record must be received. To accomplish this, the host supports an interface subsystem which formats the DML request, reformats the resultant record, handles the asychronous transmission and reception of messages to and from the backend, and controls the necessary devices for this communication link. It is assumed that each of the host modules rely on the basic operating system functions of task management and resource allocation. As in the actual system, the performance of these host functions is transparent to the processing of the application.

In order to make the model of the host environment machine independent, several conventions were employed. When scheduling was required (job, disk, or I/O), a FIFO queue was assumed for that host module. Resource allocation algorithms for TP and IOCS needs were also assumed to be FIFO. Any algorithm, however, can be substituted by changing only the host scheduling services without effecting the model's application services.

A2.2 The Backend Machine

The particular computer system selected by USACSC for the backend machine is the PDP-11/70. As a result of this prior decision a more

detailed and specific model could be built. The hardware configuration for the existing USACSC AIRMICS PDP-11/70 was modeled. Secondary storage characterization conformed to the IDMS a fixed page environment, i.e., the physical data records of the data base with a page size of 1,024 bytes and a direct file organization. Figure A2.2-1 shows the hardware configuration modeled.

The crucial aspect of modeling the backend machine for this project was to identify and characterize the interactions between the PDP-11/70 operating system and the IDMS applications. The main application oriented functions to be coordinated by the resident PDP-11/70 OS are illustrated in Figure A2.2-2. The internal functions of the backend machine are essentially the same as the generalized host except for UNIBUS management. The concept of the UNIBUS (a bidirectional data path through which all resources, including CP and memory, communicate) is central to PDP-11/70 processing. But its effect on processing is transparent to the application. The backend PDP-11/70 has an interface in the host. Because the PDP-11/70 is not dedicated to data base applications, and because USACSC has plans to support more than one data base per backend machine, each invocation of IDMS is treated as a separate subtask. The problem of concurrent updating by multiple users, however, forces the need for a single entry point or interface, called CAMP (Central Access Monitor Program). Even though entry to IDMS is controlled via the application, the actual operation of IDMS necessitates many requests for task management, resource allocation and I/O control. The model of the PDP-11/70 operates, as was assumed for the model of the host machine, in a reactive mode.

The object of the IDMS processing, retrieval of the specified physical record, necessitated the scheduling and use of I/O hardware which comes under the preview of the PDP-11/70 analog of the host's disk I/O module. Passing retrieved records to the host requires the aid of the OS to handle the sequencing of completed IDMS tasks and de-allocation of resources.

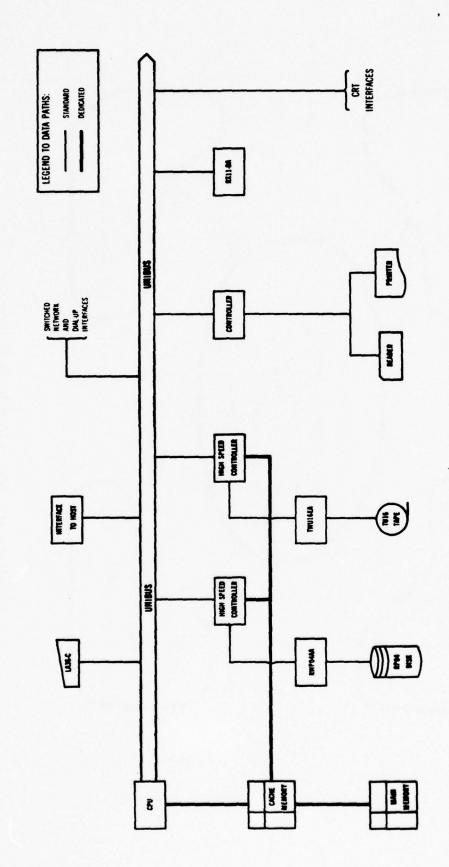


Figure A2.2-1 USACSC PuP 11/70 Configuration

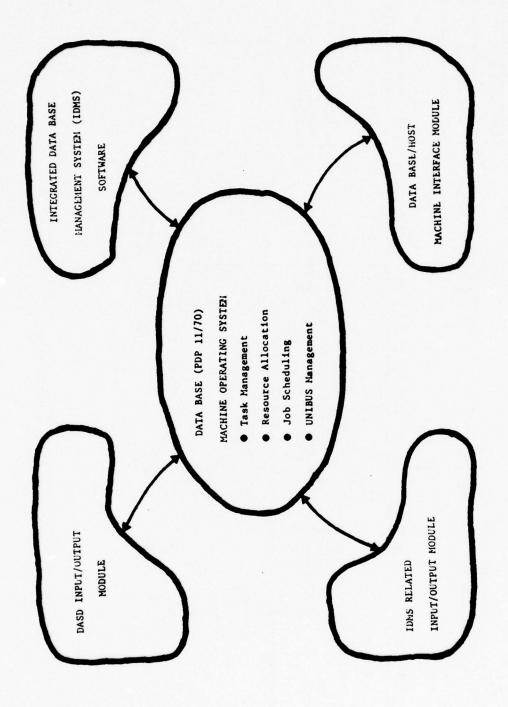


Figure A2.2-2 Hodeled Data base System Resident Software

The job and task scheduling disciplines employed by PDP-11/70 OS are priority queues based on the immediacy of use of the UNIBUS and were so modeled. All tasks performed by the PDP-11/70 priority scheme resulted in, essentially, FIFO queues.

A2.3 Integrated Database Management System - IDMS

The Integrated Database Management System is a network oriented implementation of a subset of the 1971 CODASYL Database Task Group specifications. The following activities were involved in making the conversion from sequential, batch processing to the IDMS on-line integrated database. First, the architecture of the integrated structure of the database was determined. Each of the files identified in Figure A2.3-1 became (in IDMS terminology) "record types" whose interconnections are represented by "sets". Second, the COBOL application programs for the six transaction types was rewritten to reflect IDMS processing. The major impact was in the incorporation of a centralized description of the database and in changing the I/O statements to standard IDMS Data Manipulation Language (DML) commands.

IDMS performs the following functions in response to an application DML. First it validates the request using the sub-schema definitions. Next it translates the request into a virtual page reference, determines if the page is present in its internal buffer, and requests the page from the underlying file management system when it is not already present. Finally, it places the subschema record into the application's work area and returns control to the application.

IDMS has its origins at the B.F. Goodrich Company where it was initially developed. The Cullinane Corporation (Wellesley, MA) now maintains product responsibility. IDMS operates on IBM 360/370 hardware. In addition, it is available through the Digital Equipment Corporation as DBMS-11 which operates on PDP-11/70 hardware.

The following paragraphs briefly introduce the major IDMS system components which were incorporated into the IPSS model (CUL76):

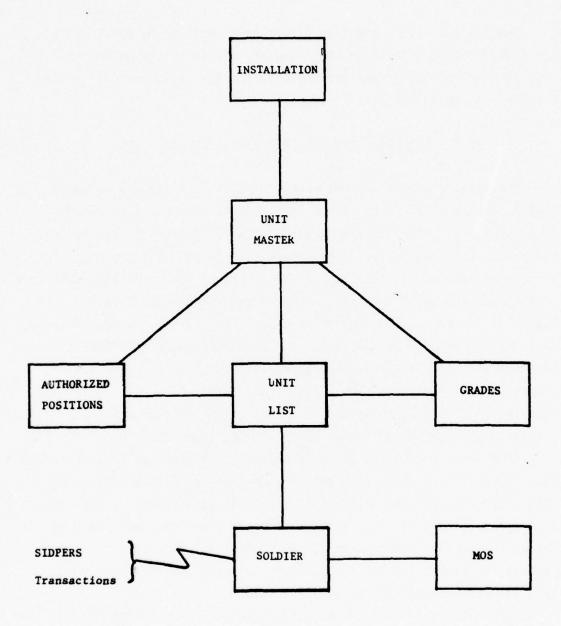


Figure A2.3-1 Modeled IDMS Schema for SIDPERS Data Base

<u>Data Description Language (DDL)</u> -- A facility for describing data independent of the programs that process it. Within IDMS the DDL consists of three parts:

- Schema DDL -- provides a complete description of the database including names and attributes of all data items, records, sets, and areas.
- 2. Device Media Control Language (DMCL) -- provides control over all the disk files which are to be on-line.
- Subschema DDL -- provides a description of those data items, records, sets and areas to be used by one or more application programs.

<u>Data Manipulation Language (DML)</u> -- A facility for accessing base. It provides a wide range of commands for requesting database services which store data in the database or returns data to the user's program.

Schema Description for SIDPERS

The STDPERS test database was designed to support the six transact—types identified earlier. The overall system architecture reflects twelve record types organized into four IDMS areas. These areas reflect functional differences in the data that is stored and its use. PERSONNEL area contains the individual solder and unit records of an installation. The MSCAREA contains diagnostic and error messages. The ADAREAL and ADAREA2 Areas contain additional personnel data. The overall structure is shown in Figure A2.3-2.

The record types belonging to each of the areas are identified in Table A2.3-1. The record occurrence and record size attributes of each of the record types are also identified in this table. The personnel area is of primary interest since most of the SIDPERS transaction activity references it record types and sets. This area consists of 3,909 records distributed among 7 record types. Although the GRADE records account for 68% of the total, the IDENT records (5,2% of the

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Figure A2.3-2 SIDPERS Data Diagram

Table A2.3-1 Record Occurrence Characteristics

	Record	Occurrence	Recor	rd Size
Record Type Within Area	Number	% of Total (within area)	Length* (Bytes)	Percent Used of Space Available Within Area
PERSONNEL AREA				
INST	1	0.0	48	.01
UNIT-MSTR	83	2,1	52	.79
IDENT	203	5.2	120	4.12
UIST	468	12.0	84	6.82
GRADE	2,656	68.0	20	11 .78
POSNO	403	10.3	140	9.45
MOSE	95	2.4	236	3.67
ADAREAL AREA				
ADD-IDENT	200	70,2	284	23.13
ADD-UNIT	85	29.8	552	18.85
MSGAREA AREA				
COMPATIBLE	155	62.0	296	14.93
VALIDITY	95	38.0	512	15.65
ADAREA2 AREA				
AALOC	9.5	100.0	96	,15

^{*}Includes IDMS Overhead of 8 bytes per record.

total) are most frequently accessed. Most of the transactions are oriented toward updating an individual soldier's data and thus the most common entry point is the IDENT record. This is accessed via the CALC set using the soldier's last name.

Figure A2.3-2 identified the organization of the personnel area record types into sets. Ten distinct set types (excluding the CALC sets) are employed to provide the inter-record traversal capabilities required by the SIDPERS transactions. Table A2.3-2 identifies these sets and displays their characteristics. The sets are identified by both an "external" name which is the set name used in the SIDPERS subset, and also by an "internal" name which is the set identifier within the IPSS models. For modeling purposes, two data items per set type were derived from analysis of the sample data base. These were:

- the probability of an owner record type having a non-null set of member records associated with it (i.e., the probability of a set occurrence), and
- 2. the number of member records in a set occurrence for each set type. Table A2.3-2 also identifies these data items for each set of the PERSONNEL AREA. Although the table shows the average number of member records, for modeling purposes probability distribution functions were used to more accurately characterize the database.

The records of the SIDPERS data base reside within specially formatted IDMS pages. These IDMS pages are virtual pages since they represent the data as known by IDMS and as such need not be reflective of secondary storage record organization. The page size itself is a function of the IDMS processing environment. Interactive SIDPERS with IDMS operating on IBM 370/165 equipment uses a page size of 3,156 bytes, whereas the page size for IDMS operating on a PDP-11/70 is projected to be 1,024 bytes. The test version of the SIDPERS data base has been loaded onto IDMS pages as shown in Table A2.3-2. The Personnel Area occupies 200 pages and is 36.6% full. The IPSS simulation model of this data base assumes a PDP-11/70 environment as the number of pages and page range was adjusted to keep the percent utilization the same while changing the page size to 1.024 bytes. This data is also reflected in Table A2.3-3.

Table A2.3-2 Set Occurrence Characteristics

Set Name		Record Type Participation		Probability of Set	Average Number of
External ID	Internal	Owner	Member	Occurrence Given Owner Record	Member Records
INST-UNIT-MSTR	INUM	INST	UNMA	1.0	83.0
UNIT-MSTR-POSNO	UMPOS	UNMA	POSN	.9	5.7
UNIT-MSTR-UIST	UMUI	UNMA	UIST	1.0	2.6
UNIT-MSTR-AWOL	UMAW	UNMA	UIST	0.0	0.0
UNIT-MSTR-GRADE	UMGR	UNMA	GRADE	1.0	32.0
POSNO-UIST	POUI	POSN	UIST	0.0	0.0
GRADE-UIST	GRUI	GRADE	UIST	.94	1.6
IDENT-UIST	IDUI	IDEN	UIST	1.0	3.0
MOSC-PMOS	PMID	MOSC	IDEN	.86	1.0
MOSC-SMOS	SMID	MOSC	IDEN	.57	2.0

Table A2.3-3 Page Utilization Within the SIDPERS Areas

	Page Size	= 3,156 Bytes	Page Size = 1,024 Bytes		
Area	Number of Pages	Page Range	Utilized	Number of Pages	Page Range
Personnel	200	1-200	36.64	617	1-617
ADAREA1	80	201-280	41,98	246	618-863
MSGA2EA	100	301-400	30.58	308	925-1233
ADAREA2	20	281 - 300	,15	60	864 - 924

DML Description for SIDPERS

A previous study of a VIABLE transaction has specified the number and type of IDMS DMLs that are issued as well as their sequence within the transaction (SHA77). These DML have been classified by type as shown in Table A2.3-4.

The function of each of these DML Commands is as follows:

FIND: Locates the selected record in the data base and returns its identity to the application program.

<u>OBTAIN</u>: Locates the selected record in the data base and returns the record as well as its identity to the application program.

MODIFY: Causes the specified record to be updated.

STORE: Causes the specified record to be stored in the data base.

<u>ERASE</u>: Causes the specified record to be physically removed from the data base.

CONNECT: Causes the specified record to be inserted into an identified set.

<u>DISCONNECT</u>: Causes the specified record to be removed from the identified set.

Table A2.3-4 SIDPERS IDMS DML Types

Category	DML
FIND/OBTAIN	
	FIND record-name RECORD FIND NEXT record-name RECORD of set-name SET FIND OWNER RECORD OF set-name SET
	OBTAIN record-name RECORD OBTAIN NEXT record-name RECORD of set-name SET OBTAIN OWNER RECORD OF set-name SET
MODIFY/STORE/ERASE	
	MODIFY record-name RECORD STORE record-name RECORD ERASE record-name
CONNECT/DISCONNECT	
	CONNECT record-name TO set-name DISCONNECT record-name FROM set-name

Within the IPSS model of VIABLE SIDPERS, each IDMS DML is represented by statements which invoke a lower level service and wait for its completion. Four parameters that are required to pass the DML information. These are:

- DML Type (FIND, OBTAIN, etc.);
- Qualifier (NEXT, OWNER, etc.);
- 3. Record Type Identified (ADDI, IDEN, etc.); and
- Set Type Identifier (GALC, GRUI, etc.).

A3. MODELING USING THE IPSS DESIGN FACILITY

A3.1 A Functional View of the Modeling Design

The purpose of the host and the backend submodels is to provide specific computer system environments for the SIDPERS application and IDMS submodels. The relationship between these submodels can be seen in Figure A3.1-1. The SIDPERS submodel was constructed by characterizing the functional SIDPERS processes required by a SIDPERS transaction from the time it enters the host system until the completion of processing. These functional processes included such activities as input validation, backend interaction, and diagnostics transmissions. The SIDPERS submodel characterized these processes, however, in a computer system free environment, concerning itself only with the impact of these various processes on SIDPERS processing. Similarily, the IDMS submodel was constructed with no assumed computer system environment. Functions such as buffer and page management, DML invocation and host interfacing were characterized as part of the IDMS submodel, at the level of complexity of the IDMS request. Both the host and backend computing environments were treated as black boxes in these models. Thus, the initial preliminary model had the SIDPERS submodel interacting with the IDMS submodels (arrows A & B).

The purpose of the host and backend submodels were to provide a specific computer environment, including hardware, software and data elements, within which the SIDPERS and IDMS submodels could operate. The operating system functions that had been assumed to be operational in the application oriented submodels, but had been aggregated, were now characterized. For example, transmitting SIDPERS diagnostic messages to the user at the terminal involves accessing the message from the host's disks (including I/O scheduling, disk allocation and I/O operations) and the actual transmission (which requires TP scheduling, TP line allocation and I/O operations). With respect to the IDMS submodel, the function of page replacement, for example, involves task scheduling, resource allocation, and memory management. It is the characterization of these support functions that the host and backend submodels provided.

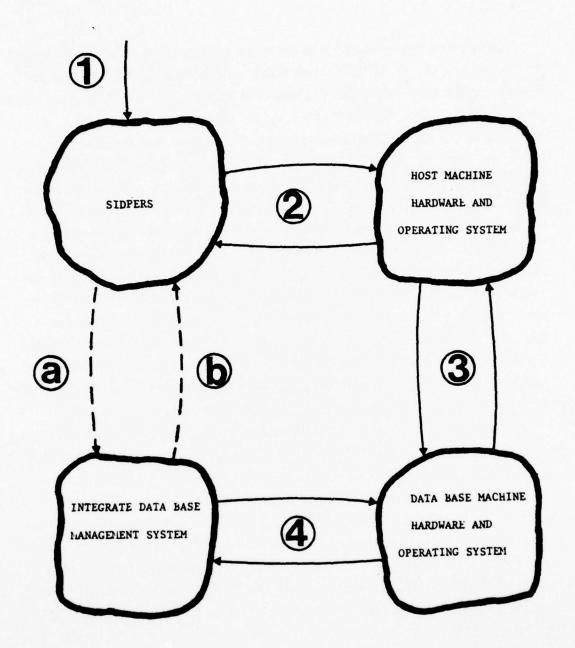


Figure A3.1-1 Relation of Submodel to Total Model

The processing relationship of these submodels is also demonstrated in Figure A2.1-1. A SIDPERS transaction enters the system through the SIDPERS application submodel (1) and continues to operate in this submodel until it requires a host operating system function. At that time it triggers off a host function to perform the needed operating system function, such as task, job or resource management (2). Soth submodels continue processing other transactions and events that may be occurring simultaneously. In essence, then, the host submodel operates in parallel with the SIDPERS model, satisfying its operating system needs and coordinating the execution of concurrent SIDPERS processes. The host submodel is essentially in a reactive mode, residing in a continual wait state until awakened by the SIDPERS application submodel requesting a host service. Communication between the host is accomplished between the host and backend submodel (3). The backend submodel then initiates the appropriate IDMS service to satisfy the DML request sent from the host (4). From then on, however, the backend and IDMS submodels operate in a parallel as does the host and SIDPERS submodels.

Previous studies of the VIABLE/SIDPERS project have not designated a particular type of computer system as the host machine. Thus, the model that was developed has been designed to be a general host machine model, including those operating system functions that would be necessary on any machine that as host. The structure of the host model, however, is such that the particular software and hardware of any computer system can be substituted for those in the initial model without modification of the existing SIDPERS application submodel and little modification of the host submodel. Currently, the test bed SIDPERS/IDMS system resides on an IBM 370/165, both the host and the backend. Thus, although the structure of the host submodel in this project is general, the particular elements are specific to the USACSC HCC IBM 370/165. The backend machine, however, has been chosen to be a PDP-11/70 and the resultant backend model will reflect the operating system philosophies of this system.

In modeling these two types of environments, a similar approach was taken. It is assumed in the IPSS methodology that operating systems,

in general, perform certain types of services regardless of the particular machine. As Figure A3.1-2 illustrates, although an operating system may be servicing many different applications simultaneously, the coordinating and servicing of these applications is completed through a stage set of underlying operating system functions. The types of system functions performed are job scheduling, task management, resource allocation, I/O supervision and telecommunication control. Thus, both the host and the backend submodels were constructed by characterizing these functions in general and then particularizing them to the specific machine, IBM 370/165 and PDP-11/70. The structure of these submodels, however, is such that at any later time the hardware and/or software can be changed to represent another system.

It should be emphasized that these characterizations of the computer systems were not dependent on the applications to be run on them, i.e., SIDPERS and IDMS. A future application of these submodels is to change the loading on the computer submodels to include other, possible other SIDPERS, applications to determine the effect on SIDPER/IDMS processing. Thus, the scheduling and allocation policies that are coded for these submodels do not reflect the particulars of either SIDPER or IDMS and are, essentially, general task and resource managers.

A3.2 IPSS Model of Host Computer

The approach taken in using IPSS to model this computer system is to first model the static elements of the system, the hardware and the data and then model the dynamic software elements. The host computer that was modeled was the USACSC HCC IBM 370/165. To construct the hardware configuration of this machine, only those elements that were impacted by the SIDPERS application were modeled. Specifically, they were the secondary storage devices upon which the IDMS error message files were kept, the backend interfacing equipment, external terminals through which USACSC personnel invoke and interact with SIDPERS, and the connecting telecommunication equipment. Other hardware typical of an IBM

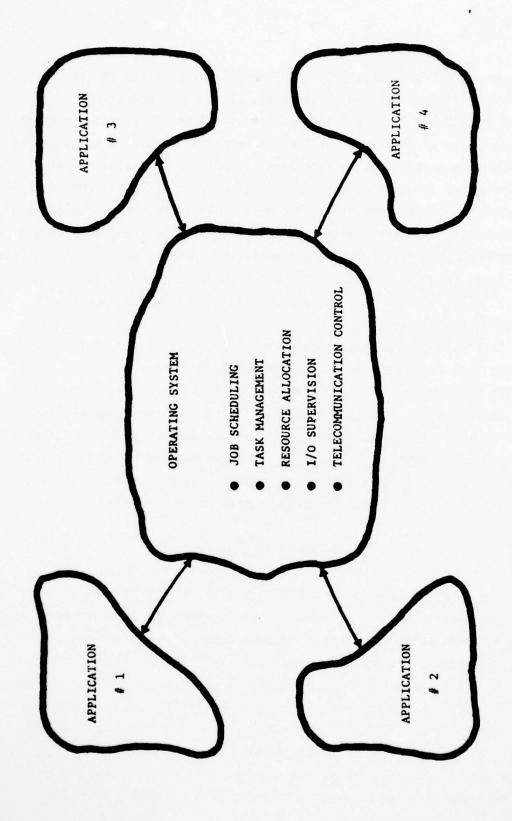


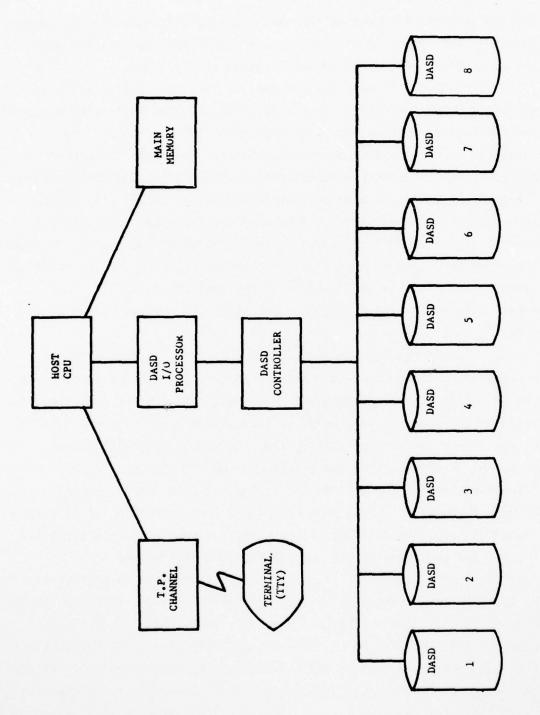
Figure A3.1-2 Modeled Operating System Functions

370/165 were not included as they would not be referenced in this model, but at a later point in time when other USACSC applications are added to the host model, additional hardware can be easily added.

Figure A3.2-1 illustrates the initial hardware configuration that was assumed for the host, IBM 370/165 computer. The particular secondary storage devices assumed were a bank of eight IBM 3330 DASDs, their associated control units, and single channel. A single CPU with 2 megabytes of main memory were assumed to characterize the host computer. The external terminals were assumed to be similar to IBM 2741 and the communication line between the host computer and these terminals was assumed to be a 2400 baud line. It should be noted that in an IPSS model the characteristics of anyone of these hardware devices can be modified without requiring the modification of the rest of the model. Thus, experiments evaluating different combinations of devices can be easily made.

The second step in constructing an IPSS model is to define the secondary storage data files. With respect to the SIDPERS application, the only data files associated with the host computer are the files of validity and compatibility error messages. The validity error file is accessed once for a compatability error by the SIDPERS application software. Although there was little data on the size and internal characteristics of these files the following was assumed: logical record length of 80 bytes, physical record length of 4000, (to fit into the VM OS page size of 4096), 1000 validity records, 200 compatibility records and placement of the two files on the same volume.

The crucial aspect of this and any IPSS model is the characterization of software processing. It has been mentioned that the operating systems of both the host and the backend computers were assumed to be in a reactive mode, responding to the needs of their respective applications. Thus, in order to determine the operating system functions that needed to be modeled in the host (and backend) IPSS submodels, the processing requirements of the SIDPERS (and IDMS) applications must be examined. The requirements for the host can be seen by examining the functional



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Figure A3.2-1 Assumed Host Hardware Configuration

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steps that a SIDPERS transaction goes through that were detailed in Section A1. Summarizing, there were five basic steps as shown in Figure A3.2-2: log-on, SIDPERS scheduling, data input and validation, IDMS/backend interaction, and log-off. At each of these steps particular operating system functions are required at verying levels of complexity.

To satisfy these needs, certain salient operating system functions were identified. These were: the scheduling and transmission of data to/from the user, the scheduling of a SIDPERS application within the host, the communication requirements of interactive data input and the necessary resource allocation and I/O scheduling to access the error message files and the scheduling transmission and repetition of messages to/from the backend computer. Underlying all these processes is the synchronization and coordination of the SIDPERS application functions and the host operating system functions which may be executing in parallel, as provided by general task management and resource allocation.

The relationship between the IPSS services that characterized the SIDPERS application and the IPSS services that performed the host functions are shown in Figure A3.2-3. Being that SIDPERS is an interactive system, constant communication between the user and the SIDPERS system, through the host, must be maintained. Thus, each of the five SIDPERS application steps, in some way, require the scheduling and transmission of information to or from the user. In the resulting IPSS submodel this interface was handled by the service XTCAM, which acted as scheduler for TP to the user. It was awakened by a SIDPERS module and returned to the wait state. The requesting SIDPERS module was then in control of the use and disposition of the TP line. The second step was to specifically scheduled the SIDPERS application for execution in the host system. This was performed by the IPSS service JOBSCH which controlled the allocation of main memory and central processor and placed the incoming SIDPERS application in the ready queue. It is general, though, so that other non-SIDPERS jobs could also be scheduled,

The third SIDPERS application step, data input and validation, was characterized in the SIDPERS application submodel by the services VALID,

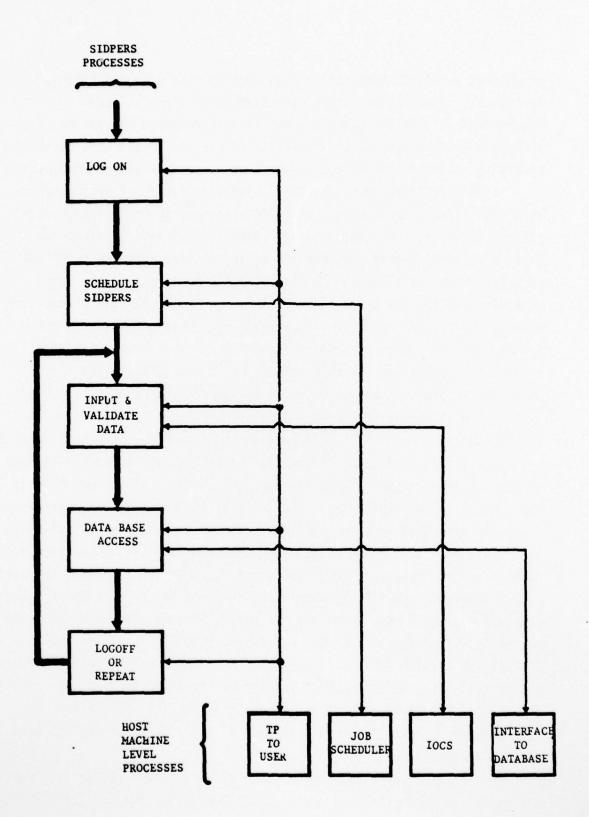
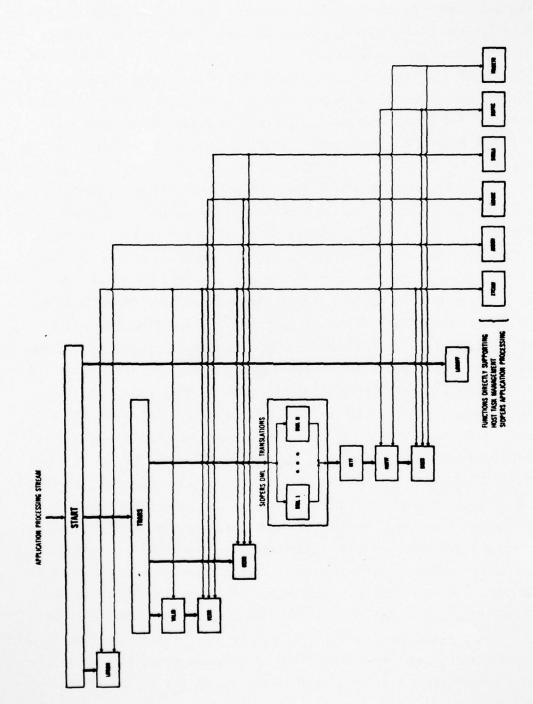


Figure A3.2-2 Relation of Host Machine Processes to SIDPERS Processes



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Figure A3.2-3 IPSS Realization of Host/SIDPERS Model

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A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS--ETC(U)
SEP 78 T 6 DELUTIS
DAAG29-77-G-0203 AD-A069 543 ARO-15356.1-A-M UNCLASSIFIED NL 2 OF 3 AD AO 69 543

VERR and CERR, representing, respectively, input process, validity error accessing and compatibility error accessing. Each of these modules involves communication with the user, thus, necessitating interaction with the host service XCAM. VERR and CERR also require the access of the host disk storage subsystem. This process is regulated by a disk scheduler and device allocator, XHDSK. The actual I/O commands are performed by an IOCS routine labeled DISKA in the IPSS submodel of the host.

The fourth step in a SIDPERS application is the DML processing, including the interface with the backend computer for the purposes of sending the DML request to the backend, and receiving the retrieved record from the backend IDMS. The SIDPERS submodel services for each transaction type and the service INTF determine the particular DML required for the transaction and perform any pre/post DML processing. The actual interface to the backend computer is performed by HINTF. This service characterizes the packaging of the DML request for transmission to the backend, buffer management (BUFALC), scheduling of the TP line to the backend (HOBETR) and allocation of that line. In the preliminary submodel of the host, the backend processing was black boxed and service BACK characterized this implementation. BACK also controlled the deallocation of resources (buffers and TP line). The final SIDPERS step of log-off parallels the log-on procedure, performing allocation (now deallocation) of resources and job scheduling.

In order to characterize the host operating systems functions as they would normally operate, asynchronously, the IPSS submodels had to operate (or simulate) asynchronously. The interface between the host and the backend is also an asynchronous process. The IPSS built-in statements of POST/WAIT EVENT allow XTCAM, XHDSK and HOBETR to remain in a wait state until needed. The IPSS built-in statements FIND IN QUEUE, WAIT IN QUEUE, and REMOVE FROM QUEUE allow these services to maintain a queue containing suspended SIDPERS transactions via any queueing discipline.

Testing of this host submodel was done in conjunction with the previously tested SIDPERS application submodel. Since the host model was always in a reactive mode, completely independent testing was not appropriate. Additionally, the use of this submodel was to determine the effect of host support for SIDPERS processes to that the SIDPERS application submodel was a necessary part of this testing.

A3.3 IPSS Model of Backend Computer

The backend computer was modeled by first characterizing the great features of that system, i.e., the hardware and data, and then modeling the dynamic software features. Since previous USACSC studies had indicated that the PDP-11/70 would be used as the backend computer, a more specific model than that built for the host could be developed. The hardware configuration for the existing USACSC AERMICS PDP-11/70 at Georgia Institute of Technology was modeled as part of the Systems Reserve Component. Since backend computer is to be used for other applications besides supporting the IDMS system, the complete configuration needed to be modified, not only those elements directly affecting IDMS.

Figure A2.2-1 illustrates the idealized hardware configuration of the USACSC AIRMICS PDP-11/70 system used in the model. Although the physical system itself was unavailable, sufficient DEC documentation existed so that detailed modeling of individual device characteristics could be achieved. Several hardware elements of this model that have significant impact on the software processing and thus software modeling, of this system are: (1) the central UNIBUS data and control path; (2) the two dedicated high speed controllers to disk and tape drives; (3) the use of cache memory by the core and; (4) communication link between host and backend. It should be noted that in this and many IPSS models, the characteristics of any one of the hardware devices can be modified without requiring the modification of the rest of the model.

Defining the characteristics of the particular data sets that reside on this system is the second modeling step. For this application, only the IDMS test data base was modeled. Each of the IDMS records was, essentially, a 1024 byte page (i.e., physical record) with varying highly logical records in each page. The IDMS test data base contained 200 pages and the IDMS file organization was direct. Other data sets for this and other applications can be easily added to these definitions.

The characterization of the backend software processing follows much the same pattern as the characterization of the host software. Both operating systems are assumed to be in reactive models, waiting for and responding to the needs of the resident applications. Thus, to determine the operating system functions to be modeled, the processing requirements of the IDMS applications were examined. A DML request in the backend follows a five step process: request arrival, IDMS scheduling, IDMS processing, physical record retrieval, and record transmission to the host. (See Figure A3.3-1). At each of these steps particular operating system functions are required.

To satisfy functional IDMS requirements for DML processing, certain PDP-11/70 operating system functions were identified. As Figure A3.3-2 shows, these were: the scheduling and transmission of data to/from the host, the scheduling of IDMS within the backend, and the necessary resource allocation, buffer allocation, I/O scheduling and I/O execution to retrieve the desired record. Underlying processes such as the synchronization of IDMS and backend functions, which may be executing in parallel, is provided for in the model by generalized task management and resource allocations procedures. Although CPU scheduling was not included as part of these functions, UNIBUS scheduling was an integral part of the synchronization process.

The relationships between the general DML processing step and their corresponding IDMS services, and between these services and those characterizing backend operating system functions is depicted in Figure A3.3-2. Both the first and the last steps in DML processing

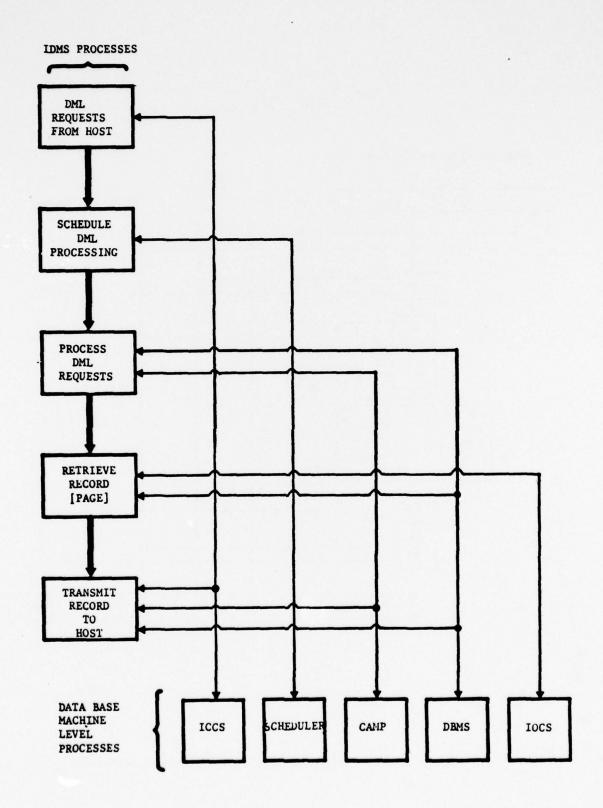


Figure A3.3-1 Relation of Data Base Machine Processes to IDMS Processes

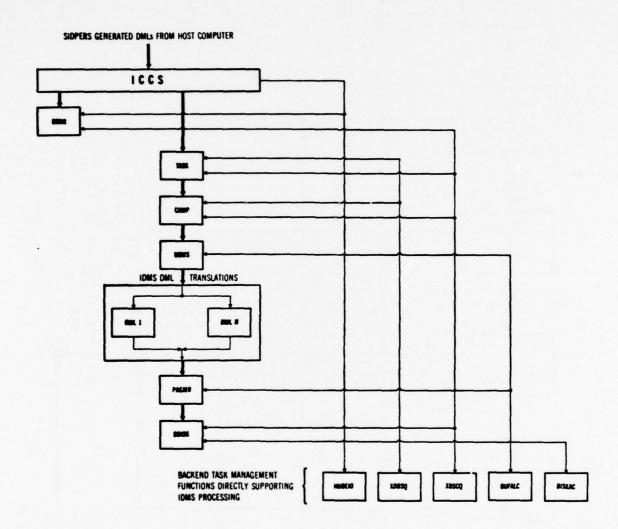


Figure A3.3-2 IPSS Realization of Back-end/IDMS Model

require communication with the host machine. This entails not only a service to characterize the backend interfacing functions such as the message printing and line protocol, but also a service to schedule the intercomputer communications. In the resulting IPSS submodel, this scheduling function was performed by the service, MUB910. It scheduled, in a FIFO manner, all requests to use the host/backend communication interface, from either the host or backend direction. It was assumed for this model that only one message could cross the interface at a time. Once the service HUBE10 scheduled and allocated the interface, the requesting service gained control of the use and disposition of the interface.

The second step in DML processing is the scheduling of IDMS itself. As mentioned before, the PDP-11/70 is not assumed to be decicated to the support of IDMS, allowing the concurrent execution of other non-IDMS applications. Thus, each invocation of IDMS, via entering DML command, must be scheduled and the appropriate IDMS software allocated. In the structure of IDMS its twin in the host application signals CAMP (Central Access Monitor Program) that a DML requires IDMS, and CAMP then schedules the invocation of IDMS itself. Only one DML can use the data base functions at a time to prevent concurrent updatings. The IPSS submodel for IDMS contained IPSS services corresponding to CAMP and the actual data base manipulation routines (DBMS). The signaling to CAMP that a DML request has arrived and needs DBMS is accomplished via the IPSS service XDBSQ. The actual allocation and deallocation occurs in CAMP.

Once the DBMS module is involved, the DML is processed in the third step via IDMS internal routing. Each of these routines has an analogue IPSS service that accomposes the high level DML into a sequence of primitive DBMS commands. IDMS individually changes these primitives to requests to its virtual data store for particular record(s) which leads to step 4.

The actual record retrieval in IDMS is a two step process. First, the record location is determined with respect to an IDMS page. The currently core resident pages are examined to see if one of them

contains the desired record. If none do, then a new page must be accessed. This page management function is accomplished via the IPSS service PAGMR (page manager). If a page must be accessed, then buffer space must be allocated (BUFACE) and the I/O request must be scheduled (XDSCO). When the I/O operations can be performed, the actual hardware devices for the PDP-11/70 backend computer are referenced via the IPSS service DISKB. Furthermore, DISKB simulates the allocation/deallocation of resources and the actual I/O operations of each and data transfer. DISKB also takes into account the dedicated data lists between main memory and auxiliary storage.

In order to characterize the backend operating system's functions as they normally operate, i.e., asynchronously, the IPSS IDMS and PDP-11/70 submodel also had to execute asynchronously. The host/backend interface (through HUBE10) is also asynchronous. To accomplish this behavior, communication with operating system functions was characterized via the posting of common IPSS Event signals (semaphores). The IPSS built-in statements of WAIT EVENT cause the XBBSQ, HUBE10, XBXCQ, XDSKQ services to remain in a wait state until needed. Conversely, once awakened by the IDMS services, these schedulers utilize the IPSS built-in statements FIND IN QUEUE, WAIT IN QUEUE and REMOVE FROM QUEUE to maintain queues of suspended IDMS DML requests.

A3.4 IDMS Models

The purpose of the model of IDMS was to characterize IDMS DML processing in order to determine:

- 1. the IDMS contribution to transaction response time,
- 2. the effect of widely varying workload definitions,
- the effect of processing DMLs concurrently (through multiple terminal loadings), and
- 4. the effect of data structure on processing efficiencies.

Two models of IDMS were prepared, one employing the current IPSS facilities and the other incorporating advanced IPSS features for data

structure characterization and data base access. These two models are similar in structure, they contain the same services and the calling sequences remains unchanged. However, the IPSS/BASIC model does not make use of the DML parameters passed from the SIDPERS transaction services. Virtual page identification is characterized through probability distribution functions. Schema definition and data access currency are not represented. It was found that these important features of data base systems are very difficult to model without the special simulation facilities provided by IPSS/DBS.

The second model of IDMS represents these features through the advanced data structure and data access facilities of IPSS/DBS. In this model, the DML parameters passed by the application services are used to traverse the data base. The data base structure is characterized through the SCHEMA facility and instances of the data base are created through special purpose IPSS/DBS built-in functions. This results in a much more accurate determination of data base processing activities in general and virtual page referencing sequences in particular. Secondary Storage functions were initially characterized by waiting an average elapsed time of 36 ms (the average time for PDP-11/70 head positioning and rotational delay). The detailed model of the PDP-11/70 was used to calculate these times.

Model Structure - Model I

The IPSS model of IDMS consists of services which interact in much the same way that the IDMS itself executes. The overall structure of the model is shown in Figure A3.4-1. The DML commands from the application services invoke the IDMS Interface service which in turn invokes the CAMP (Central Access Monitor Program) and DBMS services. The DBMS service invokes one of the 7 DML services which represent data access and retrieval. The interact with the Page Manager service and a single Schema Definition service in order to generate secondary storage references for the underlying file management system. The following is a description of the function of each of these services.

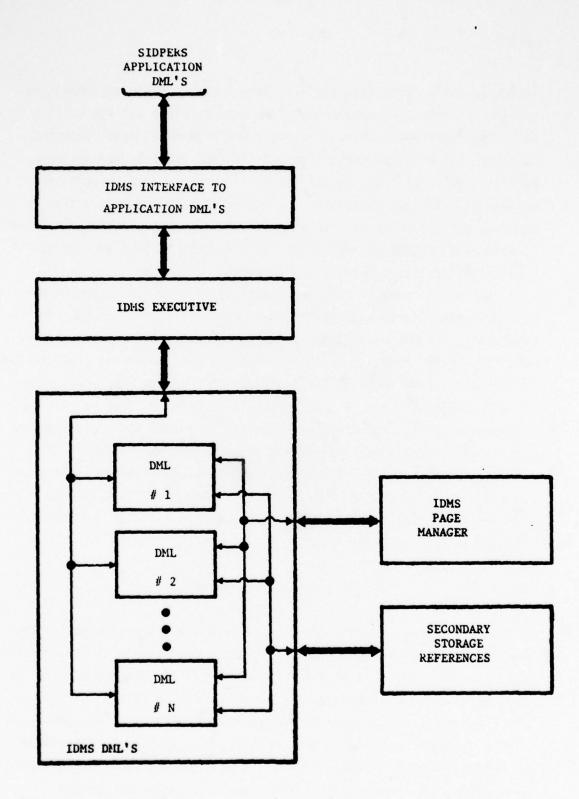


Figure A3.4-1 Schematic of the Modeled IDMS

IDMS Interface

This service represents IDMS processing activities that occur for every DML invocation. It is invoked by one of the application programs and represents the following IDMS functions:

- 1. validity checking and segmentation.
- 2, error status checking,
- 3. performing any 'before' procedures,
- 4, testing error status from the 'before' procedures,
- 5. testing the chanel indicator
- testing error status on return from performing the data base management system function,
- 7. executing any "on error during" procedures,
- 8. executing any 'after error' procedures, and
- 9. performing segmentation.

However, no information was available on the frequency or nature of these procedures within the test VIABLE/SIDPERS system. Thus, these functions were merely identified and cocumented within this service and can thus be easily expanded once more data becomes available. This service invokes the CAMP Service and waits for its completion.

Central Access Monitor Program (CAMP) Service

This service represents the IDMS program of the same name. Its function is simply to thread requests as the DBMS service can accommodate them thus holding the invoking services in the wait state. Because of the single thread of processing represented by a single user terminal, the queueing of requests at this level does not occur. However, queueing is an important performance factor for models with more than one terminal and/or concurrently executing application programs.

Data Base Management System (DBMS) Service

This service represents those functions at the very heart of IDMS, namely the materialization of an application record. This invokes the schema and subschema definitions, data base currency, the content of

buffers, and (usually) requests to the operating system for one or more data base pages. Because of the magnitude and complexity of these processes, they are modeled through several IPSS services. The DBMS service establishes the environment (e.g., page size, buffer size), collects cache paping statistics, invokes one of the IPSS DML services and waits for its completion.

DML Services

There are 7 IDMS DMLs represented in the IPSS model, each through its own service definition. These DMLs are the IDMS functions identified in Table A2.3-4. Each of these services performs the logic of the DML by:

- determining the virtual page number from the record tape and set parameters as well as currency,
- determining if the requested page is already present in the IDMS page cache, and
- initiating a service to retrieve the record if not present.

Model Structure - Model II

The IPSS/DBS model of IDMS is similar in structure to the IPSS/BASIC model reported in the previous section. As shown in Figure A3.4-1, the application DMLs invoke the CAMP service which in turn invokes the DBMS service. In this model, however, the DBMS service has the additional functions of creating an instance of the schema. This is accomplished by the CREATE SCHEMA built-in facility which copies the Data Base Structure Component schema definition into an internal work area. This service them creates a set of arrays (called routes) which represent schema access paths. These routes are the basic mechanism for creating instances of sets, maintaining currency, and simulating the traversal of the data base.

A route consists of a series of set identifiers which includes the identification of the owner and member record types. A single

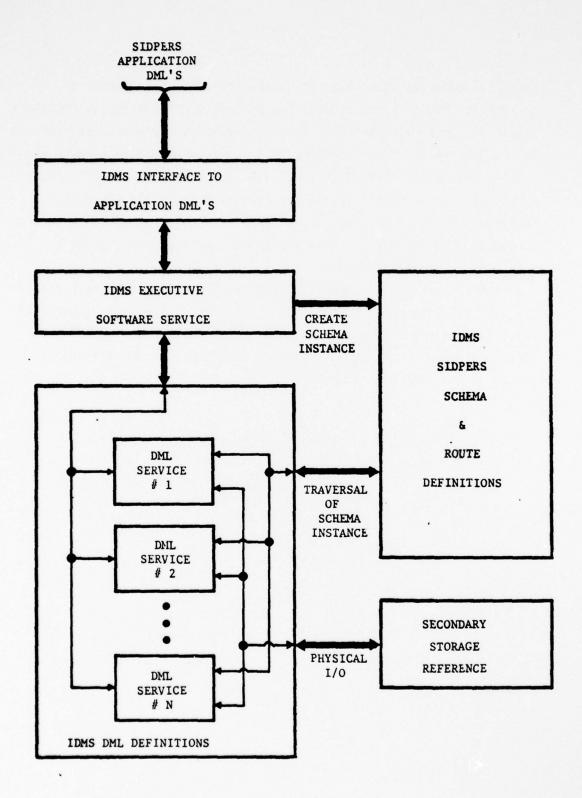


Figure A3.4-2 Schematic of IPSS/DBS Model of IDMS/SIDPERS Data Base

route incorporates those sets which are trayersed by a series of application DMLs. For example, the DML series of locating an individual soldier's record via the CALC set followed by locating his data records in the UIST record type in Figure A2.3-2, would result in a route consisting of the CALC act, IDENT-UIST set, and the IDENT and UIST record types. Instances of these sets are created dynamically based on parameters supplied for the schema facility in the Data Base Structure component. In addition, statistics are automatically collected on data base traversal activities.

The DML services in this IPSS/DBS model simulate the traversal of the data base through such built-in facilities as FIND NEXT (member in set), and determine a record's virtual page address via the GET DATA BASE PAGE command. These virtual page references become the loading on the file management system.

APPENDIX B

SUMMARY OF MODEL STATISTICS

This appendix presents an overview and explanation of statistics generated by the execution of the IPSS models synthesized for the experimentation phase of the research. All of the defined models are not reproduced here, however, the final section of this Appendix contains the IPSS source listing for the SIDPERS/IDMS workload model.

B1. SIDPERS/IDMS WORKLOAD MODEL STATISTICS

The SIDPERS/IDMS workload model consists of a collection of IPSS service facilities whose invocation sequence is hierarchical. Figure B1-1 depicts the relationships among the services and indicates their generic function. As shown in the figure, transaction arrival at a work station is represented in the START service. Since a single terminal is modeled, the service queues all arriving transactions until the proceeding one has completed. Transaction parameters are assigned (TRANS) and the data fields are inputted and validated (VALID). Once the application programs is invoked (ADDSL, ARRIV, DEPTR, DTYC, GRCH, ING) which issues a DML by invoking the IDMS interface routine (modeled by INTF) and passing appropriate parameters. IDMS processing is sequenced by CAMP which then invokes DBMS. The function of the DBMS service is to establish the IDMS processing environment and to invoke a DML service (CONECT, DISCON, ERASE, AND, MODIFY, OBTAIN, STORE). These services identify an IDMS page and request its presence in core by invoking the PAGMR service. PAGMR manages the IDMS page buffer using a modified least-recently-used page replacement algorithm. In this model, data transfer from secondary storage is simulated by a clock advance of 36 ms. Table B1-1 relates the IPSS services identified in Figure B1-1 to the corresponding SIDPERS/IDMS activity.

Statistics relative to this modeled process are contained in the following pages. An elapsed time of two hours was simulated; all times

Table B1-1. <u>Index to Modeled SIDPERS/IDMS Processes</u>

.eve1	SIDPERS/IDMS Process	IPSS Endogenous Exogenous Service Name	Source Code Page
1	Arrival of transaction at a work station	START	3
2	Transaction input and verification		
	a. Transaction input by terminal operator	TRANS	4
	b. Transaction verification by system software	VALID	10
3	Application software processing		
	 Add a soldier to the data base Process soldier's arrival at the installation 	ADDSL ARRIV	31 23
	c. Process soldier's departure from installation	DEPTR	27
	d. Process change in soldier's duty status	DTYC	18
	e. Process soldier's grade change f. Process inquiry	CRCH INQ	13 21
4	Software processing performed by the IDMS interface routine	INTF	35
5	IDMS Central Access Monitor Processor (CAMP)	CAMP	37
6	IDMS DBMS processing	DBMS	39
7	IDMS DML processing		
	a. CONNECT DML b. DISCONNECT DML c. ERASE DML d. FIND DML e. MODIFY DML f. OBTAIN DML g. STORE DML	CONECT DISCON ERASE FIND MODIFY OBTAIN STORE	50 51 52 45 46 42 47
8	Page management, retrieval of pages from auxiliary storage	PAGMR	54

SERVICE FUNCTION	transaction arrival	transaction input and verification	application processing	IDMS interface routine	DML sequencer	control	DML's	page manager
			INQ				OBTAIN STORE	
CE HIERARCHY		VALID	DTYC GRCH IN				FIND MODIFY	
IPSS MODEL - SERVICE	START	TRANS	ARRIV DEPTR D	INTF	CAMP	DBMS	DISCON ERASE	PAGMR
11			ADDSL ARE				CONNECT	

Figure B1-1. SIDPERS/IDMS Workload Model Service Invocation Sequence

are reported in milliseconds. The first page of model output shows that the exogenous service START was busy 62.4% of the time and was invoked nineteen times. The mean number of transactions at the work station at any one time was 1.1 (the maximum number was 3) indicating that a single terminal is sufficient to handle the processing load anticipated for a small data base (203 soldiers).

The average elapsed time from transaction arrival at a work station to complete is 5.64 min (338,430 ms). An examination of the transit time statistics for the endogenous service facilities reveals that the majority of this time was spent in transaction input and validation (TRANS and VALID services). This is due to the relatively long data input time required and the terminal operator error rates reflected through these services.

A significant drop in the percent busy and the busy period mean length statistics is evidenced for the INTF service. This indicates that more terminals could easily be supported. An appropriate number could be determined by experimentation with the model parameters.

An examination of the DML statistics reveals a relatively large number of executions of these services and that the average processing time is low. For example, there were 83 invocations of the FIND DML service and the average processing time was 37.3 ms. Most of this time was consumed by the PAGMR service which all the DML's invoke. The PAGMR is required to bring pages into core, and as reflected in the cache page fault statistics, the miss ratio was high (94.7%). This indicates that the data base is not well organized with respect to the DML processing. As inspection of the DML's in the application program services, however, reveals that a majority of the processing is random via CALC sets, and hence there is very little chance of more optimally organizing the data base given this processing load.

The last page of statistics for this model displays the queueing of transactions for the DBMS and TRANS services. Since the DBMS service is single-threaded (the function of the CAMP service), no queueing to it occurred. However, if CAMP was multi-threaded and the DBMS was not,

no change would be required in the model to collect DBMS queueing statistics. The queueing statistics for the TRANS service shows that the queue was never very long (maximum of 2, mean length of .6) again indicating a larger data base or heavier transaction rate could easily be accommodated.

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PAGE

INFORMATION PROCESSING SYSTEM SIMULATOR

FACILITY TYPE = ENS			
TYPE = ENS			
S DESCRIPTION PACILITY FACILITY FACI			
TICS O.2 O.4 O.1 IR ST4.2 O.4 O.1 IR ST5.4 ST6.6 AN LENCTH TICS O.4 O.1 IR ST7.6 ST7.6 ST8.7 TICS O.4 O.1 IR ST7.6 ST8.7 ST8.6 ST8.7 ST8.6 ST8.7 ST8.6 ST8.7 ST8.6 ST8.7 ST8.6 O.0 O.0 O.0 O.0 O.0 O.0 O.0 O	FACILITY 4	FACILITY S NAME INDEX	FACILITY C
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TICS 19.8 99.6 99.9 19.61 19.8 575 61 19.8 575 61 19.9 575 61 19.9 575 61 19.9 575 61 19.9 575 61 19.9 575 117925.8 117025.8 117	36.6	5.3	54.5
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STEED DEV DE NEAN LENGTH	239647.5	114307.8	77980.0	395990.6	13525.3	162139.3
CONCURATINCY STATISTICS CURRENT LEVEL	0	0	0	0	0	•
HAXIHUA LEVEL	9	1	1			
STAG DEV OF MEAN LEVEL	6.0	0.5	0.5	6.0	0.0	5.0
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	DBMS 1	TRANS 1				
BUSY PERIOD SYATISTICS	c	23.0				
NUMEER	574	16				
HEAR LENGTH	0.0	95422.4				
STND DEV OF MEAN LENGTH	0.0	144971.9				
TDLE PERIOD STATISTICS PER CENT IDLE	100.00	76.1				
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B2. HOST/SIDPERS SUBMODEL STATISTICS

Previous USACSC studies did not identify a particular machine to be used for the host environment, so a generalized host was modeled. To do this, the software functions of the host were emphasized more than the hardware functions. Thus, only a minimum hardware configuration was described and more attention was paid towards determining and modeling the host operating system requirements of the SIDPERS application.

The host operating system was assumed to execute in a reactive manner, responding to the needs of the SIDPERS application. The particular steps in the modeled SIDPERS application were discussed in the main body of this report and in Appendix A. The needs of this application were, in general, task management and resource allocation. A crucial host/SIDPERS interface was the synchronization of the execution of SIDPERS services and operating system services. The asynchronous nature of their execution compounded the problem. To model this situation, IPSS services representing task schedulers communicated with existing SIDPERS IPSS services via the posting of and waiting for global events. These events corresponded to a request for scheduling and the completion of that scheduling process.

Four schedulers were identified for the modeling of the host/SIDPERS interface. These controlled access to the host's disk subsystem, the use of the teleprocessing(TP) line to the user, the use of the teleprocessing line to the back-end and the resumption of execution after a back-end request was completed. Each of these were characterized by a separate IPSS service and their use was controlled through a queue, exemplified by an IPSS chain facility. When a SIDPERS transaction required scheduling, it was placed in one of these chains. The four chains were HDSKQ (disk subsystem), TCAMQ (user TP communication), HOBEQ (back-end TP communication) and RSPONS (back-end request completion).

The use of these chains was measured in IPSS via the queueing statistics shown. Because of the particular SIDPERS workload, no SIDPERS applications overlapped, and, therefore, there was no queue in transit time associated with any scheduler. This condition also limited the possibility of having more than one SIDPERS transaction enqueued simultaneously. If such a situation occurred in the future, the concurrency statistics would demonstrate the degree of delay.

The queue entry statistics indicate the number of times that the current SIDPERS workload requested each of the scheduling tasks. The SIDPERS load for the run depicted is shown in the utilization statistics for the IPSS services that characterize SIDPERS processing. The total number of SIDPERS transaction for this run was 9, i.e., the number of requests to VALID. The number of requests to HDSKQ corresponds to the number of validity and/or compatability errors that occurred during the current run, since each error required a sisk access. Thus, on the average, each incoming transaction had one error.

Statistics measuring the communication with the user and the back-end indicate the complexity of the SIDPERS application transactions. The number of user interactions (TCAMQ) reflects the amount of input required for the current set of SIDPERS transactions. Thus, for 9 transactions, a mean of 5.1 interactions were required per transaction. This included input of data and output of error messages and records. The number of entries in the RSPONS queue indicates that the current input stream of SIDPERS transactions generated 145 DML request to the back-end. This implies a mean of 16.1 DML requests per SIDPERS transaction types and does not take into account the DML requirements of particular transaction types. And finally, the number of entries in the back-end interface chain, HOBEQ, is twice the number of the actual DML requests since the TP interface must be acquired for both sending and receiving DML requests.

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B3. IDMS/BACK-END SUBMODEL STATISTICS

The emphasis of this model was on interfacing the IDMS functions with the PDP-11/70 backend operations. Because of the wealth of documentation available on the hardware aspects of the USACSC PDP-11/70 (and the lack of software documentation), the impact of IDMS on hardware utilization was examined. The complete configuration of the USACSC PDP-11/70 installation at the Georgia Institute of Technology was modeled and is shown in Figure B3-1.

The PDP-11/70 operating system employs a dedicated data path between central memory and high speed disks (RWPO4AA in Figure B3-1). The IDMS data base resides on these disks, allowing for rapid retrieval of IDMS pages. DML requests entering the backend from the host were eventually mapped to specific data accesses to the RWPO4AA. Thus, to measure the direct effect of the IDMS loading on the PDP-11/70 backend, the IPSS statistics on the behavior of this disk subsystem (labeled RPO4AA in the corresponding IPSS model) are examined and are found immediately following this discussion.

Utilization statistics on the RPO4AA provide the best insight into its performance relative to the IDMS loading. Busy/idle statistics indicate the percentage of time that this facility was utilized. As can be seen, the RPO4AA was used only 23.1% of the time, remaining idle for 76.9%. Such an imbalance can be attributed to the characteristics of the device exceeding the needs of the application or the actual placement of data on the disks.

Concurrency statistics measure the degree of multiprocessing that occurred on the RPO4AA. In this case, in order to maintain the integrity of the IDMS data base, only one access at a time was permitted. This is evidenced by the value of the maximum level of concurrency being 1. The value of the current levels equalling 1 implies that at the end of the current simulation run, the RPO4AA device was in use.

The last two sets of utilization statistics characterize the behavior of individual accesses to the RPO4AA. Each access involves a seek and a data transfer operation. The RPO4AA was acquired (seized) throughout both operations. The total number of references(19) and the number

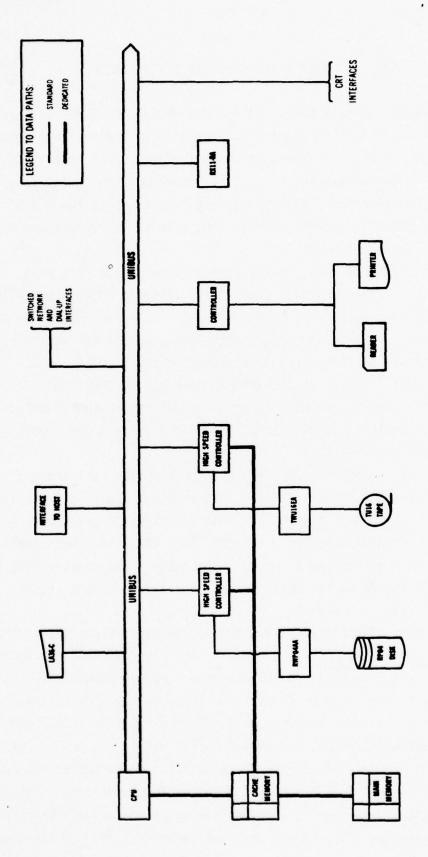


Figure B3-1 USACSC PDP 11/70 Configuration

of zero time references (0) indicate the relative closeness of successive page references to the IDMS data base. A low percentage, as is the case here, indicate that the data may be well spread out on the disk or that the individual accesses are randomly distributed. On the other hand, a high percentage of zero time seizures may indicate the compactness of the data or a deliberate ordering of accesses (e.g., sequential of within an IDMS set). The mean total time for an individual accesses to the RPO4AA, 23.2 milliseconds, includes the seek and data transfer operations. An additional mean is calculated removing the bias of the zero time seizures so that an average time per actual access can be determined. For each mean, the variability of that mean, i.e., the standard deviation, is also calculated to give some insight into the relative accuracy of the above means.

The impact of the IDMS functions can also be shown in other IPSS statistics. The IDMS DML requests are mapped into I/O requests for which IPSS has built-in primitives such as Seek and Data Transfer. Statistics are automatically calculated by IPSS on invocations of each of these I/O operations by each I/O operation (Part B1), by access mechanism (Part B2) and by data set (Part B3). The statistics by the access mechanism, RPO4AA, aid in analyzing the impact of IDMS on the backend. The sum of the mean routine times for the two I/O operations that referenced RPO4AA equals 23.2 milliseconds, the same utilization time for RPO4AA in the previous discussion. This allows us to break down the aggregate access time to RPO4AA into its component activities. It is evident that the Seek operation comprises the greatest amount of time in each DML request processing. Thus, in order to reduce this time factor affecting the Seek operation, such as placement of data files, compactness or device type, should be examined.

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B4. IPSS FUNCTIONAL MODEL OF THE SIDPERS/IDMS

The purpose of this section is to illustrate the use of the IPSS to describe a complex information processing system at the functional level. This model does not contain characterizations for the computer hardware, executive software, data base management system software or data base. At this level they were treated as black boxes, however appropriate interfaces were defined so that respective characterizations could be included in later models.

The focus of this model was to model the processes (represented as IPSS endogenous services) encountered in updating the SIDPERS data base. Figure B1-1 identified the hierachical relationship assumed for the model; Table B1-1 related the IPSS service identified in Figure B1-1 to a corresponding SIDPERS/IDMS data base update procedure. The IDMS data base management system backend was also represented in the model.

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12 GRADE CHÁNGE GOGCO 3-5 12 DUTY STATUS GOGGO370 3-7 12 INQUIRY GOGGO380 3-9 12 ARRIVAL GOGGO390 3-9 12 ARRIVAL GOGGO390 3-9 12 ADD SOLDIER GOGGC416 2-1 13 HAST INTERFACE GOGGG420 4-2 14 ACCESS MONITOR GOGGG430 4-2 22-23 FIND DML GOGGG430 4-5 22-23 FIND DML GOGGG470 4-7 22-23 FIND DML GOGGGA70 4-7 22-24 FIND DML GOGGGA70 4-7 22-25 FIND DML GOGGGA70 4-7 22-25 FIND DML GOGGGA70 4-7 22-26 FIND DML GOGGGA70 4-7 22-27 FIND DML GOGGGA70 4-7 22-27 FIND DML GOGGGA70 4-7 22-28 FIND DML GOGGGA70 4-7 22-29 FIND DML GOGGGA70 4-7 22-20 FIND DML GOGGGA70 4-7 23-20 FIND DML GOGGGA70 4-7 24-20 FIND DML GOGGGA70 4-7 25-20 FIND DML GOGGGA70 4-7 25-20 FIND DML GOGGGA70 4-7 25-20 FIND DML GOGGGA70	12 GRADE CHANGE GOGCO350 25 12 DUTY STATUS GOGCO370 27 12 ARRIVAL GOGCO390 39 12 ARRIVAL GOGCO390 39 12 ADD SOLDIER GOGCC410 29 13 HOST INTERFACE GOGCOC420 42 14 ACCESS MONITOR GOGCO420 42 22-23 GETAIN DAL GOGCO450 44 22-23 FIND DAL GOGCO640 47	12 GRADE CHANGE GOGGO360 25 12 DUTY STATUS GOGGO370 77 12 INGUIRY OCCGO380 39 12 ARRIVAL GOGGO390 20 12 ARBYURE GGGOGGO 20 12 ADD SOLDIER OCCGO420 42 14 ACCESS MONITOR OCCGO420 44 22-23 GETAIN DML GGGOGGGO 47 22-23 FIND DML GGGOGGGO 47 22-23 FIND DML GGGOGGGO 47 22-23 FIND DML GGGOGGGO 47 22-23 STORE DML GGGGGGGO 47 32-23 STORE DML GGGGGGGO 48 6.5 6.5 6.5 6.5 6.5 6.5 6.5 6.		LID (ENS)	(10)		TERMINAL INTERACT	10N000000350	36		16
12	12 DUTY STATUS GOOGO370 37 12 ARRIVAL 00000390 39 12 ARRIVAL 000000390 39 12 ADD SDLDIER 00000400 42 13 HOST INTERFACE 00000440 44 22-23 GETAIN DHL 00000440 45 22-23 HODIFY DHL 00000470 47 22-23 STORE DHL 00000470 47	12 DUTY STATUS CO000370 27 12 INQUIRY 00050380 39 12 ARRIVAL 00000390 39 12 ADD SOLDIER 00000410 27 13 HOST INTERFACE 00000420 42 15-21 DEMS 00000440 44 22-23 FIND DML 00000470 47 22-23 FIND DML 00000470 47 22-23 FIND DML 00000470 47 22-23 STORE DML 00000470 47		CH (ENS)	(13)	12	GRADE CHANGE	09000000	35	The contract of the contract o	35
12 INGUIRY 00050380 39 12 ARRIVAL 00000390 39 12 DEPARTURE 00000400 40 12 ADD SOLDIER 0000C410 71 13 HAST INTERFACE 00000420 42 14 ACCESS MONITOR 00000440 44 22-23 GETAIN DML 00000440 47 22-23 FIND UML 00000440 47 22-23 STORE DML 00000440 47	12 INGUIRY 00050380 39 12 ARRIVAL 00000390 39 12 DEPARTURE 00000400 40 12 ADD SOLDIER 00000420 42 13 HOST INTERFACE 00000440 44 15-21 DEMS 00000440 44 22-23 GETAIN DML 00000470 45 22-23 FIND DML 00000480 40 22-23 STORE DML 00000480 40 22-23 STORE DML 00000480 40	12 INQUIRY 00050380 39 12 ARRIVAL 000000390 39 12 DEPARTURE 00000400 4) 12 ADD SDLDIER 00000410 4; 14 ACCESS MONITOR 00000440 44 22-73 GETAIN DHL 00000440 46 22-23 FIND DHL 00000440 47 22-23 STORE DHL 00000440 47	10	YC (ENS)	(18)	112	DUTY STATUS	00000370	45		30
12 AKRIVAL 00000390 39 12 DEPARTURE 00000400 40 12 ADD SOLDIER 00000410 71 13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 44 22-23 GETAIN DML 00000450 45 22-23 FIND DML 00000470 47 22-23 FIND DML 00000460 47 22-23 STORE DML 00000440 47	12 AKRIVAL 00000390 39 12 DEPARTURE 00000400 40 12 ADD SOLDIER 00000420 42 13 HRST INTERFACE 00000420 43 15-21 DEHS 00000440 44 22-23 FIND DHL 00000460 45 22-23 HODIFY DHL 00000470 47 22-23 STORE DHL 00000470 47	12 AKRIVAL 00000390 39 12 DEPARTURE 00000400 40 12 ADD SOLDIER 00000420 42 13 HRST INTERFACE 00000420 43 15-21 DEMS 00000440 44 22-23 GETAIN DML 00000440 45 22-23 HIND DML 00000470 47 22-23 FIND DML 00000470 47 22-23 STORE DML 00000480 47		IO LENS!	(21)	12	INGUIRY	000000	36		
12 DEPARTURE 00000400 27 12 ADD SOLDIER 000007410 21 13 HUST INTERFACE 00000420 42 14 ACCESS MONITOR 00000440 44 22-23 GETAIN DML 00000450 45 22-23 FIND UML 00000470 47 22-23 STORE DML 00000440 47	12 DEPARTURE 00000400 27 12 ADD SOLDIER 00000410 21 13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 44 15-21 DEMS 00000440 44 22-23 GETAIN DML 00000450 45 22-23 FIND UML 00000470 47 22-23 FIND UML 00000440 47 22-23 STORE DML 00000440 47	12 DEPARTURE 00000400 27 12 ADD SOLDIER 00000410 21 13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 44 15-21 DEMS 00000040 44 22-23 GETAIN DML 00000450 45 22-23 FIND DML 00000470 47 22-23 FIND DML 00000470 47 22-23 STORE DML 00000480 48		RIV (ENS)	(23)	12	ARRIVAL	00000390	39		aç
12 ADD SDLDIER 0000CG416 7! 13 HRST INTERFACE 00000G420 42 14 ACCESS MONITOR 00000G430 44 22-23 GETAIN DHL 00000G450 45 22-23 HRDIFY DHL 00030G470 47 22-23 STORE DHL 0000G480 47 00000G460 44 27-23 STORE DHL 0000G460 47	12 ADD SOLDIER 0000C0415 7! 13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000440 44 15-21 DEMS 000000440 44 22-23 GETAIN DML 000000450 45 22-23 HRDIFY DML 00000470 47 22-23 STORE DML 00000480 48	12 ADD SOLDIER 0000C0415 7! 13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000440 44 15-21 DEMS 00000440 44 22-23 GETAIN DML 00000450 45 22-23 FIND DML 00000470 47 22-23 FIND DML 00000470 47 22-23 STORE DML 00000470 47		PTR TENS)	1211	- 21	DEPARTURE	00000000	67		46
13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 43 15-21 DEMS 00000040 44 22-23 CETAIN DML 000000450 45 22-23 FIND DML 00000040 47 22-23 FIND FWL 00000440 47 22-23 STORE DML 00000440 47 000000480 47	13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 43 15-21 DEMS 00000440 44 22-23 GETAIN DML 00000450 46 22-23 FIND DML 00000470 47 22-23 HRDIFY DML 00000470 47 22-23 STORE DML 00000480 48	13 HRST INTERFACE 00000420 42 14 ACCESS MONITOR 00000430 43 15-21 DEMS 00000040 44 22-23 GOTAIN DML 000000450 45 22-23 FIND DML 00000470 47 22-23 STORE DML 00000460 48		IDSL LENS!	(3II)	172	ADD SOLDIER	00000410	ia		40
14 ACCESS MONITOR 06000430 43 15-21 DEMS 00000440 44 22-23 GETAIN DML 060330460 46 22-23 MIDJFY DML 060330460 47 22-23 STORE DML 00000440 28	14 ACCESS MONITOR 06000430 43 15-21 DEMS 00000440 44 22-23 DETAIN DML 06000450 45 22-23 FIND DML 060330460 47 22-23 STORE DML 00000440 28	14 ACCESS MONITOR 06000430 43 15-21 DEMS 00000440 44 22-23 GETAIN DML 06000450 45 22-23 FIND DML 060330460 47 22-23 FORE DML 00060480 47		ITE (ENS)	(35)	13	HOST INTERFACE	00000420	7.5		17
15-21 DEMS 00000440 44 22-23 GETAIN DML 00000460 45 22-23 FIND DML 060,00470 47 22-23 MODIFY DML 000,00470 47 22-23 STORE DML 000,00470 47	15-21	15-21		MP (ENS)	(37)	14	ACCESS MONITOR	00000430	67		5,4
22-23 GETAIN DML 060606450 45 22-23 FIND DML 060330460 47 22-23 MPDIFY DML 06030480 47 22-23 STORE DML 06060480 48	22-23 FIND DML 00000450 45 22-23 FIND DML 000330460 44 22-23 MCDIFY DML 00000470 47 22-23 STORE DML 00000480 48	22-23 GETAIN DML 060600450 46 22-23 FIND UML 060330460 47 22-23 HODIFY DML 500306470 47 22-23 STORE DML 060606480 48	14 03	MS (ENS)	(36)	15-51	DEMS	000000	44		4.3
22–23 FIND DML 0603.304.60 47 22–23 MCDIFY DML 0000.0470 47 22–23 STORE DML 0000.0480 48	22–23 FIND DML 060:104:60 47 22–23 HnDIFY DML 000:0470 47 22–23 STORE DML 000:04:00 28	22-23 FIND DML 060:004:60 47 22-23 MODIFY DML 500:004:70 47 22-23 STORE DML 000:004:00 48		TAN (ENS)	(45)	22-23	DETAIN DML	05700000	74		777
22-23 MODIFY DML 50000470 47 22-23 STORE DML 000004R0 48 0506070	22–23 MODIFY DML 500006470 47 22–23 STORE DML 00000440 24 0	22–23 MODIFY DML 50000470 47 22–23 STORE DML 00000480 28 0506070		ND TENS)	1551	22-23	FIND DML	0000000	44		54
22-23 STORE DML 000004R0 28	22-23 STORE DML 0000004A0 2A	22-23 STORE DWL 0000004A0 2A		DIFY (ENS)	(49)	22-23	MODIFY DML	00000000	47		94
0506070	0506076	0506070		ORE LENS!	(41)	27-23	STORE DML	000000480	87		7.4
			10	2030	40		99				

	STAI	NOARD IN	STANDARD INPUT STREAM LISTING			The same of the sa	
		DATE	TL/01/60 ···				
INPUT C	ARD 0.	IMAGE	.0706070.		SEO-NO	ALT INPUT MEMBER LEV SEG-NC PEF-NG	a 3
19 CONECT (ENS)	(99)	22-23	CONNECT DML	05700000	67	44	
20 DISCON (ENS)	(53)	22-23	DISCONNECT OML	00000000	30	44	
21 FRASE (ENS)	(52)	22-23	ERASE DML	000000510	51	0.5	
22 PAGMR (ENS)	(55)		PAGE MANAGER (CACHE)00000520	1000000520	2,	53	
23 GENRP (PRUC)	(57)		PAGE # GENERATOR	00000030	63	25	
24 BufST (EXS)	(60)		PRINT PAGE STATTEV2)00000540	1000000540	75	53	
25 SNAPST (EXS)	(61)		PRINT IPSS STATTEV3 NOCHOOSSO	100000550	نَّهُ	95	
26 EXO EVENT STREAM	(41)	-	COMPONENT DEFINITION	09500000	95	35	
27 EVI (EXO EVENT)	(62)	3	INVOKE PRUCESSING	07300000	L 3	35	
28 EV2 (EXO EVENT)	(62)	54	DISPLAY PAGE STATS	000000000	ų y	1.5	
29 EV3 (EXD EVENT)	(62)	25	DISPLAY IPSS STATS	04400000	50	2.	
30 HUDEL	(65)		COMPONENT DE FINITION	00000000	07	45	
				0000000	19	07	
				00000000	6.2	.79	
RT TOMS TO	\$ \$ET		IDMS IDENTIFIER	06900000	6.3	29	
				00000000	79	3.0	
ADDI ADD-IDENT 4	410 1. CALC	IC CALC	רכ	05900000	33	3	
ADUN ADD-UNIT 4	420 2. GRUI		GRADE-UIST	03900000	66	\$9	
ALGC AALOC 4	455 3. IPUI		IDENT-UIST	000000000	F.A	97	
GRADE GRADE	350 4. THUM		INST-UNIT-MASTER	0000000	6.4	67	
IDEN TOENT	320 5. PMID		IDENT-MOSC (PMOS)	04700000	69	3	
INST INST	300 6. PPUI		Pesrio-U1ST	00100000	J.	40	
HOSC MOSC 4	460 7. SMID		IDENT-HOSC (SMOS)	01100000		01	
POSN POSNO	440 8. UMAN		UNIT-MASTER-AWOL	02700000	72	71	

. 1			
DATE 09/10/77			
INPUT CARD IMAGE10203?4^506070		TRPUT ALT ILPUT	MEMBER. REF-NC
C UIST DIST 330 9. UMGR UNIT-MASTER-GRADE	06/00/30	73	72
C UNMA UNIT-MASTER 310 10. UMPOS UNIT-MASTER-POSMO	07100000	74	73
C JI. UMUI UNIT-MASTER-UIST	051000000	54	14
RETURN	00000000	75	34
END: PROCEDURE;	27700000		7.6
	00000000	78	
	0600000	61	
EXO SERVICE: ID=START, NAME=STARTS,	00000000	RO	
SAVE AREA SIZE = 10;	00000010	8)	7.6
	00000620	26	54
C THIS SERVICE REPRESENTS THE APRIVAL OF TRANSACTIONS AT A SINGLE	00000000	. A3	36.
C WORK STATION. IT QUEUES THE TRANSACTION UNTIL THE PREVIOUS	00000000		13
C TRANSACTION IS COMPLETED.	0000000	8	14.
END: DECLARATIONS:	00000660	Fh	63
END: THITIALIZATION:	07300000	44	. 9
SEIZE: FACILITY=STAKT;	00000000	aa	685
QUEUE: FACILITY=TRANS;	04400000	96	1.6
WAIT FACILITY: FACILITY=TRAMS, STATUS=0, CONDITION=EQ;	00500000	Ųb	13
DEPART CUEUE: FACILITY=TRANS;	01400000	16	8.6
SET STATUS: FACILITY=TRANS, STATUS=!;	0250000	65	êŞ
INVOKE: SERVICE=TRANSS:	0000000	io	05
WAIT RETURN: SERVICE=TRANS:	07500000	70	13
SET STATUS: FACILITY=TRANS, STATUS=D:	05500000	26	25
RELFASE: FACILITY=START;	09600000	90	- 63

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TO SECURE

STATISTICS OF

PARTITION

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INFORMATION PROCESSING SYSTEM SIMULATOR	70A			FAGE.	ď
STANDARD INPUT STREAM LISTING					
DATE 09/10/17					
INPUT CARD IMAGE10203040506070.	80	INPUT SEG-NO	LEV SEG-NO	MERET K KIT-NO	
END: EXO SERVICE;	06500000	40		75	
	000000	96			
	06600000	66			
ENDO SERVICE: ID=TKANS, NAME=TRANSS,	000010000	001		3.6	
SAVE AKEA SIZF = In;	01010000	101		96	
3	00001050	102		14	
C THIS SERVICE REPRESENTS THE INTERACTION OF TERMINAL USER ENTERING 00001030	00001030	103		9.6	
C TRANSACTIONS INTO THE SIDPERS SYSTEM. THIS TERMINAL CUNTROL	00001000	101		66	
C PROCESSING HODGLE INCLUDES TERMINAL INTERACTION WITH THE DATA BASEODODIOSO	05010000	105		100	
C "VALID" RECORDS AND ALSO VALIDITY CHECKS. THE CHARACTERIZATION	09010000	104		101	
C DF TRANSACTION PROCESSING REPRESENTED IN THIS SERVICE WAS	06010000	101		162	
C DERIVED FROM "DESCRIPTION OF THE VIABLE TRANSACTIONS FOR THE	08010000	108		50:	
C DIMUI MODEL" BY HELMUT SCHAAFF, MARCH 1977.	06010000	100		104	
3	00001100	011		10:	
C THE TYPES OF TRANSACTIONS REPRESENTED IN THE MODEL AND THEIR	00001110	111		106	
C PERCENT UCCURRENCE ARE-	00001150	211		Tu!	
C I. GRADE CHANGE 6.68% 4. ARRIVAL 20.314	06110000	113		:08	
C MUS CHANGE IND DATA) 5. DEPARTURE 24.81%	00001140	114		109	
C 2. DUTY STATUS CHANGE 42.55\$ 6. ADD 0.77%.	60001150	115		110	
C 3. INCUIRY 4.8PX	090110000	941	The second second second second second second		
3	07110000	111	The state of the s	.:.2	
DATA SERVCE/"THAN"/	00001:80	011		113	
3	60001190	011		114	
INTEGER RETC, SEED, SUB, TLOCIIA)	00710000	120		:115	

DATE 09/10/17			
IMPUT CARD IMAGE2020304040506070	1 09	THEUT ALT IPPUT	MEMPPK NEF-NC
REAL RN, SPIALIN, SUNFM	00001210	161	116
REAL POINTS (7,2)	00001220	152	
COMMON 74K/ TLOC	00001230	123	116
DISTRIBUTION TABLE FOR THE TRANSACTION TYPES	00001240	124	411
DATA POINTS/1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 6.0,	06210000	125	150
1 0.0, 0.0668, 0.4923, 0.5411, 0.7442, 0.5923, 1.0/	00001260	961	121
	00001270	121	125
TRANSACTION PARAMETERS	00001280	128	12.5
I - PERCENT DIRECT MODE	06210000	621	124
Z - NUMBER OF DATA ITEMS PER YPANSACTION ENTRY	0001330	180	125
3 - PERCENT DATA ITEM VALIDITY EPROR	00001310	iii	lie
4 - PERCENT DATA ITEM COMPATABILITY ERROR	00001320	iĝ.	127
5 - PROCESS TIME FOR DIRECT MODE DATA ITEMS (SECONDS)	00001330	111	1,26
6 - PROCESS TIME FOR TUTORIAL MODE DATA ITEMS (SECUNDS)	00001340	781	129
7 - PERCENT OF TRANSACTIONS CANCELLED	00001350	135	130
	00001360	136	13!
	07510000	134	132
REAL TPARH(42)	0000138n	136	133
	00001390	081	134
DATA TPARM	00001400	146	135
1 2 3 4 5 6 7	00001410	iti	136
GRADE CHANGE	00001420	14.7	137
1 / 60., 10., 2.89, 6.54, 3.8, 32.3, 21.,	00001 430	143	156
DUTY STATUS CHANGE	00001440	771	66.

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STANDARD INPUT STREAM LISTING			
1 DATE 09/10/77			
INPUT CARD IMAGE2020303040506070	80	IMPUT ALT INPUT	MEMLE F. KEF-NU
2 63., 8., 3.40, 6.40, 11.1, 26.4, 7.,	00001450	\$45	140
C IMQUIRY	00001460	346	141
3 82., 3., 1.75, 10.53, 18.0, 27.7, 0.,	00001410	127	343
C ARRIVAL	00001480	148	143
59., 9., 2.00, 2.18, 16.9, 32.3, 16.,	06001460	671	75!
C DEPARTURE	00001500	1:0	145
5 57., 12., 2.60, 4.36, 13.7, 27.9, 18.,	00001510	151	34.
C ADD INDIVIDUAL	00001520	152	147
6 0., 70., 4.00, 1.43, 0.0, 20.5, 0. /	00001530	153	146
3	07 310000	711	64:
EQUIVALENCE (8SAVE(5), IP), (8SAVE(6), SEED), (8SAVE(7), SUB)	00001550	16.5	150
3	01613010	95.	151
END: DECLARATIONS:	00001570	167	152
3	00001580	, c.p	153
SEED = 1	06510000	1,50	154
00 10 I = 1,10	000010000	ikn	155
10 TLOCIII = 0	00001010	16,	156
2	00001050	291	157
END: INITIALIZATION;	06301630	:43	158
SEIZE: FACILITY=TRANS;	00001640	164	6:1
	00001650	392	160
C SELECT TYPE OF TRANSACTION	00001000	, , ,	161
AN = SUNFM(0.0, 1.0, SEED)	00001670		162
RP = SPWLIN(PUINTS,7,RN)	00001660	, 60	163

DATE 06/10/27				
INPUT CARD IMAGE2020304040506070.	60	SEC-11C	LEV SEG-RG	PFPEF REF-NL
IP = RP	06910000	169		104
	00001700	04:		165
50 SUB = ((IP-1)+7)+1	01710000	14.		166
INVOKE: SERVICE - VALIDS, PARAMETER LIST-(SEED, SUB, TPARM);	00001720	113		101
WAIT RETURN: SERVICE=VALID;	06000130	£41		166
1F(YLOC(8) .EQ. 1) GO TO 950	00001740	7Li		165
YRANSACTION IS NOT CANCELLED	00001750	544		74
GO TO (100,200,300,400,500,600),IP	00001760	91:		111
	00001770	111		172
START GRADE CHANGE PROCESSING	0961760	179		173
SO CONTINUE	06210000	041		76.
INVUKE: SERVICE=GRCHS, PARAMETER LIST=(TPARMI4),SEED);	00001860	06.		11:
WAIT RETURN: SERVICE=GRCH;	000001810	idi		:70
CD TO 900	00001820	cdi		ui
	00001830	183		17.6
	00001640	441		119
START DUTY STATUS CHANGE PROCESSING	00001650	531		180
200 CONTINUE	00001660	186		1+1
INVOKE: SERVICE=DIVCS. PARAMETER LIST=(TPARM(11), SEED);	00001670	131		162
WAIT METURN: SENVICE-DIVC:	05001680	JHE		163
60 10 900	06010000	684		.8.
	00001400	USI		19;
START INQUIRY PROCESSING	01+10000	:0:		166
300 CONTINUE	00001420	201		167

TATUMENT TO STATE TO STATE TO STATE TO STATE STA	IULA TOR			PAGE	•
STANDARD INPUT STREAM LISTING					
DATE 39/10/17					
INPUT CARD 1MAGE1020303040506070	7080	INPUT SEC-PO	ALT INPUT	MEMUCK N: F-T:0	
INVOKE: SERVICE=INUS, PARAMETER LIST=(TPARM(16), SEED);	00001730	163		14.6	
MAIT RETURN: SERVICE=ING;	00001540	751		189	
60 To 900	00001950	10:		193	
3	09610000	901		161	
C START ARRIVAL PROCESSING	00001970	161		75!	
400 CDNTINUE	00001480	108		143	
INVOKE: SERVICE=ARRIVS, PARAMETER LIST=(TPARM(25), SEED);	06610000	661		154	
WAIT RETURN: SERVICE=ARRIV:	00050000	200		195	
G0 T0 900	000002010	102		36.	
3	0000000	202		٢٠:	
C START DEPARTURE PROCESSING	000005030	202		467	
500 CCNT1NUE	00002040	702	And the second s	661	
INVOKE: SERVICE=DEPTRS, PARAMETER LIST=(TPARM132), SEED);	00000000	205		007	-
WAIT RETURN: SERVICE=DEPTR;	00005000	306	The second of the second of	107	
60 T0 900	0000000	202	The second secon	707	
3	00002060	JUE		303	
C START AUDITION PROCESSING	00002090	500		* 57	
600 CONTINUE	00002360	210		505	
INVOKE: SEMVICE=ADDSLS, PARAMETER LTST=(TPARM(39),SEED);	000002110	146		90.	
WAIT RETURN: SERVICE=ADDSL;	00000120	212		157	
60 10 900	000002130	٠١٢		305	
3	00002140	216		50.7	
900 CONTINUE	00002150	215		210	
C CHECK FOR COMPATABILITY ERROR IN TRANSACTION	00002160	912			

TARREST TOTOLOGICAL STREET					
STANDARD INPUT STREAM LISTING					
TLV01/30 09/10/11					
INPUT CARN IMAGE	٠٠٦٥٠٠٠٠٠٠6٠٠	INPUT SEC-60	LEV SEC-1.U	KEPEER KEF-NG	
END: INITIALIZATION;	00005410	24.		534	
SETZE: FACTLITY=VALID;	000005420	242		435	
2	00005430	543		236	
C GET PROBABILITY OF DIRECT MODE OF OPERATION	00005446	244		153	
DM = DARK (\$U8+01/160.	09726303	245	The state of the s	736	
C GET NUMBER OF DATA ITEMS THIS TRANSACTION	00002460	346		587	
NUMDI = DARR(SUB+1)	00002470	247		24.5	
C GET PROBABILITY OF VALIDITY ERROR	00002480	949		241	
PVERR = UARR(\$Ub+21/100.	00002490	576		275	1
C PVERR = 0.	00075000	Uac		243	
C GET PROBABILITY OF COMPATAPILITY FRAOR	01520000	25.1		.442	
PCERR = DARR (SUB+31/100.	00002520	252	The second secon	34.2	
C PCERR = 0.	00002530	243		240	
C GET DIRECT MODE PROCESS TIME IN MSFC	00002540	254	and the first the second of th	24.7	
PTDH = DARR(SUB+4) + 1000.	00002550	546		8	tez)
C GET TUTORIAL MODE PROCESS TIME IN MSEC	0000560	112		44.	
PTTM = DARR[SUB+5] * 1000.	000002570	757		943	
C GET PRIBABILITY THAT THIS TRANSACTION WILL BE CANCELLED	00002580	3.5		.65	
PCANCL = DARRISUB+61/100.	060002590	240		24,	
3	00005600	260	to the second of the second	617	1
C SELECT INPUT MUDE LEITHER DIRECT OR TUTORIAL)	000002010	.92		4.	
FN = 1UNFM (0.0, 1.0, SEED)	C0C12620	292		352	
1F(KN .GT. DH) GO TU 50	00002650	76.1		33.5	
C DIRECT HODE OF UPERATION	00002640	264		153	

DATE 09/10/77				
INPUT CARD IMAGE20304040506070.	60	INPUT SFC-NO	ALT ILPUT LEV SEG-MG	KFMcf +
C WAIT FOR JATA TO BE ENTERED (SYSTEM HOOK)	00005050	265		- 45 <i>t</i>
PROCESS: TIME=PTDM;	00002669	266		462
C DETERMINE 1F A VALIDITY EPPOR HAS OCCURPED ON THE DATA 11EM	00002670	267		260
KN = \$UNFH(0.0, 1.0, SEED)	00002680	24 A		197
IFIRM .GT. PVERR) GU TO 40	00002690	549		297
C VALIDITY ERROR - WAIT FOR CORRECTED INPUT (SYSTEM HOOK)	00002700	vLč		79:
PROCESS: TIME=PTOM;	000002710	112		777
C NO VALIDITY ERROR	00002720	212		59:
40 CUNTINUE	00002730	i12		777
GU 10 1GO	07420000	274		707
3	000002750	575	and the second s	392
3	00002760	14.2		677
C TUTORIAL MODE OF OPERATION	06052770	146		21.
50 J = I	00000780	274		11.7
69 IF(J.CT. NUMDI) GO TO 100	00002750	510		71.2
C WAIT FOR INPUT FROM TERMINAL (SYSTEM HOOK)	00002100	045		51.7
PROCESS: TIME=PTTM;	00002816	261		515
C DETERMINE IF VALIDITY ERROR HAS OCCURRED FOR THIS DATA ITEM	00002620	242		:11;
KN = \$UKFH(6.0, 1.0, SEED)	00002830	283		31.
IFIRM .CT. PVERRY GO TO 70	00002840	702		1113
C VALIDITY ERROR - WAIT FOR CORRECTED INPUT (SYSTEM HOOK)	00002850	285		47.5
PRUCESS: TIME=PTTM;	00002860	366		61.7
70 J = J + 1	00007670	787		0.7
GD 13 60	00002680	2.69		: 4:

STANDARD INPUT STREAM LISTING			
DATE 09/10/77			
THPUT CARD IMAGE20304050	Iniblit 160-17	T ALT INPUT	MEME FR KIF-NC
5	00002850 289	0	297
C DETERMINE IF TRANSACTION IS TO BE CANCELLED	00002000	000	283
C IF CANCELLED THEN RETC=1, OTHERWISE RETC=0	00002910 291		264
100 RM = SUNFMIO.0, I.O, SEED)	00002920 292	.2	:85
TLCC(R) = 0	00002930 293	3	787
IFIRM .LY. PCANCL) TLOC(8) = 1	00002540	7	287
WELEASE: FACILITY=VALID;	\$62 05620000	.5	286
END: ENDO SERVICE;	90003560	9	285
	5.65 00002570	· · · · · · · · · · · · · · · · · · ·	
	08620000	a	
	00620000	0	
ENDO SERVICE: 10=GRCH, NAME=GRCHS, SAVE AREA SIZE=10,	00000000	0	047
PARAMETER LIST=(PFRRI,SEEDI);	01060000	-	167
CDMHON ZWKZ TLOC	02050000	2	26?
INTEGER SUNFMI, SEED, SEEDI, TLOCITO)	\$75 0E0E0000	Ė	£6.
INTEGER OBTAN, FINDR, MODIF, STOR, CONCT, DISCN, ERAS	00000000	7	252
INTEGER DHY, NEXT, PRICH, DUNER	20t 02050	3	295
INTEGER ADDI, ADUN, ALOC, GRADE, TOEN, INST, MOSC, POSN, UIST, UNMA	09050000	Y	96,7
INTEGER CALC, GRUI, IDUI, INUM, PMID, PNUI, SMID, UNAM, UNGR, UMPOS, UMUI	00003070	-	14.
FOULVALENCE (SSAVE(1), PERR), (SSAVE(2), SEEU)	3080	a.	257
REAL SUNFM, RAN, PERRI	06050000		547
DATA UBTAN,FINDK,MODIF,STOP,COMCT,DISCN,ERAS/1,2,3,4,5,6,7/	00160000	c	90
DATA DMY,NEXT,PRICK, DWNERZ6,1,2,3/	01:60000		3.1
DATA ADDI, ADUN, ALOC, GRADE, IDEN, INST, HASC, POSN, UIST, UNHA	00003120 312	3	

P46t 14

INPUT CARD IMAGE2030404050607060
5. 6. 7. 8. 9. 10/ 00003130
DATA CALC,GRUI, IDUI, INUM,PHID,PCUI, SHYD,UMAM,UMCR,UMPOS,UMUI 00003140
5, 6, 7, 8, 9, 10, 11/ 00603150
000031.00
04180000
00663160
06;€7000
00003200
RETURN CODE -TLOC(I)- = 0 IF NO ERROR, I IF ERROR
06063220
- OBTAIN=1.FIND=2.MNDIFY=3.STORE=4.CONNECT=5.
RASF=7 60003240
PTOR=2,0WNER=3 00003250
21 RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS 06063260
01250000
266632290
06787990
INVOKE: SERVICE= INTES, PARAMFTER LIST=(DBTAN, DMY, IDEN, CALC); 000003300
01660000
INVOKE: SERVICE = INTF1, PARAMETER LTST=(OBTAN, DMY, ADDI, CALC); 00003320
JE3E007J
INVOKE: SERVICE= INTES, PARAMETER LIST=106TAN, DMY, MOSC, CALC); 00003340
U\$E £0000
INVOKE: SERVICE - INTES, PARAMETER LTST=(FINDR, MEXT, UIST, IDUI); 000005360

DATE 09/10/77			
INPUT CARD IMAGE 30605060	70F0	INPUT ALT INPUT	MENEER FEE-LO
WAIT RETURN: SERVICE - INTF;	00003610	363	195
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR, DMY, I DEN, CALC);	00003620	345	346
WAIT RETURN: SERVICE=INTF;	00003€30	363	555
INVOKE: SERVICE = INTES, PARAMETER [IST=(CONCT, DMY, IDEN, PMID);	000003640	716	1.5
WAIT RETURN: SERVICE=INTF;	00003650	394	193
	00003560	346	356 7
C DISCHMECT THE PERSON'S UIST RECORD	00003670	198	S BE
	00003660	346	Hot.
INVOKE: SERVICE INTES, PARAMETER LIST=(FINDR,NEXT,UIST,10UI);	000032690	:40	T
MAJT RETURN: SENVICE=INTF;	00003300	746	یان پر
INVOKE: SERVICE INTES, PARAMETER LIST=(FINDK, OWNER, DMY, GAUI)	00003710	145	36:
WAIT RETURN: SERVICE=INTF;	G00C3720	37.2	362
INVOKE: SERVICE INTES, PARAMETER LIST IDISCN. DMY, UIST, CRUI);	05750000	. درد	5.00
WAIT RETURN: SEKVICEEINTF;	09663740	718	264
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR, NEXT, UIST, IDUI);	00003750	346	30.5
WAIT RETURN: SERVICE=INTF;	06003760	378	364
INVOKE: SERVICE = INTES, PARAMETER LIST=(FINDR, OWNER, UIST, GRU1);	000003770	11.8	7.30
WAIT RETURN: SERVICE=INTF;	CC00378C	818	366
INVOKE: SERVICE INTES, PARAMETER LIST=(DISCR.DMY.UIST,CRUI);	060003790	516	600
WAIT RETURN: SERVICE=INTF;	00003600	Udi	0,1
2	00003810	loc	:7:
CONNECT UIST RECORDS TO THE GRADE-UIST SET OF THE NEW GRADE	04003820	246	:1.
3	00003130	k.)t.	17.
INVOKE: SERVICE= INTES. PARAMETER LIST=(FINDR,NEXT,UIST, 1001)	00003840	775	27.

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	¥0.		FAGE 17	
STANDARD INPUT STREAM LISTING				-
DATE 09/10/77				1
INPUT CARN IMAGE	04	THEUT ALT INPUT	MERCER Fre-fo	
WAIT RETURN: SERVICE-INTF;	00003150	385	575	1
INVOKE: SERVICE = INTES, PARAMETER LIST=(FINDR,OMY, IUEN,CALC);	000038-60	386	575	1
WAIT RETURN: SERVICE=INTF;	00003670	387	115	
INVOKE: SERVICE= INTES, PARAMETER (IST=(08) AN, NEX 1, GRADE, UMGR);	00003860	349	376	
WAIT RETURN: SERVICE=INTF;	06950000	być	97:	
INVOKE: SERVICE - INTES. PARAMETER LIST = (FINDR, DMY, UIST, CALC);	00660000	05 è	046	1
WAIT RETURN: SERVICE=INTF;	01080000	391	381	1
INVOKE: SERVICE = INTES, PARAMETER LIST=(FINDR, OWNER, DMY, GAUI);	00003920	345	382	
WAIT RETURN: SERVICE=INTF;	000003930	141	585	1
INVOKE: SERVICE INTES, PARAMETER LIST=(CONCT, DMY, UIST, GRUI);	000003540	374		1
WAIT RETURN: SERVICE=INTF;	09660000	368	285 285	m o
INVOKE: SERVICE = INTF4, PARAMETER LIST=(FINDR,NEXT,UIST, IDUI);	00003360	346		Y-K
WAIY RETURN: SERVICE=INTF;	06650000	307	100	1
TAVORE: SERVICE INTES, PARAMETER LIST = (CONCT, DMY, UIST, GRUI);	00003980	394	TQ	- 1
WAIT RETURN: SERVICE=INTF;	066 €0000	65£	100 588	Hade
C NORMAL RETURN	000040000	707		10
90 7LOC(11) = 0	01040000	403	Phi	מניי
66 10 150	02046000	462	342	
C COMPATABILITY ERROR	05047000	703	242	1-0
100 CONTINUE	00000000	707	300	
1(0011) = 1	05040000	405	355	•
150 CONTINUE	00004000	404	.96	1"
RELEASE: FACTLITY=GRCH;	31940330	77	14:	
END: ENDO SERVICE;	00004680	4Ú7	***	•

THE COMMON TAKET TOWN CARD TAKEN TEACH TOWN CARD TAKEN TOWN CARD TAKEN CARD TOWN CARD TAKEN CARD TAKEN CARD TOWN CARD TAKEN CARD TAKEN CARD TOWN CARD TAKEN CAR					
THERE SECTION ONDOGATION 417 15-017 NEPLEY SECTION 11-017 NEPLEY 11-017					
THE PROPERTY SECURITY CONTINUES	80		1	EPUE:	
######################################	06040000	404			
######################################	00004100	7.50			
LIST= PERRI, SEECI]; 00004120 412 599	00004110	411			
LIST=IPERRI,SEED1); 00004140 414 406 401 402 403 404 405 405 407 408 408 408 408 408 408 408	00004120	412		665	
	00004130	£17		007	
######################################	00004140	414	1	101	
F.5YOR.TONCT.DISCN.ERAS 00004160 4:5 40.4 WHER NUM.PHID.POINT.SHID.UMAW.UMGR.UMPUS.UMUI 00004210 4:7 4.0 RR1. (\$\$\text{AVE(12)}.\$\text{SED}) 00004210 42: 40.6 CCC42.0 42: 40.6 CCC42.0 42: 40.6 CCCC42.0 42: 40.6 CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC	00000150	517		405	
#MER #ADF., 10FW, 14ST, 440SC, POSN, UIST, UNMA	000004160	9.7		463	GE I
RADE, 105 10	000004:70	437		+0+	SBE
RRI, (\$\forall \text{FRID} \te	00004180	418		405	LSH.
### (#\$AVE(2), SEED) ##################################	0015000	610	-	705	AUP T CE
TOR.CONCT.DISCN.ERAS/1,2,3,4,5,6,7/ 60004210 422 452 459 E.1DEN.TNST.HMSC.POSN.UIST.UNMA CCCC4241 423 420 4. 5. 6. 7, 8, 9, 10/ 6604250 426 412 4. 5. 6. 7, 8, 9, 10/ 11/ 00664260 426 412 4. 5. 6. 7, 8, 9, 10, 11/ 00664260 426 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415 6. 6664260 427 415	00004500	967		137	O DD
00004220 452 465 00004220 423 410 00004250 424 412 00004250 424 413 00004260 424 414 00004260 427 414 00004260 427 415 00004260 427 415 00004220 430 417	000004510	421		404	PR
00004230 423 4.20 CCCC4241 424 CGCC425C 424 413 00004260 426 413 00004260 426 415 00004260 427 415 00004220 430 415	00004220	45.2		404	ACT
COCC4241 424 412 COCC425C 424 412 COCC42CO 424 413 COCC42CO 427 414 COCC42CO 427 415 COCC42CO 430 427 COCC642CO 430 427 COCC642CO 430 427	00004230	423		410	I Sad
00004250 424 00004260 424 00004260 427 00004260 427 00004260 427 00004220 430	(7247303	767		-111	
00004260 424 7 00664270 427 00064250 426 02664250 430 06604320 432	60604250	567		412	
3, 4, 5, 6, 7, 8, 9, 10, 11/ 00664;70 427 i	09740000	424		413	
\$\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\	000004:70	<u>L</u> 23		414	
00004250 430 00004330 433 00004330 433	00004260	927		415	
00004520 430 00004520 433	06270000	K27		416	
000004310 431	00004300	430		4.1	
000004320	01540000	167		47.8	
	000004320	432		517	
		00004150 00004150 00004160 00004160 00004160 00004210 00004210 00004210 00004210 00004210 00004210 00004210 00004210 00004210 00004210			7.5 7.5 7.5 7.5 7.5 7.5 7.5 7.5

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WINDLY BURNEY ON TO SUCCESSION WITH CONTROL OF THE	¥.			1914	
STANDARD INPUT STREAM LISTING					
DATE 09/10/77					
INPUT CARD IMAGE1020304050606070	69	INPUT SF0-M	ALT MAPUT	MENEFF.	
3	00004330	433		775	
C RETURN CODE -TLOC(2)- = 0 IF NO FRROR, 1 IF ERRUR	000004340	767		421	
2	00004350	587		777	
C PARAMETER LIST CONTAINS-	00004360	967		423	
C 1) UML TYPE - OBTAIN=1,FIND=2,MODJFY=3,STCRE=4,CONNECT=5,	00004370	167		727	
C DISCONNECT=A,ERASE=7	00004380	438		425	- 0
C 21 GUALIFIER - DAY=0.NEXY=1.PPIOR=2.CHINER=3	000004390	430		474	
C 21 RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML 1S	00004400	077		177	18
C TO UPERATE	00004410	(4)		456	IS P
C 4) SET IDENTIFIER - INDEX	000004420	142		677	GE
3	00034430	277		354	15
C COMPATABILITY CHECK	00004440	747		431	BES!
3	00004450	577		754	T QU
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN, DMY, I DEN, CALC);	00004460	797		453	10
WAIT RETURM: SERVICE=INTF;	06004470	L77		767	DOC
INVUKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN, DMY, ADDI, CALC);	000004480	448		+35	-
WAIT RETURN: SERVICE=INTF;	064440000	577		430	11.14
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN-NEXT, UIST, IDUI);	00004200	450		11.5	are.
MAIT KETURN: SERVICE=INTF;	00004510	157		-38	4
INVOKE: SERVICE INTES, PAPAMETER LIST 108TAN NEXT, UIST 1001);	00004520	755	The second secon	439	
WAIT RETURN: SERVICE=INTF;	000000	£37	and the last sea to be a second or the second of the secon	044	
INVOKE: SFRVICE= INTF1, PARAMETER LIST=(OBTAN,CMNER,CMY,UMUI);	000004540	757		17.	
WAIT RETURN: SERVICE=INTF;	05450000	557		24.7	
TANCING C CENTICE TAYER DASSAFFED STATE DASSAFFED STATES	00004560	456		**	

STANDARD INPUT STREAM LISTING					
DATE 09/10/77					
INPUT CARD IMAGE203030405050	08	INPUT SEC-NO	ALT INFUT	REF-NO	
WAIT RETURN: SERVICE=INTF;	000004570	157		144	
C DETERMINE IF COMPATABILITY ERROR HAS OCCURRED	00004580	458		4.5	
RN = SUNFMIO.0, 1.0, SEED!	00004550	654		446	
IF(RN .LT. FERR) GO TO 100	00004000	760		144	B1
C NO COMPATABILITY ERRUR	00004610	(37		347	BOM .
	02940000	462		440	PAGI
C MODIFY THE IDENT, UIST AND ADD-UNIT RECORDS	00004630	163		037	~
3	000004040	797		150	NAL:
INVOKE: SERVICE= INTFS, PARJHETER LIST=(FINDR,DMY,IDEN,CALC);	03940000	465		757	SHE
WAIT RETURN: SERVICE=INTF;	00004000	466		:53	TO
INVOKE: SERVICE= INTFS, PARAMETER LIST=(MODIF, DMY, IDEN, DMY);	00004670	194		**	DDC
WAIT RETURN: SERVICE=INTF;	00004000	446		45.5	FRA
INVOKE: SERVICE= INTFS, PARAHETFR LIST=(FINDR, DMV, UIST, CALC);	06970000	460		456	CTI
WAIT RETUKN: SERVICE=INTF;	00004100	470		157	-
INVOKE: SERVICE= INTES, PARAMETER LIST=(MODIF, DHV, UIST, DMY);	00004710	471		854	-
WAIT RETURN: SERVICE=INTF;	07450000	247		5 44	
INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR, DMV, ADUN, CALC);	00004730	173		460	
WAIT RETURN: SERVICE=INTF;	00004140	7.77		-61	
INVOKE: SERVICE= INTES, PARAHETER LIST=(MODIF, DMV, ADUN, DMY);	00004750	517		29.	
WAIT RETURIS SERVICE=INTF;	00004760	915		463	
C NORMAL RETURN	00004770	117		424	
90 TLOC(2) = 0	00004783	067		465	
GG TO 155	0604790	643	The same of the sa	466	1
C COMPATABILITY FRANK	00000	440		147	

1.7 JOHA		SEC-NO LEV SEG-TO KEF-TO	468	634 766	483 470	767	547	7H2		7.96 V.7.5 V			3 4 567 157	1 N 9/2 767	463	127 707	527 507	450	461	2.17	699	10,	cu) 485	947	
NO.		UA.	00000	0787000	00004430	00004840	06440000	09870000	0144000	00004480	00004890	0000000	01440000	00004526	000004630	00000000	05640000	09040903	01640000	66064510	06540000	00040000	01050333	00000000	0000.50.00
STANDARD INPUT STREAM LISTING	DATE 09/10/77	INPIT CARD IMAGE	100 TLOC(2) = 1	150 CONTINUE	RFLEASE: FACILITY=DTYC;	END: ENDO SERVICE;				ENDO SERVICE: ID=ING, NAME=INOS, SAVE AREA SIZE=10,	PARAMETER LIST=(PERRI, SEEDI);	COMMON /WK/ TLOC	INTEGER SUNFHI, SEED, SEEDI, TLAC(10)	INTEGER OBTAN, FINDR, MODIF, STOP, CONCT, DISCN, ERAS	INTEGER DMY, NEXT, PRIOR, DWNER	INTEGER ADDI, ADUN, ALOC, GRADF, FISH, INST, MUSC, PUSN, UIST, UNMA	INTEGER CALC, GRUI, IDUI, INUM, PMID, PNUI, SHID, UMAM, UMGR, UMPOS, UMUI	EQUIVALENCE (\$SAVE(1), PERR), (\$SAVE(2), SEEC)	REAL SUNFM.RN.PERR,PERRI	DATA OBTAN,FINDK, MUDIF, STOR, CONCT, DISCN, ERAS/1,2,3,4,5,6,7/	DATA DHY, NEXT, PRIOR, DWNER/0, 1, 2, 3/	DATA ADDI, ADUN, ALGC, GRADE, 10FN, 1WST, MMSC, PMSN, UIST, UNMA	1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/	DATA CALC.GRUI.INUM.PHID.PORT.SHID.UHAW, UMGR.UHPOS.UMUI	1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/

DATE 09/10/77				
INPUT CARD IMAGE20304050606070.	na	INPUT P	ALT JIPUT MEMBEN	PE K
PERR = PERRI7100.	04940000	÷0;	3	594
SEED = SEEDI	09050000	908	54	044
END: THITIALIZATION;	0000000	2013	57	167
SEIZE: FACILITY=IKO;	00005000	e05	57	7.67
	06000000	695	4	664
C RETURN CODE -TLOC(3)- = 0 IF NN ERROR, I IF ERROR	00150000	610	54	767
2	00005110	liii	57	544
C PARAMETER LIST CONTAINS-	00005120	215	5.9	496
C 1) DHL TYPE - 06TAIN=1,FTRN=2,MODIFY=3,STORE=4,CONNECT=5,	000005130	613	57	1815 Feor
C DISCONNECT=6,FRASE=7	00005140	514	54	S PA M ON
C 2) QUALIFIER - DHY=0.NEXT=1.PPIOR=2.OWNER=3	00005150	91.5	*	X 054
C 21 RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML 1S	00005160	616	3;	200
TO OPERATE	000003170	517	33	105
C 41 SET IDENTIFIER - INDEX	00005180	A ! ?)	7.09
	06150000	615	05	503
INVOKE: SERVICE= INTES, PARAMETER LIST=[OBTAN, DHY, I DEN, CALC);	00750200	953	9:	-0:
WAIT RETURN: SERVICE=INTF;	00005210	521	\$0°	
INVOKE: SERVICE INTES, PARAMETER LIST=(ORTAN, DMY, INST, CALC);	00005220	£22	0;	90;
WAIT RETURN: SERVICE=INYF;	00005230	623	0.	100
INVOKE: SERVICE= INTES, PARAMETER LYST=(OBTAN, DMY, AUDI, CALC);	09750000	524	0;	90,
WAIT RETURN: SERVICE=INTF;	00005250	365	,0;	,,,,,
INVOKE: SEKVICE= INTFS, PARAMETER LTST=(OBTAN,NEXT,UIST,1001);	00005260	925	2	vij
WAIT RETURNS SERVICE=INTF;	0053570	125	W	
INVOKE: SERVICE = INFFS, PARAMETER LIST=(OBTAN, DOSN, CALC);	00005280	624	316	

The same same shade and the same same same same same same same sam		AIT INDUT PEMBER	515	514	515	510	714	518	61;	625	221	225		725	525 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5 5	950				275	\$25	\$25	130	15.	532	
		SEL-NO LEV S	429	530	163	265	513	534	363	165	753	#£3	bii	075	175	542	643	775	242	975	۲۳۶	472	673	دد ن		
			06250000	00035300	00000310	00005320	00005330	00005340	00005350	00005360	00005370	00005380	00002390	00002400	00005410	00002450	00002430	00005440	000005450	00002460	000005470	00005480	05750000	00005500	01550000	
STANDARD INPUT STREAM LISTING	DATE 09/10/77	INPUT CARD IMAGE	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, ALOC, CALC);	WAIT RETURN: SERVICE=INTF;	C DETERMINE IF COMPATABILITY FARDP HAS OCCURRED	RN = SUNFHIG.0, 1.0, SEED)	IFIPERR .LT. RNJ GO TO 100	C NORMAL REYURN	90 YLDC(3) = 0	GO TO 150	C COMPATABILITY ERROR	100 YLOC(3) = 1	150 CONTINUE	RELEASE: FACTLITY=ING;	END: ENDO SERVICE;				ENDO SERVICE: ID=ARRIV, NAME=ARRIVS, SAVE AREA SIZE=10,	PARAMETER LIST=(PERR! ,SFED1);	COMMON /WK/ TLOC	INTEGER SUNFMI, SEED, SEEDI, TLOCIIOF	INTEGER CBIAN, FINDK, HODIF, STOR, CONCT, DISCN, FRAS	INTEGER DNY, NEXT, PKICK, DUNER	

DATE 09/10/77				
INPUT CARD THAGE 302030304040506070	80	THEUT ALT TREUT	TREUT MERLEH	
INTEGER CALC.GRUI.IDUI.INUM.PHTD.POUI.SHID.UMAW.UMGR.UMPOS.UHUI	00005530	£53	753	HI
EQUIVALENCE (\$SAVE(1),PERR), (\$SAVE(2),SEED)	000005540	554	\$6.3	S PA
REAL SIMFM.RN.PERR.	00005550	¥ 5 4 4 4 4 4 4 4 4 4 4 4 4 4 4 4 4 4 4	530	GE I
DAYA DBYAN.FINDR. HUDIF. SYOR, CONCY, DISCN. ERAS/1, 2, 3, 4,5,6,7/	09550000	556	153	ES B
DATA DHY.NEXT.PRIGR.DWNEK/0.1.2.3/	06005570	Lis	136	BST LSh
DATA ADDI, ADUN, ALDC, GRADE, IDEN, TWST, MMSC, POSN, UIST, UNMA	00005560	F. 5.7	567	QUA
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/	96560396	550	0.05	TOD
DATA CALC, GRUI, IDUI, INUM, PHID, PMUI, SHID, UMAM, UMGR, UNPOS, UMUI	0000000	240	341	DC
1 / 10 20 30 40 50 60 70 80 90 100 11/	0000000	:95	74;	NAC?
END: DECLARATIONS:	00005620	245	243	164
PERR = PEKRI/100.	00005630	£43	747	PER
SEED = SEED1	00005640	564	245	
END: INIVIALIZATION;	05950000	575	353	
SEIZF: FACILITY=ARRIV;	09950000	566	175	
	0000000	L)S	44	
RETURN CODE -TLOCIST = 0 IF NO ERROR, 1 IF ERROR	00005580	976	540	
	0000000	54.9	013	
PARAMETER LIST CONTAINS-	0005000	670	153	
1) DML YYPE - OBTAIN=1, FIND=7, MODIFY=3, STORE =4, CONNECT=5,	00005710	145	744	
DISCONNECT=6,ERASE=7	02720000	512	553	
2) QUALIFIER - DAY=0,NEXT=1,PRIOR=2,OWNER=3	06720000	eli	I.	
2) RECOKO TYPE - 105ATIFIER OF THE RECORD UPON WHICH THE DAL 15	00005740	74.9	304	
TO OPERATE	09153000	SLi	364	
4) SET IDENTIFIER - INDEX	00005760	576	155	

Consultant of

No. of Lot,

STANDARD INFUT STREAM LISTING					
INPUT CARD IMAGE	80	INPUT SEC-110	ALT IRPUT LEV S! G-NO	MEPEL S. ELF-NO	
	06205770	113		966	
INVOKE: SERVICE - INTES, PARAMETER LIST=(Db1AN,DMY,1DEN,CALC);	00005780	815		655	
WAIT RETURN: SERVICE=INTF;	061 50000	013		075	
INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, INST, CALC);	00002000	0 85		19;	
WAIT RETURN: SERVICE=INTF;	00005810	581		794	
INVOKE: SERVICE INTES, PARAMETER LIST = (OBTAN, DMY, ADDI, CALC);	00005620	245		563	
HAIT RETURN: SERVICE FINTF;	00005830	613		*95	
INVOKE: SERVICE = INTES, PARAMETER LIST=(DBTAN, DMY, POSN, CALC);	07950000	755		391	
WAIT RETURN: SERVICE=INTF:	00005850	593		996	
INVOKE: SEKVICE= INTES, PARAMETER LIST=(OBTAN,NEXT,UIST,10UI);	C0005860	586		294	
WAIT RETURN: SERVICE-INTF;	00002670	587		568	
INVOKE: SERVICE= INTF4, PARAMETER LIST=108TAN,NEXT,UIST, TOUI);	06005680	G D u		404	
WAIT RETURN: SERVICE-INTF;	0000 \$860	049		670	
INVOKE: SERVICE INTES, PARAMETER LIST - TOBTAN, DMY, UNMA, CALC);	000355000	005	-	11.6	
WAIT RETURN: SERVICE=INTF;	01650000	105		21.5	
INVOKE: SERVICE = INTES, PARAMETER LIST = (OBTAN, DMV, ADUN, CALC);	00005520	213		.73	
WAIT RETURN: SERVICE-INTF;	0000530	693	The state of the s	514	
INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, ALOC, CALC);	000000	703	And the second section of the second section is a second section of the second section	57.5	
WAIT RETURN: SERVICE=INTF;	06450000	Sos		37.5	
KN = SUNFMIO.0, 1.0, SEED)	00005460	Jos		111;	
IFIRM .LT. PERRI GO TO 100	01450000	Lús		576	
	28560300	80%		27.5	
END COMPATABILITY CHECK - NO COMPATABILITY ERROR DETECTED	06550000	603		095	
MODIEY TORNY AND ADD-IDENT RECORDS	00090000	YCU		193	

		Thru MITER SEG-NO KEF-NO	582	583	554	THI PRO	S P.	265 741	397	ALS:	QUANTE OS.	TO D	PI PI	SAC:	756	555	956	142	844	65.4	600	103	500	٠٥.	793	
		INPUT ALT	601	£42	603	604	605	903	€04	809	609	9.0	611	612	613	414	615	616	617	618	619	127	1,29	424	+23	434
		03	01090000	0000066.20	000000000	04090000	0000000	09090000	00000000	00000000	06090000	00:90000	000000110	000000120	26190000	06000240	000000150	000001100	06036170	0000001 00	06:40000	0000005000	00006210	000000250	000000230	0000000
STANDARD INFUT STREAM LISTING	DATE 09/10/77	INPUT CARD IMAGE		INVOKE: SERVICE= INTES, PARAMETER LIST=(DBTAN, DHV, ALOC, CALC);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LYST=(FINDR, DHV, IDEN, CALC);	WAIT RETURN: SERVICE - INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(MODIF, DHV, IDEN, DHY);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LYST=(OBTAN, DHV, ADDI, CALC);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(MODIF, DMV, ADDI, DMY);	WAIT RETURN: SERVICE=INTF;	3	C STORE THE NEW UIST RECORD AND ITS CONNECTION IN THE IDENT-UIST	C SET AND IN THE UNITHIT-DIST SET	5	INVOKE: SERVICE= INTES, PARAMETER LIST=(SYOR, UMY, UIST, DMY);	WAIT RETURN: SERVICE=INTF;	INWIKE: SERVICE= INIFS, PARAMETER LIST=(GETAN, DMY, GRADE, UMGR);	WAIT RETURN: SERVICE=INTF;	INVOKE: SEKVICE= INTES, PARAMETER LYST=(CONCT,DMV,UIST,POUI);	WAIT RETURN: SERVICE=INTF;	INVOKE: SEKVICE INTES, PARAMETER LIST=(DBTAN,DMV,UNMA,CALC);	WAIT RETURN: SERVICE=INTF;	VEUTRE CLEDITE INTEL BABANETED IVEY TOWN THE THEIR

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			374	
STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD IMAGE	0,	INPUT ALT INFUT	PEMF CA	
WAIT RETURN: SEAVICE=INTF;	00006250	625	606	
INVUKE: SERVICE= INTF *, PARAMETER LIST=(CONCT.DMY, UIST, GRUI);	00000540	626	103	
WAIT RETURN: SERVICE=INTF;	00000570	627	407	1
RN = \$UNFMIG.0, 1.0, SEED)	000000	629	5,33	
IFTRN .67. 0.201 GO TO 100	000006290	629	616	
INVOKE: SERVICE - INTES, PARAMETER LIST=(DISCN,DMY,UIST,POUL);	00000000	630	611	
WAIT RETURN: SERVICE=INTF;	000006310	(;)	ci2	1
INVOKE: SERVICE - INTES, PARAMETER LIST=(DISCN,DMY, UIST, UMUI);	000000320	563	613	10
WAIT RETURN: SERVICE=INTF;	00000930	633	719	PHON
INVOKE: SERVICE= INTES, PARAMETER LIST=(DISCN.DMY, UIST, GAUI);	000000340	4.4	573	100
MAIT RETURN: SERVICE=INTF;	000006350	425	9.1	2V .
INVOKE: SERVICE= INTES, PARAMETER LIST=(ERAS,DMY,UIST,DMY);	00000360	626	117	i ka
WAIT RETURN: SEKVICE=INTF;	000009370	637	uSH e:s	451
C NORMAL RETURN	00000380	dět	5:0	QUA
90 TLOC(4) = 0	05690000	0;9	074	LIT
GU TO 150	000000	64.5	10.73	YPA
C COMPATABILITY ERROR	0000000	44!	1 223	ACI
100 CONTINUE	00000430	642	6.3	100
TLOC(4) = 1	000000	6443	-55-	-
150 CONTINUE	000006440	644	377	
RELEAS: FACILITY=ARRIV;	05.00000	4.65	377	
END: ENDO SERVICE:	(0000 400	tak	121	
	0000000	L73		
	04796000	479		

STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD IMAGE20304050606070	บล••••	JNPUT SEC-NO	ALT BREUT	MEMPEK REF-NU
	25490002	671		
ENDO SERVICE: ID-DEPTR, NAME-DEPTRS, SAVE AREA SIZE=10.	000000	450		624
PARAMETER LIST=(PERRI, SEEDI);	01590000	1.99		677
COMMON /WK/ TLOC	000006520	652		630
INTEGER SUNFMI, SEED, SEEDI, TLOCIIO)	0000000	653		ار الله عالم الله الله الله الله الله الله الله ا
INTEGER UBTAN.FINDR.MODIF,STOR,CONCT.NISCN,ERAS	00006540	459		634. 5
INTEGER DMY, NEXT, PRIOR, DWNER	0000000	653		633 5 5
INTEGER ADDI, ADUN, ALGC, GRADE, IDEN, INST, MOSC, POSN, UIST, UNMA	00006560	4.56		c34 7 52
INTEGER CALC, GRUI, TOUI, INUM, PHIC, POUI, SMID, UMAW, UMGR, UMPUS, UMUI	02590000	153		573
EGDIVALENCE (\$SAVE(1), PERR), (*SAVE(2), SEED)	000005 60	9.51		S TO
KEAL SUNFM, RN, PEKR, PERRI	00000000	659		100
DATA USTAN,FINDR,HUDIF,STOR,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/	00000000	660		638
DATA DMY, NEXT, PRIUR, DWNER/N, 1,2,3/	01990000	199		050
DATA ADDI.ADUN, ALDC.GRADE . IDEN, INST. HISC. POSN, UIST, UNHA	00016620	299		343
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/	00006630	663		171
DATA CALC, GRUI, TOUI, INUN, PWID, POUT, SMID, UMAN, UMGR, UMPOS, UMUI	000000000	949		279
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/	0000000	145		C43
END: DECLARATIONS;	0000000	949		950
PERR = PERRI/100.	00000c7C	647		(1)
SEED = SFEDI	00000680	448		17.3
END: INITIALIZATION;	05990000	640		£4.7
SEIZE: FACILITY=DEPTR;	006036763	670		7 t i
	01130000	1173		54.7
RETURN CODE -TLOC(5)- = 0 IF NO ERROR, I IF ERROR	02490000	172		16.3

	1108		1746	CF //
STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD IMAGE	03	JEPHIT ALT	ALT BEPOT NEWLEI LEV SEK-NU KEF-IK	
3	00006730	677	159	
C PARAMETER LIST CONTAINS-	07/90000	474	755	
C 1) DML TYPE - OBTAIN=1,FI40=2,MOD1FY=3,STORE=4,CCNNECT=5,	00006750	675	636	
C UI SCONNECT=6,ERASE=7	09296000	67.6	459	
C 2) OUALIFIER - DHY=0,NFXT=!,PRIOR=2,OWNER=3	0000000	11.3	559	7.2
C 2) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS	00006 789	678	9;9	Part (
C TO OPERATE	06290000	619	153	2027
C 4) SET IDENTIFIER - THOEX	000000	680	653	140
3	0000000	. 89	577	t and
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN, DMY, IDEN, CALC);	02493000	Ç693	046	1
WAIT RETURN: SERVICE=INTF;	000006630	693	197	ie
INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, INST, CALC);	00000640	767	29,	MOO
WAIT RETURN: SERVICE=INTF;	000000850	685	663	-
INVOKE: SERVICE = INTES, PARAMETER LIST=(OBTAN, DMY, ADDI, CALC);	00000000	9+9	539	
WAIT RETURN: SERVICE=INTF;	00006670	687	\$30	1
INVOKE: SERVICE INTES, PARAMETER LIST=(OBTAN, DMY, UNMA, CALC);	000008860	457	79,	
WAIT RETURN: SERVICE=INTF;	069990000	690	1.12	-
INVOKE: SERVICE = INTES, PARAMETER LIST=(ORTAN,DMY,ADUN,CALC);	00490000	009	5.93	
WAIT RETURN: SERVICE=ANTF;	000006410	601	500	
INVOKE: SERVICE = INTES, PARAMETER LIST=(OBTAN, DMY, UNMA, CALC);	07690000	400		
WAIT RETURN: SERVICE=INTF;	00000000	663	11.3	
INVOKE: SERVICE - INTES, PARAMETER (IST=(08TAN,DMY,ADUN,CALC);	00000000	404	67.5	
WAIT RETURN: SENVICE-INTF;	00000000	100		
INVOKE: SERVICE INTES, PARAMETER LIST=(OBTAN, DMY, ALOC, CALC);	000000	101	76.3	

(

		1 11 11.01 -10 1FV SEL-1.0			1					193	793	543	17.7	6.65	999	6+7	643	CHY	069	163	692	263	469	543	049	6.23	F 698
		INPINT SFO-PO		169	648	669	700	101	71.2	٤٥٠	706	705	346	101	708	100	١١٤	hir	11.2	213	714	غالا	7:6	414	716	oic	720
		E0	0.000	01690000	00000000	06690000	99010000	01070000	0201000	06057030	00001040	00001050	00007000	01610000	00007080	06070000	00110000	01110000	07140000	00007130	04170000	0512000	000001160	000001170	00007180	06117000	007.2000
STANDARD INPUT STREAM LISTING DATE 09/10/77	DATE	INPUT CARD IMAGE	OTAK BEVIELS CEDUVE - VINE	MAII KEIUKN: SERVICE=INIF;	C DETERMINE IF COMPATABILITY FRRME HAS OCCURRED	RN = \$UNFM(0.0, 1.0, SEFD)	IFIRM .LT. PERRY GG TO 100	C NO COMPATÁBILITY ERROK	INVOKE: SERVICE= INTES, PARAMETER LIST=(Obtan,NEXT,UIST,1DUI);	MAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN,NEXT,UIST,IDUI);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTF4, PARAMETER LYST=(FINDR, DMY, 1DEN, CALC);	WATT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LTST*(MODIF, DMY, IDEN, DMY);	WAIT RETURN: SÉRVICE=INYF;	INVOKE: SERVICE= INTFS, PARAMETER LYST=(FINDR,NEXT,UIST,1001);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,IDUI);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(MODIF, DAY, UIST, DMY);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INIFS, PARAMETER LIST=(MODIF,DMY,UIST,DMY);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR, DMY, AUDI, CALC);	WAIT RETURN: SERVICE-INTF;	INVOKE: SERVICE: INTES, PARAMETER LISTE (MODIF, DMY);

STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD JMAGE	Ευ	INPUT ALT INPUT	PEPEEP KET-1.C	
MAIT KETURN: SERVICE=INTF;	000007710	121	567	
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR, DMY, ADUN, CALC);	00007220	77.2	750	
WALT RETURN: SERVICE=INTF;	00007230	نكد	101	1
INVOKE: SERVICE= 147F5, PARAMETER LIST=(MODIF, DMY, ADUN, DMY);	00001240	124	702	
WAIT RETURN: SERVICE=INTF;	00007250	725	205	
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR,NEXT,UIST,IDUI);	00007260	726	704	
WAIT RETURN: SERVICE=INTF;	07570000	tět	705	
C WORMAL RETURN	00007280	724	700	FRU
96 YERCISI = 0	00001290	621	191	M OO
GO TO 150	0001300	VEL	76.6	T X
C COMPATABILITY ERKOR	01670000	731	901	1.
100 CONTINUE	00007320	732	310	1
YLOC(5) = 1	00007330	753	111	1
150 CONTINUE	00007340	734	317	000
RECEASE: FACILITY=DEPTR;	00007350	725	n3	DC
END: ENDO SERVICE;	00007360	736	711	-
	00000330	737		-
	03670000	dil		
	95843399	614		
ENDO SERVICE: 10=ADDSL, NAME=ADDSL\$, SAVE AREA SIZE=10,	00763903	740	32	
PAKAMETER LIST=(PERRY, SFEDI);	00001410	741	116	
CCMMON /WK/ TLOC	02420000	int.	7:7	
INTEGER SUNFHI, SEF U., SEEDI, TLOCIIA)	06 74 30	143	1118	
INTEGER OBTAN FILIDE HODIE STIPS CONCT. DISCN. FRAS	0772000	7.4	314	

DATE 99/10/77			
INMIT CARD THAGE20204040506070	80	SEC-NO LEV SEC-NO	RESOLF REF-MO
INTEGER UMY, NEXT, PRIUR, OWNER	06007450	745	071
INTEGER ADDI, ADUN, ALDC, GRADE, TOEN, THST, HOSC, POSN, UIST, UPINA	09523303	746	721
INTEGER CALC. GRUI, IDUI, INUM, PHID, POUI, SHID, UMAM, UMGR, UHPOS, UHUI	00001410	747	766 237
EGUIVALENCE (\$SAVE(1),PERR), (\$SAVE(2),SEED)	60007480	748	72.55
REAL SUNFM, RN, PERR, PERRI	06067496	740	124
CATA DBTAN,FINDR,MUDIF,STOR,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/	00001500	75.0	72.5
DATA DHY,NEXT,PRIOK,OWNER/0,1,2,3/	01510000	151	72.6
DATA ADDI, ADUN, ALOC, GRADE, TOFN, INST, HOSC, POSN, UIST, UNHA	00007526	752	127
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/	00007536	753	972
DATA CALC, GRUI, IDUI, INUM, PHID, PRUI, SMID, UMAM, UMGR, UMPOS, UMUI	00007540	7:4	521
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/	00001550	755	730
END: DECLARATIONS;	09512000	756	111
PERR = PERRI	07270000	151	732
SEED = SEEDI	000075.60	759	261
END: INITIALIZATION;	06670000	156	134
SEIZE: FACILITY=ADDSL;	00001600	760	135
3	000001010	76!	736
C RETURN CODE -TLUCISI- = 0 IF NO ERROR, I IF ERROR	00007620	767	131
2	00007630	763	738
C PARAMETER LIST CONTAINS-	00001640	764.	73.5
C II DHL TYPE - OGTAIN=I.FIND=2.HODIFY=3.STORE=4.CONNFCT=5.	000007050	76.5	071.
C DISCUNGECT=6,FPAKE=7	000076.60	76.5	74.1
C ZJ QUALIFIER - DAY=0.NEXT=1.PRIOR=2.OWNER=3	06647676	767	241
C ZI KECOKO TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE OML IS	00007680	745	14:

No. of Lot

33						20	HIS	la co	· · · · · · · · · · · · · · · · · · ·	100	15		4/1				4.00							00u 641	•
r 26F	P.FMLFR KET-P.C	746	74.5	746	727	748	74.7	750	75.1	752	75.5	10.	532	756	157	128	13.c d	760	76.1	762	763	164	765	74.6	167
	ALT INPUT														THE RESIDENCE OF COMPANY AND PARTY OF THE PA						A STATE OF THE PERSON OF THE P	The state of state of the state			
	INPUT SFQ-MO	249	770	177	CIT	773	774	311	776	lie	366	611	740	164	CoL	7.63	746	285	ANE	787	71.8	hil	206	101	707
LATOR	0360	00001590	0001100	01110000	021720	00001730	0711000	05110000	00007760	0177700	09660000	06770000	000001000	000007110	00001820	00007830	00007640	000007150	09420303	00001870	00007560	06920000	000.40000	05007510	00001920
INFORMATION PROCESSING SYSTEM SIMULATOR STANDARD INFUT STREAM LISTING DATE 09/10/77	INPUT CARD IMAGE 303040506070	TO OPERATE	4) SET IDENTIFIER - INDEX		INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, IDEN, CALC);	WAIT RETURN: SERVICE-INTF;	INVOKE: SERVICE= INIFS, PARAMETER LIST=(OBTAN, DMY, UNNA, CALC);	WAIT RETURN: SERVICE=INTF;	INVOKE: SERVICE= INTES, PARAMETER LIST=(OBTAN, DMY, MOSC, CALC);	WAIT RETURN: SERVICE=INTF;	I = \$UNFMI(I,2,5EED)	IFII .6T. 1) GO TO 100	INVOKE: SERVICE= INTES, PARAMFTFR LYST=(OBTAN, DMY, MOSC, CALC);	WAIT RETURN: SERVICE=INTF;	TOO COPTINUE	INVOKE: SERVICE= INTES, PARAMETER LISTER; DHY, IDEN, DHY);	WAIT RETURN: SERVICE=INTF;	INVUKE: SERVICE= INTFS. PARAMETER LYST=(STOR, DMY, ADDI, DMY);	WAIT RETURN: SERVICE=INTF;	DETERMINE IF COMPATABILITY CHECK HAY OCCURRED	RN = SIMFHIG.O. 1.0. SEED)	IFIRM .LT. PERRY GO TO 300	NO CCHPATABILITY ERROR	INVOKE: SERVICE= INTES, PARAMETER LYST=(STOR, GHY, UIST, DHY):	WAIT RETURN: SFRVICE=INTF;

DATE 09/10/77				
INPUT CARD TMAGE 3030	80 81	SEC-IN LEV SEU-10	NEF-NO	
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR,NEXT,GRADE,UMGK);	00007930	743	71.5	
WAIT RETURN: SERVICE=INTF;	0962000	761	346	
INVOKE: SERVICE= INTFS, PARAMETER LISY=(CONCT. DMY, UIST, GRUI);	02520000	364	170	
WAIT RETURN: SERVICE=INTF;	00007500	904	101	
INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDK, OMY, MOSC, CALC);	01.510000	404	711	THI:
WAIT RETURN: SERVICE=INTF;	00001980	794	:11	S PA
INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT.DMY, TOEN, PMID);	0644,0000	464	174	GE I
WAIT RETURN: SERVICE=INTF;	0000000	0v3	321	SBI
I = SUNFAI(1,2, SEED)	00000000	103	776	est
1FIT .67. 1) 60 TJ 200	00008020	PAS	1.66	OUA ?
INVOKE: SERVICE= INTES, PARAMETER LIST=(FINDR, DMY, MGSC, CALC);	0000000	ė U s	17.6	10 D
WAIT RETURN: SERVICE=INTF;	00006040	£07.	-11	oc Ph
INVOKE: SERVICE= INTES, PARAMETER LIST=(CONCT, DMY, IDEN, SMID);	05080000	\$0¢	296	
WAIT RETURN: SERVICE=INTF;	000008060	Pne	181	-
ZOO CUNTINUE	00000000	107	195	
NORHAL RETURN	0000000	нун	16:	
1LUC(6) = 0	0608000	6.0	192	
GU TU 350	0018000	ماد	185	
COMPATABILITY ERROR	000000110	kii	186	
300 CCMTINUE	021-0000	275	7117	
TLGC(E) = 1	00006130	Fis	188	
350 CINTINUE	00008140	7.0	169	
RELFASE: FACILITY=ADDSL;	00000150	- 419	23.6	
END: ENDO SERVICE:	00008160	A18	151	

	¥			1 AGE
STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD IMAGE10202030404050606070	09****	THPUT AL	ALT INPUT	MEMBER AFF-ND
	000ce170	7.4		
	00000	818	-	
ENDO SERVICE: ID=INTF, NAME= INTFS, SAVE AREA SIZE*10,	000006190	618		761
PARAMETER LYST = (P! .P2.P3.P4);	00780000	926		641
C IDMS INTERFACE ROUTINE	01280000	164		154
COMMON /WK/ TLOC	00008220	622		3.5
INTEGER ERRST, CIND, VALERR	00008230	423		79.6
INTEGER DAL, OPER, RT, SET, PI, P2, P3, P4	000006240	F24		161
INTEGER SUNFHI, SEED, SEEDI, TLOC(170)	00006250	964		758
INTEGER UBTAN FINDK , MUDIF, STOR , CONCT , DISCN, ERAS	09793000	459		552
INTEGER DMY, NEXT, PRIOR, OWNER	00006270	127		603
INTEGER ADDI, ADUN, ALOC, GRADE, IOEM, INET, MOSC, POSN, UIST, UNMA	00:308280	Pos		. BC !
INTEGER CALC, GRUI, IDUI, INUM, PHIO, POUI, SHID, UMAM, UMGR, UMPOS, UPUI	04280000	P29		102
EGUIVALENCE (SSAVE(5), DML),	00006300	830		£03
1 (\$\$AVE(6), OPER), (\$\$AVE(7), RT), (\$\$AVE(8), SET)	00008310	155		*00
REAL SUNFH, RN	00008320	ć.o		105
DATA CBTAN,FINDR, HODIF, STOR, CORCT, DISCN, ERAS/1,2,3,4,5,0,7/	000008330	R33	-	80£
DATA CMY,NEXT,PRIDR,OWNER/0,1,2,3/	000000	454		£67
DATA ADDI, ADUN, ALUC, GRADE, IDEN, INST, MNSC, POSN, ULST, UNMA	00006253	pse		404
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/	03580000	364		2000
DATA CALC, GRUI, IDUI, INUM, PMID, PMUI, SMID, UMAM, UMCR, UMPUS, UMUI	0168333	467		F10
1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/	000008380	ARA		2:12
END: DECLARATIONS;	0600000	926		513
Phi - B	000000	078		

DATE = 10 DATE	NATION N	STANDARD INPUT STREAM LISTING	-		
OPTR = P2	OPPR = P2	1:			
NT = P3	SET = P4	INPUT CARD IMAGE	90		KEMEEN KEF-NC
SET = P4	### ### ##############################	OPFR = P2	00008410	. 54 t	614
SET = P4 GRONG#20 P43 Lib	SET = P4 00008430		00006420	276	
STEEF FACILITY=INTE;	SETZET FACILITY=LIATION; SETZET FACILITY=LIATE SETZET FACILITY=LIATE STEP 1. PERFORM FUNCTION VALIDITY CHECKING AND SEGNENTATION 00006450 644 124 ISYSTEM HOOK! CAOOGE-50 644 124 STEP 2. TEST ERROR STATUS AFER VALIDITY CHECKING CAOOGE-50 644 124 STEP 2. TEST ERROR STATUS AFER VALIDITY CHECKING CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM ANY BEFORE PROCEDURES CAOOGE-50 644 124 STEP 3. PERFORM STATUS FROM THE FRENCE PROCEDURES CAOOGE-50 644 124 STEP 4. TEST THE ERROR STATUS FROM THE FRENCE PROCEDURES CAOOGE-50 645 124 STEP 4. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CANCEL INDICATOR CAOOGE-50 645 124 STEP 5. TEST THE CAOOGE-50 645 124 STEP 5. TEST THE CAOOG		00008430	673	636
STEEF FACILITY STREET FACILITY CHECKING AND SEGREFUTATION 00006450 644 619 620 620 620 620 620 620 620 620 620 620	STEEF FACILITY=INTF; SCHEMTATION STEPFACE STATUS AFTER VALIDITY CHECKING GOOGB450 S44 E.20	END: INITIALIZATION;	07790000	778	617
STEP 1. FERFURN FUNCTION VALIDITY CHECKING AND SEGNENTATION	ISYSTEM HOOK) UNLERR = 0	SETZE: FACTLYTY=INTF;	00008450	845	616
1595FEH HGOK1	SYSTEM HOOK VALERR = 0	STEP 1. PERFURH FUNCTION VALIDIT	00008460	846	61.9
STEP 2. TEST ERROR STATUS AFTER VALIDITY CHECKING G0008490 F49 F22 IFIVALERR .AE. 0 60 TV 80 G0008500 F50 F23 INU VALIDITY ERROR STATUS AFTER VALIDITY CHECKING G0008510 F51 F24 STEP 3. PERFCKH ANY 'BEFORE' PROCEDURES C0008510 F71 F24 STEP 4. PERFCKH ANY 'BEFORE' PROCEDURES C0008520 F72 F25 ISYSTEM HOOK CIND 1S THE 'CANCEL' TROICATOR C0008520 F72 F25 ERRST = 0	STEP 2. TEST ERROR STATUS AFTER VALIDITY CHECKING C0008490 P49 B22 IFTVALERRE. 01 GO TO 60 E23 INVALERRE. 01 GO TO 60 E23 INVALERRE. 01 GO TO 60 E23 INVALERRE. 01 GO TO 60 E23 STEP 3. PERFOREE. 01 GO TO 60 E23 ISTEP 3. PERFOREE. 01 GO TO 60 E23 ISTEP 3. PERFORE STATUS INDICATOR C0008520 E54 E27 CIND IS THE "ERROR STATUS INDICATOR C0008520 E55 E25 CIND IS THE "CANCEL" INDICATOR C0008520 E55 E25 ISTER 4. TEST THE ERRUR STATUS FROM THE "HEFORE" PROCEDURES C0008520 E56 ISTER STE. 01 GO TO 80 C0008520 E60 E25 IFTERRSTE. 01 GO TO 80 E60 E25 E2		000008410	178	6.20
STEP 2. TEST ERROR STATUS AFTER VALIDITY CHECKING C00004590 P29 E23 IFIVALERR .NE. 01 GO TO 80 C00004590 P29 E23 NO VALIDITY ERRORS GCURRED C00004500 PF1 P24 STEP 3. PERFORM ANY "BEFORE" PROCEDURES C00004520 PF2 P25 ISTEP 3. PERFORM ANY "BEFORE" PROCEDURES C00004520 PF2 P24 CIND	IFTVALERR .h.e. 01 GO TO 80 823	VALERR = 0	09780000	673	321
IFTVALERR .NE. 01 GO TO BO IFTVALERR .NE. 01 GO TO BO NOTELIDITY ERRORS OCCURRED NOTOCESSO RFS L24 STEP 3. PERFURN ANY 'BEFORE ' PROCEDURES NOTOCESSO RFS L27 CIND 15 THE "ERROR STATUS" INDICATOR NOTOCESSO RFS L27 CIND 15 THE "CANCEL" INDICATOR NOTOCESSO RFS R28 ERRSY = 0	IFTUALERR *NE. O J GO TO BO	TEST ERROR STATUS AFTER	66480000	678	
NOT VALIDITY ERRORS GCCURPED	NO VALIDITY ERRORS OCCURRED	TETVALERR .NE. 0) GO TO 60	0000 200	ASO	
SYSTEM HOOK ISYSTEM HOOK NOTE - EKAST IS THE "ERROR STATUS" INDICATOR 00000520 054 026	SYSTEM HOOK 15YSTEM HOOK 15YSTEM HOOK NOTE = ENAST IS THE "ERROR STATUS" INDICATOR		00000010	[ted	1
15Y3TEH HOOK) NOTE = ERAST IS THE "ERROR STATUS" INDICATOR OCCOGSSO FS	(SYSTEM HOOK)	STEP 3. PERFORM ANY "BEFORE" PR	02380000	8 . 5	
NOTE - EKAST IS THE "ERROR STATUS" INDICATOR	NOTE		06666530	Fr3	1
STEP 4. TEST THE ERROR STATUE FROM THE "HEFORE" PROCEDURES 0000085E0 855 828 828	ERRST = 0 CIND 1S THE "CANCEL" TNDICATOR 006/08/550 A55 A2H CIND = 0 CIND = 0 00008/500 PF6 L25 STEP 4. TEST THE ERROR STATUS FROM THE "MEFORE" PROCEDURES 00008/500 PF6 B31 ISYSTEM HOOK! IFICERRST .NE. 01 GO TO 80 PROFEDURES 00008/500 P60 B53 NO ERRORS OCCURRED P60 P60 P50 NO ERRORS OCCURRED P60 P60 P60 STEP 5. TEST THE CANCEL INDICATOR P34 P34		04580000	7 30	
STEP 4. TEST THE ERROR STATUS FROM THE "HEFORE" PROCEDURES	CIND = 0	CIND IS THE "CANCEL" T	00008550	455	1
STEP 4. TEST THE ERROR STATUS FROM THE "MEFORE" PROCEDURES 000008570 RET 530 ISYSTEM HOOK! 150 TO 80 REP	CIND = 0 000008570 PROTEDURES 000008570 PAGE PAGE STEP 4. TEST THE ERROR STATUS FROM THE "MEFORE" PROCEDURES 000008570 PAGE B31 IF(ERRST .NE. 0) GO TO 80 NO ERRORS OCCURRED NO ERRORS OCCURRED FE1 F34 STEP 5. TEST THE CANCEL INDICATOR	ERRST = 0	000008560	9.4	
STEP 4. TEST THE ERROR STATUS FROM THE "MEFORE" PROCEDURES	STEP 4. TEST THE ERROR STATUS FROM THE "MEFORE" PROCEDURES	CIND = 0	07380000	647	
IFTERRST -NE. 01 GO TO 80 E32 IFTERRST -NE. 01 GO TO 80 B33- NO ERRORS OCCURRED FE1 F34 STEP 5. TEST THE CANCEL INDICATOR CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC	ISYSTEM HOOK)	STEP 4. YEST THE ERROR STATUS FROM THE	30006580	999	1
IFTERRST -NE. 0) GO TO 80	IFTERRST .NE. 01 GO TO 80 NO ERRORS OCCURRED STEP 5. TEST THE CANCEL INDICATOR 635		00008540	pro	
NO ERRORS OCCURRED FE1 STEP 5. TEST THE CANCEL INDICATOR 652 IF(CIND .NE. 0) GO TO 100 663 DML MOT CANCELLED 60006140 614	NO ERRORS OCCURRED STEP 5. TEST THE CANCEL INDICATOR	IF(ERRST .NE. 0) GO TO 80	00008000	966	1
STEP 5. TEST THE CANCEL INDICATOR 1FICIND .NE. 0) GO TO 100 DML MOT CANCELLED COOUGRAGO F1.4	STEP 5. TEST THE CANCEL INDICATOR		000008¢ 10	FEI	134
TFICING .NE. 0) GO TO 100 Erg			00005620	F52	685
DML NOT CANCELLED	CCOURAL 30 Fr 3	1F(CIND .NE. 0) GO TO 100	06000000	Fr3	136
	DML MOT CANCELLED		00008146	715	153

STANDAKD INPUT STREAM LISTING			1011
DATE 09/10/77			
INPUT CARD IMAGE	F0	JNP"T ALT ALPUD	McPf.17 F.F-146
PERFURM DATA BASE MANAGEMENT FUNCTION	00008650	865	36.3
ALL DAL'S ARE ASSUMED TO REQUIRE DRMS SERVICE	0007.8660	846	900
ISYSTEM HOOK FOR FRONT-END PACK-END COMMUNICATION)	00008670	PA7	0 + 2
INVOKE: SERVICE=CAMPS, PARAMETER LIST=(DML+UPER,RT.SET);	000008060	FER	.74
WAIT GETURN: SERVICE=CAMP:	0000086.00	460	545
STEP 7. TEST ERROR STATUS	00190000	670	143
TFTEKRST .EQ. 0) GU TO 100	00000110	169	5,3
ERROR OCCURRED IN THE EXFCUTION OF THE OML	00008720	ذذه	649
STEP 6. EXECUTE ANY "ON ERROR NURING" PROCEDURES	00008750	573	546
TSYSTEM HODK)	00006740	77.3	۲٦:
	000000150	P75.	646
STEP 9. EXECUTE ANY "AFTER EPRINR" PRICEURES	09290000	474	0,40
(SYSTEM HOOK)	00035770	£_0	04.
STEP 10. PERFORM SEGMENTATION	00008760	474	139
(SYSTEM HOUK)	000000	76.5	7;,,
	000008000	ben	653
RETURN CONTROL TO THE APPLYCATION PROGRAM	CCC36419	444	767
RELEASE: FACILITY=INTF;	COGOREZE	P#7	4.
END: FNOU SERVICE;	00008130	C) 00.	656
The same of the sa	00006840	744	
	OCOURE 50	870	
EMNO SERVICE: ID=CAMP, NAME= CAMPS, SAVE APLA SIZE=10,	COOCEFEC	FSA	553
PARAMETER LIST = (PI.PZ.P3.P4);	00000000	pa.1	E ''
SECTION SECTIONS AND SECTION SECTIONS SECTIONS	31.10000	0 44	• • • •

DATE 09/10/27				
1020304050606070.	98	3 C-1 C	LEV SFC-NC	RET-NU
THIS PROCRAM "THREADS REQUESTS AS (THE) DBMS CAN ACCOMBUATE	06008000	234		be f
THEM - HOLDING THE NECESSARY USFRS IN THE WAIT STATE!	00680000	600		101
	01030000	164		1.02
COMMON YWK/ TLGC	00006920	692		563
INTECER DML.OPER.RT.SET.P1.P2.P3.P4.YLCC(10)	00000430	£58		*92
EQUIVALENCE (\$SAVE(1), DAL), (\$SAVE(2), OPER), (\$SAVE(3),RT),	04686303	758		5)4
1\$5AVE(4),SET)	05580000	163		
END: DECLARATIONS:	09690000	958		199
DMC = PI	02430000	143		304
OPER = P2	00000000	403		BEST
RT = P3	06583700	604		343
SET = P4	00060000	005		87.
END: INITIALIZATION;	01060000	105		213
SEIZE: FACILITY*CAMP;	04009620	2vb		613
QUEUE: FACILITY=DBMS;	06050000	463		763
WAIT FACILITY: FACILITY=DSMS, STATUS=0, CONDITION=EQ;	05009040	504		27.5
DEPART QUEUE: FACILITY=DFMS;	05060000	2 US		1.76
SET STATUS: FACILITY=DRMS, STATUS=1;	00000000	306		11.1
INJOKE: SERVICE=DUMSS, PARAMETER LIST=(DML,CPER,RT,SET);	01060000	1.00		673
WAIT RETURN: SERVICE=DBMS;	09060000	åvo		875
SET STATUS: FACILITY=DBMS, STATUS=0;	04040000	, our		nt 0
RELEASE: FACILITY=CAMP;	00007100	016		£6!
END: FNDS SERVICE;	00000	1115		244
	00009120	516		

- Contractor

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No. of Lot

1000000

STAMBARD INPUT	STREAM LISTING		-		
DATE	11/01/60				
INPUT CARD IMAGE	02	90	TNPUT SFR-ND	ALT INPUT	MEMET'N REFER
	93	06140000	6.5		
ENDS SERVICE: 10=08MS, NAME= DRMC4, SAVE AREA SIZE=10,		060609140	716		£ 6.3
PARAMETER LIST = (P1,P2,P3,P4);	00	00000150	5.6		#81
C ****** DATA BASE MANAGEMENT SYSTEM FUNCTIONS******		02150000	416		1.96
COHHON JAKZ TLOC	00	07190000	Lio		#84
COMMON /PE/ PAGET	00	09150000	615		179
COMMON /PAGAR/ FSTPN.LSTPN.AVFOIS.ANSAV.KDREC.BUF.RNRET.INDX.NVAG 00059190	SUF, RNRET, INDX, NFAG OO	051600	619		3.3.1
INTEGER FSTON, LSTON, AVEDIS, RNSAV, RDREC, PUF (10,3)		00760000	0.5		184
INTEGER DML, DPER, KT, SET, PI, P2, P3, P4, TLOC(110), INDX		012-0000	155		944
INTEGER 1, J. RN. SUNFMI, KETP, TOT!, TOT?, PAGFT(101	00	02260000	225		147
EQUIVALENCE (185AVE(5), DML), (185AVF(6), OPER), (\$5AVE(7), KT),		00008230	623	1	268
(\$\$AVE (\the) , \$ET)	00	00069240	7.0		€6#
C BUFFER - BUFILD - FIRST SURCRIPT IS BUFFER NUMBER		000005250	364		46.0
C - SECOND SUBSCRIPT	00	00005260	420		467
C (1) PAGE NUMBER (0 TO NUMBER UF PAGES IN DB)		00000000	1:6	The state of the s	340
C (2) LEAST RECENTLY USED (LRU) INDICATOR (0-NPAG)		30C69280	600		263
C LRU PAGE HAS VALUF O, NEWEST PAGE	= NPAG	06760000	020		954
(3) MUST REWRITE INDICATUR (0-1)		00004300	01.5		55.2
C REWRITE PAGE WHEN = 1	00	01660000	100		000
C PAGE FAULT STATISTICS AREA - PAGET(10) -	00	00009320	633		451
C (1) MIMBER OF FACES CONCURRENTLY IN THE PAGE BUFFER	The state of the s	C050933C	ich		405
C (2) PAGE SIZE (BVTES)	00	00000340	760		503
C (3) PAGE WRITES	00	0.260000			12,
C (4) CACHE MISSES ON WRITE	00	000003:40	100		:05

77 J974	The same of the sa		MEPUFF FIF-1.0	405	20%	50e	515	11 m	RUM	PAGI OGE	2 15	RES	T Q	UALITU	עם ניין	315	515	075	125	??\$	4:3	777	5.5	7.20	124	325	576	
	And the second s		TEPUT ALT TEPUT	937	8 i.o	615	076	641	دلان	670	444	575	944	£75	675	075	CKA	0K V	596	9.3	756	OKK	150	Cer	g i j	CEO	046	
SYSTEM SIMULATOR	FAM LISTING	77/01/77	.60	00009370	09630000	SED 00009350	00004400	00000430	00000450	00000430	07760000	0575000	09460900	01.760000	04760000	04760000	000055000	01460000	02540000	000003630	THE 16H370/16500009540	04450000	39560000	01.260000	085,0000	04340000	0000000	60Pr
INFORMATION PROCESSING SYSTEM SIMULATOR	STANDARD INPUT STREAM LISTING	DATE 09	INPUT CARD IMAGE202030304040506070	C (5) PAGE READS	C (6) CACHE MISSES ON READ	C (7) NUMBER OF PAGES REWRITEN RFFORF PAGE SPACE REUSED	C (8) NIMBER OF PAGE REPLACEMENTS WITHOUT KENKITTING	C (9) MIMBER OF PAGE NUMBERS GEPERATED (3Y GENER)	C (10) TOTAL DISTANCE BETWEEN PAGE REFERENCES	END: DECLARAYTONS;	DML = P1	OPER = P2	RY = P3	SET = P4	IF(PAGFT(2) .EG. 1024) GO TO 18	DO 19 I = 1 ₉ 10	bo 16 J = 1,3	10 BUF(1,3) = 0	00 15 1 = 1,10	15 PAGFT(1) = 0	C PAGE SIZE IS 1024 BYTES ON THE POPITYTO AND 3156 ON THE IEM370/16500009540	PAGFT(2) = 1024	FSTPH = 1000	LSTPN = 1200	AVEDIS = 45	RNRFT = 1000	HPAG = 5	1020304050

7									F.	Rout		-1	3 -		E IX	מע	THE	S P.	AGE	151	SES!	UQU	ALI TO	TY P	TAC	
		MENL LA Kzi-Nr	050	431	754	433	434	454	936	153	436	454	075	176	346	545	77	645	7.	2.5	73,	345	055	isi	104	.53
		ALT INFUT											The second second second second													
		TP-PLIT SFC-ND	1961	296	676	496	.95	946	647	940	696	315	ies	245	673	51.5	545	4.6	665	960	616	080	163	646	CRP	700
MULATOR G		.7060	01960000	02960000	06960000	000006640	00000460	09960000	0608070	00000660	04960000	001.50000	01160000	02260000	00009730	07160000	05260000	0260300	021,40000	00000780	05260900	00000000	00034010	00006450	06860000	000008640
STANDARD INPUT STREAM LISTING	0ATE 09/10/77	INPUT CARD IMAGE2020304040505070	PAGFT(1) = NPAG	INDX = 0	18 CONTINUE	END: INITIALIZATION;	SETZE: FACILITY=DBMS;	C INVOKE THE REQUESTED DML SFRVICE	GO TO 1160,200,300,400,500,600,700),DML	C DBTAIN DML	IDD CONTINUE	INVOKE: SERVICE=OBTAIS, PARAMETER LIST=(DML,OVER,RT,SET);	WAIT RETURN: SERVICE=OBTAIN:	60 10 960	C FIND DML	200 CONTINUE	INVUKE: SERVICE=FINDS, PARAMETER LIST=(DML,OPER,RT,SET);	WAIT RETURII: SERVICE=FIND;	60 10 960	C MODIFY DML	300 CONTINUE	INVOKE: SERVICE=MODIFS, PARAMETER LIST=(DML, UPER, RT, SET);	WAIT RETURN: SERVICE=MODIFY;	60 13 900	C STOKE DHL	400 CONTINUE

STANDARD INPUT SIREAM LISTING					
, DATE 09/10/77					
.010970.	F.	14-DUT	ALT INFOI	MENTEL	
PARAMETER LIST = (P1,P2,P3,P4);	04001000	1008		5.16	
C execest DBMS OBTAIN DML PROCESSING FUNCTIONS+8****	00101000	0101		11.5	
COMMON /WK/ PLGC	61101330	1001		576	
COMMON 7PE7 PAGET	30010120	1012		46.	
COMMUN /PAGAR/ FSTPN, LSTPN, AVENIS, RNSAV, POREC, BUF, RNRET, INDX, NPAG GGG;3:30	06:61000	1013		1385	
INTEGER FSTPN.LSTPN.AVEDIS.RNSAV.RDRFC.RUF(10.3)	0710100	1014		461	
INTEGER DAL, OPER, RT, SET, PI, P2, P3, P4, TLOC(10), 1NDX	00010120	1015		795	-7
THTEGER I.J.RN.SUNFHI.RETP.TOTI,TOT2,PAGFT(10)	00121200	1016		534	
EGUIVALENCE (8SAVE(1), OHL), (\$SAVE(2), OPER), (\$SAVE(3),RT),	000,0170	1017		264	
I (SSAVE(4),SET)	00010100	10.6		586	P
END: DECLARATIONS;	05.01000	7101		316	
UML = PI	00010500	1020		735	4 3E
UPER = P2	00010210	1201		404	Mar.
RT = P3	000 10220	1022		284	
SEF = P4	06 30 1000	1023		955	140
END: INITIALIZATION;	07.01000	1024		165	101
SELLE: FACILITY=UBTAIN;	00010250	inze		266	AD C
2	0001076	1076	And the second s	5.54	
C GENERATE A RANDOM PAGE NUMBER 1001-1200 - RETURNED VALUE IS KNKET	27:01000	Livi		755	
20 CALL GENRP	00010580	1076		344	
C GET THIS PAGE IF NOT ALREADY PRESENT IN THE CACHE	000:0500	١٧١٠		344	
INVOKE: SCRVICE=PAGMRs, PARAMETER LIST=[1];	0001000	0E0 i		List	
MAIT RETURM: SERVICE=PAGMR:	00010310	1031		***	
C DETERMINE IF THE NEXT PAGE IN THE CALC SET NEEDS TO BE ACCESSED	00010320	1032		250	

DATE 09/10/77					
INPUT CAPD THAGE 30305050	7080	Or ti-vio	ALT TAFUT LFV SEC-NO	PEMEFA AFF-NO	
C THIS IS THE CASE ON "FIND" AND/OR "ACQUIRE" FOR CALC SETS	06601000	E É U		1360	
C WHEN THE RECURD COULD NOT BE PLACED ON THE PAGE DETERMINED BY	00010340	1636		1001	
C THE CALC ALGORITHM	0601000	1035		3001	120
RN = \$UNFHI(1,1000,2)	00010360	1036		1063	Jan 3
JFIRM .67. 900) GU TO 20	07501000	1001		1004	1432 24-Y
C NEXT PAGE IN THE CALC SET NEED NOT BE ACCESSED	000:0360	1030		1505	10.
C PROCESSING OF "OBTAIN" FINISHED	0601030	1039		1006	CALL
RELEASE: FACILITY=OBTAIN;	00010400	U701		1001	TQL
END: ENDO SERVICE;	00010410	1001		166c	10
	02491000	2701			шос
	00010430	8738			
ENDO SERVICE: 1D=FIND, NAME= FINDS, SAVE AREA SIZE=10,	00010440	7		1009	
PARAMETER LIST = (PI,P2,P3,P4);	00010450	5701		1010	
C ******* DBMS DRTAIN DML PROCESSING FUNCTIONS******	00010460	1046		7107	
COMMON ZWK/ TLOC	04.501000	1041		1012	
COMMON 7PE/ PAGET	00010480	875 i		1015	
COMMON /PAGAR/ FSTPN.LSTPN.AVFDIS.RNSAV.RDREC.BUF.RNRET.1MDX.NPAG COOTO450	05401000	6701		7101	
INTEGER FSTPN, LSTFN, AVEDIS, RNSAV, RDREC, BUF(10,3)	000:000	1050		1015	
INTEGER DML, OPER, RT, SET, P1, P2, P3, P4, TLOC(10), INDX	00010010	101		:016	
INTEGER I, J. RN. BUNFMI, RETP, TOTI, TOTZ, PAGFT(10)	00010520	2 su i		7:01	
EQUIVALENCE (18SAVE(1), DML). (4SAVE(2), OPER), (18SAVE(3),RT).	0001000	1053		1015	
I (\$SAVE(4),\$ET)	00010540	1054		1015	
END: DECLARATIONS:	00010550	1.055		7620	
- DE	00010560	1046		1051	

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STANDARD INPUT STREAM LISTING					
74/0/17 04/10/77					
INPUT CARD IMAGE203040506070	60	It-bill SEC-NO	ALT IMPUT	MFMPFK F.F-NC	
OPER = P2	0001000	1067		1012	
RT = P3	0001000	1058		1623	
SET = P4	06501000	1056		1024	
END: INITIALIZATION;	0001000	0+0		15.5	
SEIZE: FACILITY=FIND;	00010610	1061		11.26	
C GENERATE A RANDOM PAGE NUMBER 1001-1200	0201000	1062		1627	
20 CALL GENRP	000,000	1063		1026	
C GET THIS PAGE IF NOT ALREADY PRESENT IN THE CACHE	0001000	1064		5201	
INVOKE: SERVICE=PAGHRS, PARAMETER LIST=(1);	00010020	3701	And the second control of the second control	1650	
WAIT RETURN: SERVICE=PAGMR:	0001000	10.66		1001	
C DETERMINE IF THE NEXT PAGE IN THE CALC SET NEEDS TO BE ACCESSED	0001000	1067		7601	
C THIS IS THE CASE ON "FIND" AMDZOR "OBTAIN" FOR CALC SETS	000100k0	1069		1033	
C WHEN THE RECORD COULD NOT BY PLACED ON THE PAGE DETERMINED BY	05901303	590;		1034	
C THE CALC ALGORITHM	000010000	المان		103:	
RN = \$UNFHI[1,1000,3]	0001000	·¿vì		1636	
TFTRN .GT. 9001 GD TO 20	00010100	1072		1037	
C NEXT PAGE IN THE CALC SET NEED NOT BE ACCESSED	06701000	1073		1629	1
C PROCESSING OF "FINISHED	(201003	4571		1234	
SAD CONTINUE	06616750	1016		1040	
RELEASE: FACILITY=FIND;	00010160	3000		1641	
END: ENDO SERVICE;	00101000	101		1643	
	0410100	3201			
	24601000	9671			
ENDO SERVICE: 10=MODIFY. NAMF = MODIFS. SAVE AREA SIZE=10.	0010100	1080		104.	

PAGE 46			INPIT ALT INPIT MEMICK SEC-OF LEV SEC-AL REF-AL	1,081			1083	1044 1044 国社	1085 1045 1045	1086 X	1087	1088 165.	1045	1000 1053	100C 100C	1092	1003	1653	1005	1005	1097	10c1	1009 ; 005	1160	1362	1102 100:	1103	The second secon
SIMULATOR	Inc		7060	01001000	00010630	00010850	00010830	00010040	X.NPAG 00010650	00010860	0001010	00010880	06901000	20501000	01401000	07601000	60010530	00010040	05501002	00010960	000:000	00010580	04601000	00011000	00011010	00011020	00011030	50011045
INFORMATION PROCESSING SYSTEM SIMULATOR	STANDARD INPUT STREAM LISTING	DATE 09/10/17	INPUT CARD IMAGE1010203030405050	PARAMETER IST = (P].P2.P3.P4);	u	C ******* DBMS BCDIFT DML PRICESSING FUNCTIONS******	COMPON ZWKZ TLOC	COMMON 7PE7 PAGFT	COMMEN /PAGAR/ FSTPN, LSTPN, AVEDTS, RNSAV, ROREC, BUF, RNRET, INDX, NPAG 00010650	INTEGER FSTPN, LSTPN, AVEDIS, RNSAV, RDREC, BUF (10,3)	INTEGER OML, OPER, RT, SET, PI, P2, P3, P4, TLOC(10), INOX	INTEGER I,J,RN,SUNFMI,RETP,TOTI,TOTZ,PAGFT(10)	EQUIVALENCE (\$SAVE(1), DML), (\$SAVE(2), OPER), (\$SAVE(3), RT),	1 (\$\$AVE(41,\$ET)	END: DECLARATIONS;	DM[= p]	OPER = P2	AT = P3	SET = P4	END: INITIALIZATION;	SEIZE: FACILITY=MODIFY;	C WALGREEN DATA - 50% READ & 100% WRITE MISSES	C (SYSTEM HOOK)	PROCESS: TIME=14.7;	RELEASE: FACILITY=MODIFY;	END: ENDO SERVICE:		FNDG SERVICE: ID=STORE, NAME STORES, SAVE AREA SIZE=10,

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		The second secon	
STANDARD INPUT STREAM LISTING			
DATE 09/10/77			
INPUT CARD IMAGE20304040506070	90	SEC-NO LEV SEC-NO	MEMILIA K.T-N
PARAMETER LIST = (P1,P2,P3,P4);	00011050	1105	10¢7
C ******* DBMS STORE DML PROCESSING FUNCTIONS*****	00011000	1106	1066
COMMON /WK/ TLDC	00011000	1107	1965
COMMON /PE/ PAGET	000:1080	1108	3030
COMMON / FACAR / FSTPN. LSTPN. A VEDIS, RNSAV, KDREC, BUF, RNRET, INDX, NPAG 00011090	6 00011090	:100	1001
INTEGER FSTPN, LSTPN, AVEDIS, RNS AV, RDR EC, BUF (10,3)	00011100	61,1	1072
INTEGER DHL, OPEK, RT, SET, PI, P2, P3, P4, TLOC(10), INDX	01111000	1111	1072
INTEGER I.J. RN. SUNFMI, RETP. TOTI, TOT 2, PAGET (10)	00011120	1.12	1674
EQUIVALENCE (\$SAVE(1), 0ML), (\$SAVE(2), 0PER), (\$SAVE(3),RT).	0001130	11113	1675
1 (\$SAVE(4);SET)	07111000	11114	1076
END: DECLARATIONS;	00011150	Siti	1077 TO 1
DMC = PI	07;11000	1116	SOM G
OPER = P2	06611170	1117	1075
RT = P3	00011180	1116	1660
SET = P4	06011190	1110	1001
END: INIVIALIZATION;	00011500	uzil	
SEIZE: FACILITY=STOKE;	01711000	1121	
C DETERMINE IF THE "SPACE MANIGEMENT PAGE" NEEDS TO BE ACCESSED	002211000	22.1	DDC
C (ON 'STORE' WHEN NOT ENDUCH SPACE AVAILABLE IN PRESENT PAGE)	00011730	1123	
RN = \$UNFHII1,1000,41	06011240	124	1066
JFIRM .LT. 9901 GD TG 200	06211000	11.56	1067
C MUST ACCESS THE SPACE MANAGEMENT PAGE	000:1:500	1511	1018
C GET THIS PAGE IF NOT ALREADY IN THE CACHE	0001173	tin)	10.89
RIMET = FSTPN	000:1249	1150	1040

Name and Address of the Owner, where

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		The second secon	The same of the sa
STANDARD INPUT STREAM LISTING			
DATE 09/10/77			
IMPUT CARM IMAGE 3050	7080	SEC-10 LEV SCHOOL	F. P. L. F.
PARAMETER LIST = (P1,P2,P3,P4);	00011530	1153	1115
******* DBMS CGNNECT DML PROCESSING FUNCTIONS*******	00011540	1154	1114
COMMON /WK/ TLOC	00011550	i les	1115
COMMON /PE/ PAGET	00011560	95,1	3716
COMMON /PAGAR/ FSTPN.LSTPN.AVFOIS.RNSAV.ROREC.BUF.RNRET.INDX.NPAG	PAG 00611570	1919	1:17
INTEGER FSTPN.LSTPN.AVEOIS, RNSAV, ROREC, BUF(10,3)	00011560	9966	11:16
INTEGER DAL, OPER, 4T, SET, PI, P2, P3, P4, TLOC(1C), INDX	06511000	1169	200
INTEGER 1, J.RN, SUNFHI, KETP, TOT!, TOT2, PAGET (10)	00011600	07.1	77 DE 1
EQUIVALERCE (\$SAVE(1), OHL), (\$SAVE(2), OPFR), (\$SAVE(3), RT),	00011610	1911	1 004
1 (SSAVE(4),SFT), (SSAVE(7),RN), (\$SAVE(8),1)	000:1620	1162	1162
END: DECLARATIONS;	00011630	1163	F211
DHL = P!	00011640	1164	. +211
OPER = P2	00011650	1163	5211
RT = P3	00011660	1164	11.50
SEY = P4	00011670	1167	11177
CNO: INITIALIZATION:	00011680	1148	1126
SETZE: FACILITY=CONECT:	00011:90	1169	6311
ASSUME UNIFORM DISTRIBUTION OF 1-5 MEMBER RECORDS TO BE CONNECTED	TED 00011700	vLii	1.30
RN = \$UNFHI(1,5,5)	01711000	1111	1151
I = 0	05711000	2211	1132
150 1 = 1 + 1	00011730	٤٤١٠ .	1133
1F (1 .GT. RN) GG TO 500	0003117-0	72.1	1134
CALL GENRP	000:17:0		1135
INVOKE: SERVICE=PAGHRS, PARAMETER LIST=(1);	00011760	1176	1136

DATE 09/10/17				
INPUT CARD IMAGE 3030505050	08	TNFUT AL	ALT INPUT P	REF-NO
WAIT RETURN: SERVICE=FAGMP;	00011770	1.77		1:37
GG TO 150	00011780	1176	-	1136
500 CONTINUE	00011750	1179		1:36
RELEASE: FACILITY=CONECT;	00011600	1180		1140
END: ENDO SERVICE:	CGC11E20	1911		1711
	00011620	1182		
	00011830	1153		
ENDO SERVICE: 10=DISCON, NAME= DYSCOS, SAVE AREA SIZE=10,	00011640	7611		1!42
PARAMETER LIST = (PI.P2.P3.P4);	06011650	1195		1143
C ******* DBMS DISCONECT PRECESSING FUNCTIONS******	00011600	1146		11:44
CCHMON 74K/ TLOC	00011870	Lyii		111-5
COMMON /PE/ PAGET	000:1660	11.48		1140
COMMON /PAGAR/ FSTPN.LSTPN.AVEDIS, RNSAV, RDREC, BUF, RNRET, INDX, NPAG 000:15"0	6 000:16"0	1190		1147
INTEGER FSTPN, LSTPN, AVEDIS, RNS AV, RDREC, BUF110,3)	00011300	1190		1144
INTEGER DHL, DPER, RT, SET, PI, PZ, P3, P4, TLOC(10), INDX	00011910	ilei		1145
INTEGER I.J.RN. SUNFMI, RETP. TOT!, TOTZ, PAGFT(10)	02511000	2611		7.50
EQUIVALENCE (SSAVE(1), DML), (\$\$4VE(2), OPER), (\$\$AVE(3),RT),	06211000	1143		1::1
I (\$\$AVE(4), \$5T), (\$\$\$VE(7), \$N), (\$\$AVE(6),1)	06011540	7011		1152
END: DECLARATIONS;	0:611000	1196		1153
JMC = PY	09611000	9511		*500
OPER = P2	00011770	Luii		11:5
RY = P3	0001190	Boil		115.
SET = P4	05511000	1100	MARY COST AL DESIGNATION OF	1157
END: INITIALIZATION:	00071000	1 :ne		1156

, ,				
STANDARD INPUT STREAM LISTING				
DATE 09/10/77				
INPUT CARD IMAGE2030	80	IN-PUT SED-NO	LEV SEC-NU	MEMBER AND KENTER
SETZE: FACILITY=DISCON;	00012010	1021		1159
C ASSUME UNIFORM CISTRIBUTION OF 1-5 REFORDS IN BE DISCONNECTED	00015050	1,702		1160
RN = \$IRVFHI(1,5,6)	000:2030	1203		13.61
0 = 1	07071000	1204		1262
150 I = 1 + 1	00012050	1205		1165
IF (I .GT. RN) GO TO 500	00015000	1205		1164
CALL GENRP	02022070	1021		11.65
INVOKE: SERVICE=PAGMRS, PARAMFTFR LIST=(1);	00015060	126P		1:66
WAIT RETURN: SERVICE =PAGMR;	00012000	1200		1167
GO TO 153	0012100	1210		Liee
SCO CONTINUE	00012110	1511		11.65
RELEASE: FACILITY=DISCON;	00012120	1212		11.15
END: ENDO SEAVICE;	00012130	1713		1111
	00012140	1214		
	00012150	Sici		
ENDG SERVICE: ID=ERASE, NAME= FRASES, SAVE AREA SIZE=10,	00012160	\$12.		2411
PARAHETER LIST = (P1,P2,P3,P4);	00012170	12.7		:17:
C ******* DBMS EKASE DML PROCESSING FUNCTIONS******	00012160	1210		. 174
COMMON /WK/ TLUC	00012150	1219		1175
COMMON /PE/ PAGET	00012200	122i		11.76
COMMON /PAGAR/ FSTPN.LSTPN.AVFDIS.ANSAV.ROREC.BUF.RNRET.INDX.NFAG GCG12219	\$ 00012219	1221		11177
INTEGER FSTPN.LSTPN.AVEDIS.RHSAV.RAREC. EUF(10.3)	00012220	1221		1178
INTEGER OPL. UPER, KT. SET. PI. P 2, P 4, TLOC (10), INDX	00012230	1223		1175
INTEGER I . J. A. H. SUNFMI, RETP, TOTT, TOTZ, PAGFT (10)	00012240	1224		1160

FAGE

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING	STANDARD INPUT STREAM LISTING					:
657 0 57 700						
0A1E 09/10/77						
INPUT CARD IMAGE2030304r506070	99	TUPNI TUPNI	LEV SEG-NU	REMEE'N		
150 I = I + 1	C0012490	1249		1205		
IF (I .CT. RN) GO TO 500	00012500	1250		1200		
CALL GENRP	00012510	1561		1207		
INVOKE: SERVICE=PAGHRS, PARAMETER LIST=(1);	07612520	125.2		1208		
WAIT RETURN: SERVICE=PAGHR;	00012530	1267		1205		
60 10 150	00012540	1254		1510	TH	
530 CONTINUE	00012550	1255		17.11		
RELEASE: FACILITY=ERASE;	00012560	1256		1212	OF.	
END: ENDU SERVICE:	06612570	1257		1213	S	1
	00012580	1258			BES	BO
	00012550	1254				a co
ENDO SERVICE: ID=PAGMR, NAME= PAGMRS, PARAMETER LIST=(RW),	00012600	1260		1514	LIA	X
SAVE AREA STZE=TO;	00012610	1921		1515	000	-
C ***** BUFFER MANAGEMENT PRUTTME****	00012620	1262		1216		النطا
C (IDMS " HODIFIED LEAST-RECENTLY-USED ALGORITHM)	00012630	. 263		1.17	1	
C THIS ROUTINE IS RESPONSIBLE FOR BRINGING THE REQUESTED PAGE	90012c40	1366		11.16		
C (I.E., RARET) INTO THE PAGE BUFFER. THIS MAY INVOLVE REWRITING	60012659	136F		1215		90
C ONE OF THE PAGES IN THE PUFFER.	00012563	1756		1550		_
C INPUT PARAMETER - 1=KEAD, 2=WPTTF	60C12r70	1247		1521		-
CDMMON ZWK/ TLDC	09012680	1:68		1424		
COMMON /PE/ PAGET	00012690	1260	The state of the s	1.13		
COMMON /PAGAR / FSTPN, LSTPN, AVEDIS, RNSAV, RUREC, BUF, RNRET, THD X, NPAG 00012700	000121000	0221		1,50		1
INTEGER FSTPM,LSTPM, AVEDIS, RNSAV, ROREC, BUF (13,3)	31721330	1521	-	11.25		
INTEGER DML.OPER.RT.SET.PI.PZ.P3.P4.TLOCITO. INJX.RM.RDWT	00012720	1277		11.26	1	

FAGE :4				_					7	HIS	PAG	L	BE	st q LSH	WAL:	ITI DO	P.A.	CTI				:)	
•		PLT TREUT MENUEL TV SEC-RU FTF-RU	1771	17.26	677;	1,50	16.31	1132	123.	1:34	37.1	17:0	15.37	1621	1239	1.40	1721	2421	1243	1.4.	1745	1640	12.1	1:46	14.49	1,50
er digge contrabilità de la conse		SEO-NO LEV	1273	1274	1275	1276	1237	1279	1279	1260	Laci	1282	1263	1584	1.60	1286	1287	1269	1289	1:00	1501	1252	1,207	1 300	1501	9061
LATOR		060	00012720	00012740	00012750	00012760	00012770	09677000	00012790	00012100	00012810	60012620	60012630	00012840	00012850	00012660	00012870	00012680	00012890	00012600	00012510	00012520	00012930	00012640	00012950	09671000
INFORMATION PROCESSING SYSTEM SIMULATOR STANDARD INPUT STREAM LISTING	DATE 09/10/77	INPUT CARD IMAGE1020303040505050	INTEGER 1, J.RM, SUNFMI, RETP, TOT7, TOT2, PAGF1110), 1X	EQUIVALENCE (SSAVE(1), DML), (STAVE(2), OPER), (SSAVE(3),RT),	1 (\$5AVE(4), \$ET), (\$5AVE(5), \$BI), (\$5AVE(9), IX), (\$5AVE(10), \$DWT)	C PAGE FAULT STATISTICS ARFA - PAGET(10) -	C 11) WINBER OF PAGES CONCURRENTLY IN THE PAGE BUFFER	C (2) PASE SIZE (GYTES)	C (3) PAGE WRITES	C (4) CACHE MISSES ON WRITE	C (5) PAGE REAUS	C 16) CACHE MISSES ON READ	C (7) NUMBER OF PAGES REWRITEN REFORE PAGE SPACE REUSED	C (8) NUMBER OF PAGE REPLACEMENTS WITHOUT REWAITTING	C (9) NUMBER OF PAGE NUMBERS GERFRATED (BY GENER)	C (10) TOTAL DISTANCE BETWEEN PAGE REFERENCES	END: DECLARATIONS:	RN = RURET	KDWT = RW	END: INITIALIZATION;	SEIZE: FACILITY= PAGMR;	IF(RDM) -LY. 1) ROWY = 1	IFTROWT -GT. 2) RDWT = 2	C UPDATE NUMBER OF READS/WRITES	IFTRDHT .Eq. 1) PAGFT(5) = PEGFT(5) + 1	IFTRDWI .EQ. 2) PAGET(3) = PAGET(3) + I

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INFORMATION PROCESSING SYSTEM SIMULATOR	æ			1974	55
STANDARD INPUT STREAM LISTING					
DATE 09/10/77					
THPUT CARD IMAGE202030404050606076	<u>jų</u>	Jiput SE G-PiC	ALT INPUT	MEMDER LEF-10	
C DETERMINE IF PAGE IS IN BUFFER POOL	00751000	1297		1221	
00 15 I = 1,NFAG	00012580	1208		1252	
]X = 1	00012950	1290		1253	
7F(60F(1,1) .Eq. RN) GO TO 7A	00023000	1 300		1754	
TO CONTINUE	000:3010	1301		1255	
C PAGE WAS NOT IN THE BUFFER POOL	00013020	1302		1256	
IFTROWT .EQ. 1) PAGFT(6) = PACFT(6) + I	00013030	1303		1257	
IFIRDWT .EQ. 21 PAGFT(4) = PAGET(4) + 1	00013043	1304		1258	
C SIMULATE RECORD RETRIEVAL FROM DATA BASE - ELAPSED TIME = 36 MS.	00013050	1305		1,39	DH TR
C (SYSTEM HOOK)	09013000	1306		17.60	IS P
PROCESS: TIME=36.0;	07051000	1307		1261	AGE
C STORE THE PAGE IN THE BUFFER POOL	00013060	1306		1262	IS
1x = 0	0606 1000	1369		1263	BES
30 IX = IX + 1	00(13100	1310		1264	TQU
1F(1X .GT. NPAG) GO YO YO	00013:10	1311		1,65	IALI NO
1F(BUF(1X,2) .NE. 0) 60 T0 30	00013120	1312		1266	DOC
C DETERMINE IF THE PAGE MUST BF WRITEN BEFORE IT IS REPLACED	00013130	Fiel		1267	RA
1F(BUF(1X+3) .EG. 0) GO TO 40	00013140	7.6.1		loc	A PA
C PAGE TO BE REPLACED HAS BEEN MODIFIED - REWAITE	00013150	Siĉi		121.9	- CY
PAGET(7) = PAGET(7) + 1	00013160	13.6		1:10	13
C SIMULATE RECORD STORAGE TO DATA BASE - ELAPSED TIME = 36 MS.	07161000	1317	A COLOR OF THE PERSON NAMED IN COLOR OF THE P	n:n	
C (SYSTEM HOOK)	00013160	9161		2772	
PRICESS: TIME=36.0:	06:8:90	1300		1473	
# W #3	06.51000	025.	1	17.74	

UAIS 07/10/11				
INPUT CARD IMAGE2020304040506070	F0	SEO-NO LE	LEV SEU-NO S	MERCER
40 PAGET(8) = PAGET(8) + 1	0.0013210	1321		1.75
	C061322C	1322		1276
45 00 50 J = 1.NPAG	00013230	1323		11.77
1F(BUF(J,2) .GT. 0) BUF(J,2) = RUF(J,2) - 1	00013240	1324		3276
50 CONTINUE	00013250	1995		1,76
STORE THE NEW PAGE IN THE PUFFER POIN.	00013263	1326		12.60
BUF(IX+I) = RN	06513270	1327		1261
EUF(IX,2) = NPAG - 1	00013260	1328		THY FR
BUF(IX,3) = 0	06781000	6221		1.03
	00013300	044.		1.64
	01561000	1331		1,05
	60015220	13:5		1286
70 INDX = 1X	00013:30	1333		1.67
IFTROWT .EQ. 1) GO TO 80	00013340	1334		1448
TURN MUST-WRITE SHITCH "ON"	00013350	1336		1.69
BUF(INDX+3) = 1	00013340	1886		0621
80 CONTINUE	00013270	(LEI		1.51
RELEASE: FACILITY=PAGMR;	00013.80	133₽		767.
END: ENDO SERVICE;	05551000	beci		12.93
	00013400	Unil		
	00013410	1361		
PROCEDURE: NAME=GENRP, TYPE=SUFROUTINE;	12420	1342		1531
THIS PROCEDURE GENERATES A UNIFORMLY DISTRIBUTED RANDOM PAGE	00013430	6361		1.55
NUMBER OVER THE SPACIFIED PAGE RANGE WITH THE DESIRED AVEGAGE	00013440	4251		4,

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STANDARD INPUT STREAM LISTING	STANDARD INPUT STREAM LISTING	-	7 1784		
DATE 09/10/77					
INPUT CARD IMAGE	PO	IMPUT SEC-NO	ALT FIAUT	PERBYA Ket-LC	
C DISTANCE BETWEEN SUCCESSIVE PAGES. USING THIS ALGORITHM.	00013450	1345		1673	
C *CLUSTERED * PAGE REFERENCES CAN RE GENERATED REGARDLESS UF THE	00013460	1346		1.56	!
C SIZE OF THE DATA BASE.	00013470	±76 i		1299	-
3	00013480	4761		ison	
C FSTPN - LOW ORDER BOUND ON GENERATED PAGE NUMBERS	0613490	645:		13.01	
C LSTPN - HIGH ORDER SCUND ON GENERATED PAGE NUMBERS	00013500	1350		1302	
C AVEDIS - THE GESTRED AVERAGE DISTANCE BETWEEN SINCESSIVE PAGE	0135100	1361		1305	
C REFERENCES	00013520	1352		1304	
C RNRET - THE COMPUTED PAGE NUMBER	06351000	1363		1265	
COMMON ZWKZ TLOC	00013540	7366		1306	F
COMMON 7PE7 PAGET	06351,000	1355		1367	ROU
COMMON /PAGAR/ FSTPN.LSTPN.AVEDTS,RNSAV,KOREC,BUF,RNKET,INDX,NP4L	00013560	13.6		1506	PAC
INTEGER FSTPN.LSTPN.AVEDIS,RMS.AV.RDREC,BUF(10,3)	01351000	1357		3369	Y
INTEGER DML, OPER, RT, SET, PI, P2, P3, P4, TLOC(10), INDX	00013550	1366		1310	PEAL
INTEGER I.J.RN.SUNFHI.RETP.TOTI.TOTZ.PAGFT(10).K.L	00013590	1369		13.11	SHI
INTEGEK KNI, KNZ, VARZ	00013600	1340		1512	TX
REAL SUNFM, V, RN3	01961000	1361		1:13	bad (
C INITIALIZE	00013520	1362		1114	3
MAXPN = LSTPN - FSTPN	06013630	1363		1515	-
IF(MAXPN .LE. 0) GO TO 100	096:3640	3.36.4		1:16	+
RNI = RNRET	00013650	1365	The second secon	1317	
C CCMPUTE KN2 USING THE FULL PANGE OF PAGES	00013660	1346		1316	
IO RNZ = SIMFMIFSTPN+LSTPN+BI	06013670	1367		1:14	
VARI = RNI + AVEUIS - FSTPN	00013660	136.9		1326	

8.			*						T	HIS	PAG	E IS	BE	ST Q	UAL.	DO	PR	LETI						1		••••		
P AGE			MEMBET KLF-NO	1321	1522	1323	1:24	1325	1256	1327	1528	1329	1:50	1:31	1532	1333	1534	1:35	17.30	1537	1338	1334	1:40	1341	1542	. 34.	1364	
			ALT TIPUT LEV SEC-NO																									
			Sec-40	1360	1370	Itei	1372	1373	761	1375	1376	1377	1378	biel	1300	:38;	1382	1363	1364	1345	1386	1381	1 788	638.1	13.0	1501	1361	
ATOR			04	06981900	00013760	01161000	06013720	00013720	07613740	09013750	COC 13760	07751000	00013780	0615:000	00013600	00013816	00013520	00013830	00013840	00013850	0961366	00013670	00013680	06013890	00581300	00013610	00013720	60
INFORMATION PROCESSING SYSTEM SIMULATOR	STANDARD INPUT STREAM LISTING	DATE 09/10/77	INPUT CARD IMAGE10	IF (VARI . GT. MAXPN) VARI = MAXPN	VAR2 = RNI - AVEUIS - FSTPN	IFIVARZ .LT. 0) VARZ = 0	C DON'T RECOMPUTE RN2 IF NOT OUTSIDE RANGE RNI+-AVEDIS	IFIRNZ .GT. ANI+AVEDISI GO TO 20	IFTRNZ .LT. KNI-AVEDISI 60 TO 20	GO TO 50	20 CONTINUE	C Y * PROBIRNZ > RNI+VAR) OR PROBIRNZ < RNI-VAR)	IF(RNZ .GT. RNI) Y = 1 (VARI/(MAXPN*1.0))	IF (RNZ .LT. RNI) Y = (VAR2//MAYPHAY.O))	C TEST BOURDS ON Y - ALGORITHM CAULD LARP - EXIT IF OUT OF KANGE	IF(V .LY. 0.) GO TO 50	IF(V .6T. 1.) 60 TO 50	C GENFRATE RN3 TO TEST AGAINST Y	RN3 = \$UNFM(0.0,1.0,8)	36 CCNTINDE	IFIRM3 .GE. VI GO TO 50	C RECOMPUTE RN2 - KANGE CRITERIA NOT MET (ON AVERAGE)	K = RNI - AVEDIS	IFIK .LT. FSTFNI K = FSTPN	L = RN + AVEDIS	JFIL .6T. LSTPN) L = LSTPN	RNZ - SUNFMICK, C. 6.1	1030405070

*								-			18		E IS	BE	ST Q SHE	UAL.	I TY	PRA	er:		43				
1461	MIMEE'S	1342	1.346	1547	1346	1245	5401	1351	1357	1:53	1354	1355	1356	1:57	1256	1254			1.00	13:01	3362	; e ;	1.64	1.45	1:66
	יואריו בון אריופאן ביי ארי	1 343	1 304	1306	9601	1961	139R	77E I	1400	luzi	1405	أبراغ	707;	1405	9071	1401	1400	6071	141.7	1411	1412	14:3	וקוע	1474	1416
ſſſŔ	80	06613930	RETURN00013940	00013950	09681000	00013970	00013980	06581000	00014000	00014010	00014050	00014030	07071000	06014050	00071000	06014070	000034060	04071000	00014100	00014110	00014150	00014130	00014140	00014150	06614160
INFORMATICN PROCESSING SYSTEM SIMULATOR STANDARD INPUT STREAM LISTING	INPUT CARD THAGE	101F = (KN2-RN1)	C AVERAGE DISTANCE CRITERIA IS SATISFIED - UPDATE VARIABLES & R	50 RNRET = RNZ	PAGFT(5) = PAGFT(9) + 1	PAGFY(10) = PAGFT(10) + IARS(PNI-RN2)	IDIF = (RN2-RN1)	C BYPASS SECOND WRITE OF GENRP DATA	IFIRN3 .LT. Y) GG TO 500	20 10 500	130 CONTINUE	C ERROR RETURN ON FSTPN LSTPN VALUES - FSTPN > LSTPN	RNRET = \$UNFMI(LSTPN,FSTPN,8)	500 CURTIMUE	KETÜRN	FND: PROCEDURE;			EXO SERVICE: ID=RUFST, NAME=PUFSTS,	SAVE AREA S12E = 10;	C PRINT BUFFER PAGE FAULT STATISTICS	COMMON /PE/ PAGFT	INTEGER PAGET (10), 7071, 7072	REAL MISA	END: DECLARATIONS;

AD-A069 543

OHIO STATE UNIV RESEARCH FOUNDATION COLUMBUS
A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS-ETC(U)
DAA629-77-6-0203

ARO-15356.1-A-M

NL

3.03

AD-604.7

DESCRIPTION OF THE PROCESSING SYS-ETC(U)
DAA629-77-6-0203

ARO-15356.1-A-M

NL

END
ONL

DATE 09/10/77				
INPUT CARN IMAGE 3030505070.	70	1NFPT SFC-1-0	ALT HAUT	MLNLF"
END: INITIALIZATION;	00014170	7117		1367
WRITE(6+10)	00014160	1416		1308
10 FORMAT("1", 20X, "CACHE PAGE FAULT STATISTICS"//)	06014190	1410		17.69
WRITE(6,20) PAGFT(1)	00014500	1420		1576
20 FORMAT(" ',10x, CACHE SIZE (# PAGES): ',12)	00014210	1271		เรน
WRITE(6,30) PAGFT(2)	00014220	1422		1:12
30 FORMAT(' ',10x, 'PAGE SIZE (RVTES): ',14)	00014230	1423		1373
HISR = PAGFT(4)/(PAGFT(314).0)	00014240	1454		137.
WRITE(6,40) PACFT(3), PACFT(4), MISR	00014250	1456		1375
40 FORMATI'0',10X, PAGE WRITE REFERENCES: ",16,	00014260	1756		137c
1 3X, CACHE MISSES: ", 16, 3X, "MISS RATIO: ", F6.4)	00017510	Lč7i		11577
MISR = PAGFT(6)/(PAGFT(5)*1.0)	CGC 14280	1629		15.78
WRITE(6,50) PAGET(5), PAGET(6), HITER	60014290	1429		1:79
SO FDRMATI' ',10x, PAGE READ REFERENCES: ',16,	00014300	1430		1360
1 3X, CACHE MISSES: ", 16, 3X, "MISS PATIO: ", F6.4)	00014510	1431		17.61
TOTI = PAGET(3) + PAGET(5)	66614320	1422		12.62
TOT2 = PAGFT(4) + PAGFT(6)	06014330	1433		13.63
HISR = T0T2/(10T1+1.0)	00014340	1414		1364
WRITE(6,60) 7011,7072,MISR	00014350	1416		1365
60 FORMATIFO ,10X, PAGE YOTAL REFERENCES: ",16,	00014360	1436	The state of the s	1286
1 3X."CACHE HISSES: ", 16, 3x, "HISS PATIO: ", F6.4)	60614370	נביזו		1:17
WRITE(6,70) PAGFT(7)	00014350	1436		44.
TO FORMATI'S . 10X PAGES REWRITEN REFORE SPACE REUSED: "16)	04541930	1634		*:
URITEIS. 801 PAGETIBI	00014400	1440		05.

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			73 1971
STANDARD INPUT STREAM LISTING	STING		the same of the same state of
DATE 09/10/77			
INPUT CARD IMAGE10102030405050	70+0	INFUT A	ALT TIPUT NEGETY LEV SEG-60 PEF-70
80 FORMATI 10x, . PAGES REPLACED WITHMIT REWKITING: ". I.6.)	00014410	1441	1361
MISR = PAGFT(101/(PAGFT(9)+1.0)	00014450	1442	2551
WRITE(6,90) MISR	00014430	1443	1393
90 FORMATI'O', 15X, MAVERAGE DISTANCE BETWEEN SUCCESSIVE PAGE",	00014440	144.5	1354
1 REFERENCES: ",F9.4//)	66614450	574;	1355
END: EXO SERVICE;	00014400	1446	3561
P. (3)	00014470	1447	
2.5	00014480	677 i	
END SERVICE: ID=SNAPST, NAME=SNAPSS,	06771000	1 149	1357
SAVE AREA SIZE = 10;	00617800	1450	1398
S C PRINT STANDARD IPSS STATISTICS AT THE DESIGNATED INTERVALS	00014+10	1761	1239
G 3 END: DECLARATIONS:	00014520	1452	0071
ENG: INITIALIZATION;	00014530	6571	1401
CALL SSNAP	00014540	7371	1402
END: END SERVICE;	05471070	5 3 7 i	140.5
45	00014540	1456	
	00014570	1571	
END: SYSTEM RESOURCES;	03614580	1450	140.
	06541000	1440	
	00014600	1460	A training against the control of th
DEFINE EXOCENDUS EVENT STREAM: NAME=FX, COMPILE=YES,	00014610	1971	
DISPOSITION=KEEP;	00014620	3 17 1	
	00014430	1771	The second secon
EXOCEMBLE EVENT: TOMEND. TIME = PROCETTMENT:	00014640	1464	

60	SEC-FO LEV SEC-EU	-c) PEIEFF
00014650	1466	2
00014660	3371	.5
00014670	1467	
00014680	1448	1
06971000	1459	41
00014100	1470	3
0:45:000	1451	4
02441000	1472	, b
00014720	1473	,
00014140	1474	10
00014750	9271	:
00014160	1476	1.2
000:4110	2676	13
00014789	9171	1,
0614190	1479	
00014800	i LAC	9 ;
00014-10	1471	
0001420	1462	2
00014830	1441	7-
06014840	7471	3.7
00014850	1485	
00014660	1456	22
02371000	1411	ż
0001480	9471	•
	00014650 00014670 00014670 00014680 00014760 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770 00014770	

-

0170700 0110					
:					
INPUT CARD IMAGE	7060	THEUT SEQ-NC	417 11:PUT	REPLEK K: F-MC	
PROCEDURE: NAME=TIME2, TYPE=SURGOUTINF;	00015130	1:13		37	!
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH THE PAGE	00015140	7151		4.7	
C CACHE STATISTICS AKE TO RF PRINTED	00015150	1516		0.7	
REAL*R ITIME	09151000	16.6		1.5	
DATA 111ME/7.2000010+06/	0215170	1611		2;	
CALL SSMFAL(ITIME, SYSCOM, 34)	00015160	9151		55	THI FRO
REYUKN	00015160	51.1		75	SPI
END: PROCEDURE;	00015:00	1530		*;	GE.
	00015210	.291			IS B
PROCEDURE: NAME=YIME3, TYPE=SU3ROUYINE;	00015226	223.		75	EST LSt
C THIS PROCEDURE ESTABLISHES THE TYME AT WHICH A SNAPSHOT	00015230	1523		1.5	QUA
C OF THE STANDARD 1PSS STATISTICS ARE TO BE PRINTED	06015240	7551		5.5	LIT TO D
KFAL*8 IYIME	06751000	1626		35	DC
DATA 111HE/7.20+06/	CCC15266	1631		23	
CALL \$SREAL (ITIME, SYSCOM, 34)	00015270	1631		6.1	
REYURN	00015280	1528		7.7	
END: PROCEDURE;	00015250	6631		[7]	
	00015300	0411			
END: EXOGENDUS EVENT STREAM;	00015310	1631		1.	
	02551000	1432			
	00015330	list			
DEFINE MNDEL: TO=SIOPRS.	00015340	7171			
COMPONENTS=(CF(SYSTEM.CF),FX(SYSTEM.FX));	065645000	1 5 5 6			
	00015360	9631			

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(1	•	,	1		(,			1	1	- 187 -	as la	1		1		,,,,,,		
	\$3					-			1				THIS P	Y Pur	LSH Q	UALI1	Y PR	ACTION	•		
	FA65			Mr Pt f k	۵.	5	4		•	7		5									
				ALT INPUT										7						1	
				SFC-ND	1537	1538	1639	1540	1751	1542	1543	1544									
	*			Po	00015370	00015280	06251000	00015400	00015419	02451000	00015430	00015440									
	SIMULATO	TIRG		70	ŏ	00	00	ŏ	00	0	ō	00									
	NG SYSTEM	TREAM LIS	11/01/60						00010.1.												
	INFORMATION PROCESSING SYSTEM SIMULATOR	STANDARD INPUT STREAM LISTING	DATE 09/10/77	IMAGE					-(TINE=288	: (ON=5											
•	NFORMATIO	STANDA	1	10.					NFF. STOP	LASILAVIS											
	1			30.	=EV1;	=EV2;	=£V3;		T, TRACE=	INATE=10.											-
				20	START, EX	SUFST. EX	SNAPST. EX		REAM=STAR	ERRUR=(TERMINATE=10, STATISTICS=NO);											
				INPUT	FQUATE: CF=START, EX=EVI;	FQUATE: CF=8UFST. EX=EV2;	EQUATE: CF=SNAPST, EX=EV3;		SINULATE: STREAM=START, TRACE=OFF, STOP=(TIME=28800010.),	13		END: MODEL:									
					FQ	10	EQ		18			ENDS									

APPENDIX C

RESULTS OF THE RESEARCH ACTIVITIES

This appendix contains the conclusions drawn from the research activities. These conclusions are based in a large part from the experiences gained with the IPSS SIDPERS/IDMS modeling activities. Recall that the purpose of the modeling activities was to demonstrate the feasibility and suitability of the Information Processing System Simulator (IPSS) with respect to modeling large complex information systems. The system modeled to demonstrate these capabilities was a host/backend processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system functions of task management, job scheduling and resource allocation.

The research proposal identified 164 different extensions to the then existing IPSS. As a result of this research, 139 new language and statistics gathering features have been incorporated into the IPSS. Other extensions not in the original proposal were also identified as useful features and are being added to the IPSS. The results are reported in the following three sections of this appendix:

Section	Area Covered
C1	Data Base Systems - specific extensions
C2	Information System - related performance statistics
C3	Computer Networks - specific extensions

The vehicle used for assessing the usefulness of the suggested IPSS extensions was to model the complex systems architecture described in previously in this report. For evaluative purposes, the proposed extensions to the IPSS are each assigned one of the following assessment dispositions:

Assessment Conclusion				Ca	te	eg	ory/Disposition
Not Examined							NE
Examined and Rejected							
Examined and found to be useful						Ť	
- implemented as proposed							EU-I
- implemented but merged into							
another facility							EU-IM
 defined and not implemented 							EU-D
- defined and not implemented							
but used in paper model .							EU-DP

Dispositions D and DP are intended to show the depth to which an IPSS language statement under consideration has been examined. Paper models demonstrate effectively the definitional capability of IPSS to characterize information system activities of interest, and as such are a firm indication of the desire of implementing the statement into IPSS.

C1. IPSS LANGUAGE STATEMENTS FOR CHARACTERIZING DATA BASE SYSTEMS

The IPSS system was designed to analyze information processing systems in their totality. To date, IPSS mechanisms for the definition and evaluation of hardware, computer system software and storage level data structures have been extensively developed and tested. The next phase in the continued development of IPSS was the incorporation of features for the description, manipulation and evaluation of logical data structures of the kind normally encountered in modern data base systems.

The goal for this aspect of the research was to identify the usefulness of incorporating a set of generalized data base-related characterization facilities into the IPSS. Data Base Systems (DBS) characterization is provided through two components: one definitional in nature which is specifically designed to represent logical data structures, and the other procedural to represent application and DBS

software processing. These components have been named the Data Base Structure component and the Data Base Access component, respectively. The relation of these two components to the existing model architecture is shown in Figure Cl-1. For purposes of identification they have been titled the IPSS/DBS submodel architecture. The proposed language for characterizing DBS facilities for the IPSS/DBS are identified in Table Cl-1. Also shown are the results of their evaluation as part of this research effort.

Cl.1 The Data Base Structure Component

The purpose of the Data Structure component is to provide the modeler with a set of facilities which allows one to define logical data structures and to characterize the relationships among them. This component is defined independently of the data base access component (See Section C1.2) and thus these facilities can be applied to a variety of DBS architectures and application environments.

The Data Base Structure component facilities serve the following functions:

- They define templates for schemas consisting of record types and sets, instances of which can be transient entity definitions in the DBA component;
- They allow data base structures to be viewed as a collection of contiguous logical areas and to map these onto non-contiguous domains of a DBS's storage structure and also onto other logical data structures;
- They specify allocation policies of record type records to the sets in which they participate; and
- 4. They allow the characterization of access paths between record types of a schema.

The Data Base Structure component permits the modeler to investigate the effects on system behavior caused by alternate set, record type, and access path definitions. The definitional facilities provided allow the modeler to investigate a wide spectrum of logical data structure organizations and allocation policies.

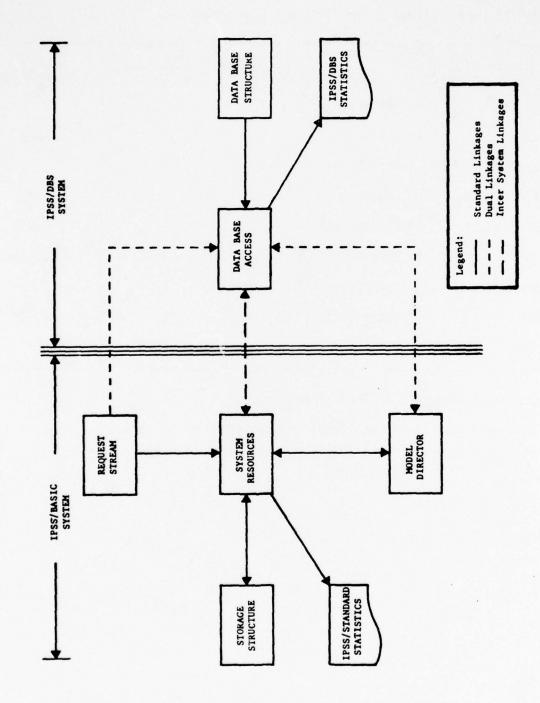


Figure C1-1 IPSS System Architecture with UBMS Extensions

Table C1-1. IPSS/DBS Facilities

		Status	in IPSS	
Statement	Туре	Current	Additional	Disposition*
BUFFER POOL	Definitional	X		
DML SERVICE	Executional		X	EU-I
ENDOGENOUS SERVICE	Executional	X		
EXOGENOUS SERVICE	Executional	X		
EXTENT	Definitional		X	EU-IM
REALM	Definitional		X	EU-I
RECORD TYPE	Definitional		X	EU-I
SCHEMA	Definitional		X	EU-I
SET TYPE	Definitional		Х	EU-I

*Disposition codes are defined as follows:

EU-I : Examined, found Useful - Implemented as proposed;

EU-IM: Examined, found Useful - Implemented by Merging

into another facility.

Schema Definition

The DEFINE SCHEMA - END SCHEMA statement sequence defines a logical data structure of inter-related elements. Within this sequence are the definitions of sets, record types and extents.

Set Definition

The SET TYPE statement characterizes connections between record types in a schema. This statement identifies a collection of record types, names the collection, specifies the type of linkage between the record types, identifies the type of participation (either owner or member) of the record types in the set, and characterizes the relative frequency of set occurrences.

Record Type Definition

The RECORD TYPE statement characterizes those properties of a record type which are independent of secondary storage. The physical attributes are expressed in terms of the record format, the number of record occurrences, and record origin. An association with an extent is also provided for linearization. These physical characteristics are independent of the record location in secondary storage. (The IPSS Storage Structure component performs the required mappings from logical records to storage structure representation.) This Data Base Structure component facility completes the record type definition of the Data Base Access component whose major purpose is to define the source of the records.

Extent Definition

The EXTENT statement defines one or more extent facilities. The extent provides the interface between a record type's logical address space (defined via the RECORD TYPE statement) and the Realm specification defined via the REALM statement. Any number of extents may be defined.

Realm Definition

The REALM facility is used to linearize the address space of a schema. The REALM statement provides the external identifier to the REALM facility and provides the space management area associated with the REALM. Record types may be assigned to realms through the EXTENT statement. As many realms as desired may be defined for each schema definition. IPSS views a realm to be the equivalent of a virtual data set.

Cl.2 The Data Base Access Component

The Data Base Access component is the central component in an IPSS/DBS defined model. Within it are contained the definitions of all the required DBS and application level software. Hardware resources are limited to main memory buffers used by the DBS software and user work areas. All DBS-related entity type facilities are defined within this component. This subsection describes those facility definition statements which can appear within a Data Base Access component definition. These represent extensions to the basic IPSS language.

In addition to the facilities identified below, the modeler can include any number of standard IPSS procedure facilities and Fortran Subprogram definitions. These defintions can appear anywhere in a Data Base Access component definition with the exception that they cannot be imbedded within any other facility definition. This implies that a Procedure or Subprogram facility definition cannot appear within:

- 1. Another Procedure,
- 2. Another Subprogram,
- 3. An Endogenous Service, or
- An Exogenous Service.

They may, however, be referenced from within any of the above. Procedure and Subprogram syntax and semantics are discussed in Volume I of the IPSS Snytax & Semantics document (Section 3.6).

DML Service Definition

The DML SERVICE - END DML SERVICE statement sequence corresponds in function to the standard IPSS Endogenous Service definition. It has been named DML service in the DBA component to more accurately reflect DBS service request processing.

Record Type Definition

The RECORD TYPE statement defines a record type facility within the Data Base Access Component and identifies the source of the records which are created for the specified record type facility.

Schema Definition

The SCHEMA statement identifies a schema facility to the Data Base Access component. The named Schema is associated with a schema definition in the Data Structure component through EQUATE statements in the Model Stream Component.

C1.3 IPSS/DBS Oriented Built-In Facilities

IPSS/DBS Built-in Procedure Facilities provide the modeler with a large number of predefined capabilities whose purpose is to reduce the time needed to develop data base management system models. These procedural facilities free the modeler from the detailed mechanics of the simulation process. The built-in procedures have been grouped into three categories based on features or functions common to each group they are:

- 1. I/O-oriented procedures,
- 2. Schema-oriented procedures, and
- Data structure traversal-oriented procedures.

Table C1-2 contains an alphabetized listing of all the IPSS/DBS Built-in procedures and denotes the dispositions of the statements.

All Built-in procedure statements are converted to Fortran subroutine calls with the parmeters associated with that statement transposed to subprogram parameters. Formal descriptions for these IPSS facilities

Table C1-2 IPSS/DBS Built-in Procedures

		Stat	us in IPSS	
Statement	Туре	Current	Additional	Disposition
ACQUIRE RECORD	Executional		X	EU-DP
ALLOCATE SCHEMA EXTENT	Executional		X	EU-IM
ALTER SET	Executional		X	NE
CONFIGURE CEQ	Executional	X		
COPY SET	Executional		X	NE
COPY ROUTE	Executional		X X X X	EU-DP
CREATE ROUTE	Executional		X	EU-I
CREATE SCHEMA	Executional		X	EU-I
CREATE SCHEMA EXTENT	Executional		X	EU-IM
CREATE TRANSACTION	Executional	X		
DESTROY ROUTE	Executional		X	EU-I
DESTROY SCHEMA	Executional		X	EU-I
DESTROY SCHEMA EXTENT	Executional		X X X	EU-IM
DISABLE EVENT	Executional	X		
OO TO LABEL	Executional	X		
ENABLE EVENT	Executional	X		
ND DO	Executional	X X		
FIND IN QUEUE	Executional	X X X		
IND MEMBER	Executional		X	EU-I
IND NEXT BUFFER	Executional	X	^	
IND ROUTE	Executional	^	X	EU-DP
REE BUFFER	Executional	x	^	20 01
SET BUFFER	Executional	Î		
GET SAVE AREAS	Executional	Î		
IDENTIFY OWNER	Executional	^	y	EU-DP
IDENTIFY SET OCCURRENCE	Executional		X X	EU-DP
INITIATE TASK	Executional) x	^	LU-Dr
		1 0		
MERGE QUES	Executional	X X X		
MODIFY BUFFER POOL	Executional	1 0		
MODIFY QUEUE	Executional	^	v	EU-DP
MODIFY REALM	Executional		Ŷ	EU-DP
MODIFY ROUTE	Executional		?	EU-IM
MODIFY SCHEMA EXTENT	Executional		÷	EU-DP
MODIFY SCHEMA RECORD TYP			X X X X	EU-DP
MODIFY SCHEMA SET	Executional		^	E0-0F
PLACE IN QUEUE	Executional	X		
POST SEMAPHORE	Executional			
POST SIGNAL	Executional	X		
PROCESS TIME	Executional	X		
PUT SAVE AREAS	Executional	X		FIL TM
RELEASE SCHEMA EXTENT	Executional		X	EU-IM
REQUEST DML SERVICE	Executional		X	NE
REQUEST SERVICE	Executional	X		
RESUME SCAN	Executional	X		
SET FACILITY STATUS	Executional	X		
SET INDIRECT REFERENCE	Executional	X		

Table C1-2 IPSS/DBS Built-in Procedures (Continued)

	T	Stati	is in IPSS	T
Statement	Туре	Current	Additional	Disposition
SET PRIORITY	Executional	X		
SET SYSTEM PARAMETER	Executional	X		
START QUEUE STATISTICS	Executional	X		
START USAGE STATISTICS	Executional	X		
STOP QUEUE STATISTICS	Executional	X		
STOP SIMULATION	Executional	X		
STOP USAGE STATISTICS	Executional	X		
STORE RECORD	Executional		X	NE
TERMINATE SERVICE	Executional	X		
TERMINATE TASK	Executional	X		
WAIT DML SERVICE COMPLET	Executional		X	NE
WAIT FACILITY STATUS	Executional	X		
WAIT RECORD	Executional		X	NE
WAIT SEMAPHORE	Executional	X		
WAIT SERVICE COMPLETE	Executional	X		
WAIT SIGNAL	Executional	X		

Examined, found Useful - Implemented as proposed Examined, found Useful - Implemented by Merging into EU-IM:

another facility

Examined, found Useful - defined and not implemented but EU-DP:

used in paper model

Not examined. NE

are not presented in this report. IPSS Syntax and Semantics, Vols. I and II presents the following items of information for each Built-in Procedure:

- Statement syntax,
- 2. The function of the statement,
- 3. Semantics for the statement's parameters,
- 4. A narrative description of the execution mechanism,
- IPSS generated Fortran source code which replaces IPSS statements and
- 6. Examples of the use of the statement.

Cl.3.1 IPSS/DBS I/O Related Built-in Procedures

The following Built-in Procedures are used to initiate I/O activity and to wait for the I/O to be completed. The ACQUIRE and STORE RECORD procedures determine the source of the record from the DBA component record type facility definition. If a service is specified as the source, then the service is automatically invoked. WAIT RECORD statements are defined to enable service processing to be suspended pending the I/O operation. This section also contains the definition of the statement for initiating and waiting for a DML Service. These procedures can be referenced from within either Procedure facilities or Exogenous and Endogenous Service facilities.

Acquire and Store Record

The ACQUIRE and STORE RECORD Built-in Procedures invoke a service which simulates the acquisition and storage of the specified record to and from the service in which the ACQUIRE or STORE statement appears. The identity of service to be invoked is contained in the definition of the specified record type facility in the DBA component.

Wait Record

The function of the WAIT RECORD Built-in Procedure is to place the specified transaction into the wait state unit1 the completion of the Endogenous service which the specified transaction invoked via an ACQUIRE or STORE RECORD statement. Service for the suspended transaction resumes after the specified wait criteria (count and request number) are satisfied. Each completed Endogenous Service routine satisfying the request number criterion decrements the suspended transaction's Wait Count by one and the delayed transaction is reactivated when the count reaches zero and is returned to the active state.

Request DML Service

The REQUEST DML SERVICE Built-in Procedure is used to initiate the execution of a DML Endogenous Service.

Wait DML Service Complete

The function of the WAIT DML SERVICE COMPLETE Built-in Procedure is to place the named transaction into the wait state until the completion of the specified Endogenous DML Service which it invoked. Service for suspended transaction resumes after the specified wait criteria (count and request number) are satisfied. Each completed Endogenous DML Service routine satisfying the request number criterion decrements the suspended transaction's wait count by one and the delayed transaction is reactivated when the count reaches zero and is returned to the active state.

C1.3.2 IPSS/DBS Schema Related Built-in Procedures

The following Built-in procedures are used to change the current state of facilities defined in the DATA Structure component of an IPSS/DBS model. These procedures can be referenced from within either Procedure facilities or Exogenous and Endogenous Service facilities.

Allocate and Release Schema Extent

The effect of executing the ALLOCATE SCHEMA EXTENT Built-in Procedure is to simulate the allocation of a logical address space to the Extent subfacility. The allocation causes the extent to take on a displacement attribute which specifies an extent's displacement relative to a Realm facility. This is done by placing the calculated displacement of the extent start into the Schema facility's extent table. The displacement is a real number whose dimension is in terms of the reference units implied for the Realm. This information is used by the Acquire and Store Record Built-in procedures (when the desired record type has an origin of Realm) to determine the logical address of a record type record. If a Schema Extent is not allocated, it is assumed to have zero capacity relative to its ability to hold schema records. The RELEASE SCHEMA EXTENT Built-in Procedure returns allocated address space.

Create and Destroy Schema

The purpose of the CREATE SCHEMA Built-in Procedure is to create for the named Schema facility a copy of its associated Data Structure component definition. The copy is placed in a pre-specified area of the Transaction Area. Once copied, simulation functions involving the Schema facility can occur. The definition remains valid until the issuing of a DESTROY SCHEMA statement for the same Schema facility.

Create and Modify Schema Extent

The function of the CREATE AND MODIFY SCHEMA EXTENT Built-in Procedure is to alter the current attributes associated with a Schema Extent facility. The attributes are first assigned via the extent definition in the Schema facility definition of the Data Structure component. A Schema Extent once created can only be modified or destroyed and can only be "created" when it is not in a "created" state. The procedure can only be invoked after a Schema has been created.

Destroy Schema Extent

The DESTROY SCHEMA EXTENT Built-in Procedure deletes the specified schema extent definition from the named Schema facility. The destroy procedure releases the logical address space associated with the extent by setting its size attribute to zero and removing its Realm facility association.

Modify Realm

The MODIFY REALM Built-in Procedure enables the modeler to change the attributes associated with a Realm facility definition. The procedure changes only those attributes explicitly identified in the Modify Realm Statement.

Modify Schema Record Type

The MODIFY SCHEMA RECORD TYPE Built-in Procedure changes the parameter values of the specified record type facility definition. The original parameters are specified in the RECORD TYPE statements in both the Data Structure and DBS components. The procedure changes only those attributes explicitly identified in the MODIFY RECORD TYPE statement.

Modify Schema Set

The function of the MODIFY SCHEMA SET Built-in Procedure is to modify parameters associated with a Schema's Set subfacility. The original parameters are specified via the Data Structure component's Set facility definition. The procedure changes only those attributes explicitly identified in the MODIFY SCHEMA SET Statement.

Cl.3.3 IPSS/DBS Data Structure Trversal Oriented Built-in Procedures

The following Built-in Procedures are used a) to create routes or individual sets based upon the schema definitions in the Data Structure component, b) to modify routes or sets by specifying the elements to be changed, and c) to destroy an existing route. The procedure statements are

divided into two groups, those which operate on sets, and those which operate on routes. These procedures can be referenced from within either Procedure facilities or Endogenous and Exogenous Service facilities.

Copy Set

The COPY SET Built-in Procedure creates a set occurrence from the specified set definition. The set occurrence is placed into an output array which the modeler designates. The set may either be copied from the corresponding set definition in the Data Structure component or from an existing route. In either case the modeler may change the set traversal currency indicator during the copy process.

Alter Set

The ALTER SET Built-in Procedure modifies a previously created set occurrence. The specified set occurrence may either be the output area of a previously executed Copy Set procedure or may exist within a route. This statement can be used to change the owner and member record types of the set as well as the record occurrences for these record types.

Find Member

The FIND MEMBER Built-in Procedure identifies a member record occurrence of a set relative to the specified set traversal currency. The output of this procedure is an updated member record occurrence and/or member record type of the set's traversal currency. A NO MEMBER EXIT parameter is provided for those cases in which it is not possible to update this currency; for example, FIND NEXT MEMBER when the number of occurrences associated with the member record type is zero. The FIND MEMBER examines the set and the record type facility definitions in the Data Structure component in order to determine the existence of the desired record. The statement types are as follows: a) FIND FIRST MEMBER specifies that the first member record occurrence for the designated set type is to be located and its identity placed in the set's traversal currency. By "identity" we mean the member record type, the

record occurrence within the set and the record occurrence within the record type. The success of this operation depends on a data structure which directly links the set traversal currency before execution of this procedure to the first record in the set; b) FIND PRIOR MEMBER specifies that the record to be identified is the record occurrence which immediately precedes the one specified by the set traversal currency. The success of this option likewise depends upon, a direct link between the record in the set traversal currency and the prior record occurrence; c) FIND NEXT MEMBER specifies that the record to be identified is the record occurrence which immediately follows the record occurrence of the set traversal currency. The successful completion of this procedure depends upon the existence of a direct link between the record specified by the set traversal currency and the next record occurrence in the set; d) FIND LAST MEMBER specifies that the last record occurrence for the specified set type is to be located and its identity placed in the set's traversal currency. The success of this option depends upon a data structure which directly links this record to the set traversal currency before execution of this procedure.

Identify Owner

The IDENTIFY OWNER Built-in Procedure identifies the record type which is the owner of the specified record type and identifies the set in which they both participate. This procedure examines the set and record type facility definitions in the Data Structure component of the current schema to determine this relationship. All the sets which own a given record type can be obtained by specifying IDENTIFY OWNER followed by successive IDENTIFY NEXT OWNER statements.

Identify Set Occurrences

The IDENTIFY SET OCCURRENCE Built-in Procedure identifies the sets which the current member record occurrence owns. By current member record occurrence we mean the traversal currency as reflected in the

specified set. This procedure examines the set and record type facility definitions in the Data Structure component of the current schema to determine ownership of a non-empty set of records. All of the sets owned by a record can be obtained by specifying IDENTIFY SET OCCURRENCE followed by successive IDENTIFY NEXT SET OCCURRENCE statements.

Create Route

The CREATE ROUTE Built-in Procedure determines the existence of a path between two record types of a schema. The beginning and ending nodes of the path are specified in terms of record type identifiers. The function of this procedure is to determine if these two record types are connected and if the connections meet all the set constraints. The set declarations of the identified schema are examined in the order of their declaration until either a path is found or all the declarations for the schema have been explained. All the paths between two record types can be located by specifying PATH = FIRST followed by successive PATH = NEXT statements. When a path is found, its sets are placed in the \$TRANS array beginning with the location returned as the value of the route identifier. If no path is found, the value of the route identifier will be a zero.

Destroy Route

The DESTROY ROUTE Built-in Procedure deletes the specified route from the \$TRANS array. This allows the re-use of the space occupied by the route for other modeling activities. The execution of this procedure does not affect schema definitions in the Data Structure component.

Modify Route

The MODIFY ROUTE Built-in Procedure changes the route definition by either: a) reestablishing the current set, b) deleting or inserting sets, or c) concatenating another route to the existing one. A route consists of a sequence of set traversal tables. This procedure can change this

sequence by deleting, inserting or concatenating set traversal tables relative to any of the sets in the route. Only one type of change can be specified for each Modify Route Statement.

Copy Route

The COPY ROUTE Built-in Procedure creates a route and copies the specified route data into it. The sets in the output route can be modified during this process. These modifications consist of: a) re-establishing the current set, b) inserting or deleting sets, or c) concatenating another route to the output route. These modifications can be specified either by naming specific sets or by referencing the sets in the route relative to the set currency.

Find Route

The FIND ROUTE Built-in Procedure uses a previously created route and the current schema definition to establish specific record occurrences within the sets of the route. A set other than the first set of the route may be designated as the beginning set for the procedure to establish record occurrences. A set other than the last set in the route may also be specified.

C2. IPSS PERFORMANCE MEASURES

The proposal identified 120 extensions to the currently available model statistics. The purpose of this section of Appendix C is to present the results of the research in the area of added statistical outputs from IPSS models.

The Information Processing System Simulator provides a modeler with a number of statistics concerning the behavior of modeler defined entities and IPSS supplied built-in services. These statistics fall into eight general categories listed below. The evaluative results of the research is found in Subsections C2.2 through C2.8 which correspond to these statistics categories:

- 1. Operational Statistics,
- 2. Request Stream Statistics,
- 3. I/O Activity.
- 4. Queueing Statistics,
- 5. Utilization Statistics,
- 6. Wait Statistics,
- 7. Service Statistics, and
- 8. Task/Activity Statistics.

In addition to these "automatic" statistics, the modeler can employ the complete facilities of the Fortran language to develop his own statistics. Statistics are printed automatically at the conclusion of each model simulation unless explicitly inhibited.

Time constraints prohibited the implementation of any new statistics. However, definitions for a number of the proposed statistics were completed and are reported in the appropriate subsection. These statistics employ both IPSS facility specific variables and those global to an entire model. For each defined statistic the variables required to support their collection and calculation are also discuss along with the method for data collection including the specific IPSS statements responsible for the activity. Note: all variable names used in these definitions are for publication purposes only.

C2.1 IPSS Operational Statistics

These statistics provide an overview of the operational characteristics of the current simulation execution. The Clock and Termination Statistics involve the length of the run, in terms of both simulation time and CPU time, and the reason for termination. The Transaction Statistics provide insight into the degree of transaction activity generated by the current simulation run. Statistics on the number of transactions generated and the average life span of a transaction are collected by transaction type. The results of their evaluation are shown in Table C2.1-1. Table C2.1-2 identifies the required data to meet the definitional requirements for these statistical extentions to the IPSS. The explicit statistic definitions are shown in Table C2.1-3.

Table C2.1-1 IPSS Operational Statistics

Statistic Type	Specific Statistic	Currently Available	Proposed Extention	Disposition*
Clock	Absolute elapsed simu- lation time Elapsed simulation time since last reset Elapsed simulation time of current model exe- cution Elapsed CPU time (IPSS run time)	x x x	X X	EU-DP
Termination	Scaling factor Number of errors Termination status	Х	x	EU-DP EU-DP
Transaction (By Trans- action type)	Total number generated Current number in system Maximum number in system Starting number in system Total number terminated Mean time of existence	X	X X X	EU-DP EU-DP EU-DP

*Distribution codes used in this table are defined as follows:

DU-DP: Examined, found useful - Defined, not implemented but used in paper models

Table C2.1-2 IPSS Operational Statistics - Required Model Statements

Variable Name	Description of the Variable	Method of data Collection
	VARIABLES GLOBAL TO AN IPSS	MODEL
CLOCK1	Absolute elapsed simulation time	Automatically by simulation driver
CLOCK2	Elapsed simulation time since last reset	Automatically by simulation driver
CLOCK3	Elapsed simulation time of current model execution	Automatically by simulation driver
CPUTIM	Beginning CPU time of current IPSS run	SVC call by simulation driver
SCALEF	Scaling factor	Specified by user
ERRM	Total number of errors in current model execution	Updated by internal IPSS error routine
TERMST	Termination status	Specified by user
	VARIABLES WHICH ARE TRANSACTION	SPECIFIC
TNGENA	Total number generated	Updated by simulation driver upon generation of new transaction
TNCURA	Current number in system	Maintained by driver in SYSINF
TNMAXA	Maximum number in system	Either user specified or system default
TNSTRTA	Starting number in system	Either user specified or system default
TNTRMA	Total number terminated	Updated by driver upon termination

Table C2.1-3 IPSS Operational Statistics - Definition

Stastic Type	Specific Statistic Name	Definition
Clock	Absolute elapsed simulation time	CLOCK1
	Elapsed simulation time since last reset	CLOCK2
	Elapsed simulation time of current model execution	CLOCK3
	*Elapsed CPU time (IPSS run time)	CPUTIM
	*Scaling Factor	SCALEF
Termination	Number of errors	ERRM
	*Termination Status	TERMST
Transaction	*Total number generated	TNGENA
(by trans- action type)	Current number in system	TNCURA
action type,	Maximum number in system	TNMAXA
	Starting number in system	TNSTRTA
	*Total number terminated	TNTRMA
	*Mean time of existence	(CLOCK3/TNGEN _A)

^{*}Statistic currently exists in IPSS models.

C2.2 IPSS Request Stream Statistics

These statistics describe the behavioral characteristics of the Request Stream facilities. For each Exogenous Event a profile of the interarrival time mechanism is presented. The event of the IPSS built-in Endogenous Services, ENABLE/DISABLE Event, on the arrival of events is also summarized. The collection of these statistics is an automatic by-product of the simulation drive mechanism. Table C2.2-1 presents the evaluation results. Table C2.2-2 identifies the required model data and Table C2.2-3 the actual statistic definition.

Table C2.2-1 IPSS Request Stream Statistics

Statistic Type	Specific Statistic	Currently Available	Proposed Extension	Disposition'
General	Number generated Number terminated	х	X	EU-DP
Arrival Time	First generated arrival		X	EU-DP
	Average interarrival time Last arrival time		х	EU-DP
	generated		Х	EU-DP
Enable/Disable	Number of enables		X	EU-DP
	Total time enabled		X	EU-DP EU-DP
	Average enable time Number of disables		Ŷ	EU-DP
	Total Time Disabled		Ŷ	EU-DP
	Average disable time		x	EU-DP

*Disposition Codes used in this table are defined as follows:

EU-DP: Examined, found useful - Defined, not implemented but useful in paper models

Table C2.2-2

IPSS Request Stream Statistics - Required Model Statements

Variable Name	Description of the Variable	Method of Data Collection
EXGEN	Number generated	Updated by driver
EXTRM	Number terminated	Updated by driver
EFARR	First generated arrival time	Specified by user (either directly or via Procedure)
ELARR	Last (latest) arrival time generated	Updated by driver
ENEN	Number of enables	Updated by ENABLE statement
ETLEN	Time of last enable/disable	Updated by both the ENABLE and DISABLE statements
ETTEN	Total time enabled	Calculated by DISABLE statement as (CLOCK3 - ETLEN) + ETTEN
ENDSA	Number of disables	Updated by DISABLE statement
ETTDSA	Total time disables	Calculated by ENABLE statement as (CLOCK3 - ETLEN) + ETTDSA

Table C.2.2-3

IPSS Request Stream Statistics - Definitions

Statistic Type	Specific Statistic Name	Definition
General	Number Exogenous requests generated *Number Exogenous requests terminated	EXGEN EXTRM
Arrival time	*First generated Exogenous event arrival time	EFARR
	*Mean interarrival time *Last arrival time	(ELARR - EFARR)/EXGEN ELARR
Enable/Disable	*Number of enables *Total time enabled *Mean enable time *Number of disables *Total time disabled *Mean disable time	ENEN ETTEN (ETTEN/EXEN) ENDSA ETTDSA (ETTDSA/ENDSA)

^{*}Statistic currently exists in IPSS models

C.2.3 IPSS I/O Activity Statistics

All of the proposed new statistics in the area of I/O activity were concerned with the concept of locality. Statistics were proposed at the Data Set level and within the Data Set, by File, Extent, Volume, Cylinder and Track, and at the Area level, and within Area by Segment and by Volume. All of these statistics are to be gathered by the GET ADDRESS statement. Major modifications however, will be required to the existing IPSS code supporting the IPSS GET ADDRESS Built-in Procedure in order to gather the required data. The purpose of these statistics is to measure the degree of hardware and file management level I/O activity that occurred during the current model. Measures of I/O routine performance are collected by I/O routine type, by associated access mechanism, and by referenced Data Set and File. Statistics showing the degree of locality of consecutive logical and physical references by Data Set and Area were also proposed.

A second set of statistics of similar intent was proposed for the IPSS/DBS subsystem. Since these statistics are closely associated with other I/O activities, the decision was made to delay the examination of this class of statistics until the IPSS/DBS subsystem was operational. However, the I/O-related statistics are defined in Table C2.3 for report completeness.

Table C.2:3 IPSS I/O Activity Statistics

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition
I/O Invocations	Number of calls	X		
(By I/O Statement)	Percent zero-time calls	X		
	Mean time per call	X		
	Mean time for non-zero			
	calls	X		
	Standard deviation of			
	time per call	X		
	Standard deviation of			
	time per call			
	(excluding zero-time			
	calls)	X		
	(Mean and standard			
	deviation of various			
	physical attributes,			
	i.e., cylinders crossed			
	characters read, etc.)	X		
I/O Invocations	Du 1/0 tune for each			
	By I/O type for each			
(By Access Mechanism)	Access Mechanism: Number of calls	v		
	Number of zero-time	X		
	calls	v		
	Percent zero-time calls	X X		
	Mean routine time	x		
	Mean routine time (ex-	^		
	cluding zero-time			
	calls)	X		
	Standard deviation of			
	mean routine time	X		
	Standard deviation of	, and the second		
	mean routine time (ex-			
	cluding zero time call	s) X		
I/O Invocations	By I/O type for each Data	a		
(By Data Set and	Set and File:			
File)	Number of calls	X		
	Number of zero-time call:			
	Percent of zero-time cal	ls X		
	Mean routine time	X		
	Mean routine time (ex-			
	cluding zero-time call	s) X		
	Standard deviation of			
	mean routine time	X		
	Standard deviation of			
	mean routine time (ex-			
	cluding zero-time call	s) X		

Table C.2.3 IPSS I/O Activity Statistics (Continued)

		Currently	Proposed	
Statistics Type	Specific Statistic	Available	Extension	Disposition
Locality (Data Set) (Conditional)	By Data Set: Number of consecutive references for each			
	file By File: Number of consecutive references for each		X	NE
	Extent By Extent: Number of consecutive references to the same		X	NE
	Volume By Volume: Number of consecutive references to the same		X	NE.
	cylinder Average number of		X	NE
	cylinders crossed By Cylinder: Number of consecutive references to the same		X	NE
	track Average number of track		X	NE
	crossed By Track: Number of consecutive references to the	3	X	NF
	same block Average number of		X	NE
	blocks crossed		X	NE
Locality (Area)	By Area: Total number of referen	CPS		
(Conditional)	to area Number of consecutive r		X	NE
	ences to same area Average number of conse		X	NE
	references Average length between consecutive references		X	NE
	terms of reference unit By Segment: Total number of referen	s	X	NE
	to segment		X	NF

Table C.2.3 IPSS I/O Activity Statistics (Continued)

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition
	Number of consecutive			
	references to same			
	segment		X	NE
	Average number of			
	consecutive references		X	NE
	Average length between			
	consecutive references			
	in terms of reference			
	units		X	NE
	By Volume:			
	Number of consecutive			
	references to same			
	volume		X	NE
	Average number of			
	consecutive reference	es	X	NE
	Average length between			
	consecutive referenc	es		
	in terms		X	NE

^{*}Disposition codes used in this table are:

NE: Not evaluated

C.2.4 IPSS Qeueing Statistics

An extensive number of queueing statistics are currently gathered during an IPSS model's execution. Several additional extensions to these statistics were proposed and evaluated. These are now being incorporated into the IPSS statistics generation procedures. Table C.2.4-1 identifies both the existing and proposed statistics and the result of the evaluative effort. The purpose of Table C.2.4-2 is to identify the specific data for these statistics. Table C.2.4-3 provides definitions for the statistics.

Table C.2.4-1 IPSS Queueing Statistics

Statistics Type		Currently Available	Proposed Extension	Disposition
Busy Period	Total number of busy			
busy . Cr rou	periods	X		
	Number of zero time busy			
	periods	Χ		
	Percent zero time busy			
	periods	, X		
	Total busy time	X		
	Percent busy	X		
	Mean length of busy	^		
		X		
	period			
	Mean length of busy perio	u		
	(excluding zero time		X	EU-DP
	periods)		^	EU-DP
	Standard deviation for			
	non-zero time busy		X	EU-DP
	periods		^	בט-טף
	Maximum length of a		v	בון סס
	busy period		X	EU-DP
	Minimum (non-zero) length			511 DD
	of a busy period		X	EU-DP
	Length of the current			
	busy period		X	EU-DP
Idle Period	Total number of idle			
	periods	X		
	Number of zero-time			
	idle periods	X		
	Percent zero-time			
	idle periods	X		
	Total idle time	X		
	Percent idle	X		
	Mean length of an idle			
	period	X		
	Mean length of an			
	idle period (excluding			
	zero~time periods)		X	EU-DP
	Standard deviation for		^	20 0.
	non-zero time idle			
	periods		X	EU-DP
	Maximum length of an		^	LU-UF
			X	EU-DP
	idle period		^	LU-UF
	Minimum (non-zero)			
	length of an		v	EU-DP
	idle period		X	EU-UP

Table C.2.4-1 IPSS Queueing Statistics (Continued)

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition
Statistics Type	Specific Statistic	Available	LACEISTOIL	DISPOSTETOR
Queue Length	Mean length	X		
(During a Busy	Maximum length	X		
Period)	Current length	X		
	Standard deviation			
	of mean length	X		
Queue	Total number of			
	queue entries	X		
	Total number of			
	completed queue			
	entries	X		
	Number of zero-time	^		
	queue entries	X		
	Percent of zero-time	^		
	queue entries	X		
	Mean queue transit	^		
	time	X		
		۸		
	Mean queue transit			
	(excluding zero-	v		
	time transit)	X		
	Standard deviation of			
	non-zero time	· ·		
	queue transit	X		
	Maximum transit time		X	EU-DP
	Minimum (non-zero)			
	transit time		X	EU-DP

Disposition codes used for this Table are:

EU-DP: Examined, found useful - Defined not implemented but useful in paper model

Table C.2.4-2 IPSS Queueing Statistics - Required Model Statements

Variable Name*	Description of the Variable	Method of Data Collection
QCNCE	Current number of concurrent elements in queue	Update by Queue and Depart Queue Statement
QMXNCE	Maximum number of current concurrent elements in queue	Update by Queue statement Queue
QSNCE	Sum of number of concurrent elements in queue	Updated by Depart Queue statement
QSNCES	Sum of (number of concurrent elements in queue)**2	Updated by Queue and Depart Queue statement
QSQE	Total number of queue elements	Update by Queue statement
QSZQE	Total number of zero-time queue elements	Updated by Depart Queue statements
QSTT	Sum of time per transaction	Updated by Depart Queue statement
QSTTS	Sum of (time per trans)**2	Updated by Depart Queue statement
QCTI	Cumulative Time integral	Updated by Queue/Depart Queue statement
QTQDQ	Time of last occurrence of Queue or Depart Queue invocation	Updated by Queue and Depart Queue statement
QSBP	Total number of busy periods	Updated by Queue statement
QSBPL	Sum of length of busy periods	Updated by Depart Queue statement
QSBPLS	Sum of (length of busy periods)**2	Updated by Depart Queue statement
QSIP	Number of idle periods	Updated by Depart Queue statement
QSIPL	Sum of length of time for each idle period	Update by Queue statement

Table C.2.4-2 IPSS Queueing Statistics - Required Model Statements (Continued)

Variable Name*	Description of the Variable	Method of Data Collection
QSIPLS	Sum of (length of idle periods)**2	Updated by Queue statement
QTBIP	Time of last Busy or Idle period start	Updated by Queue statement and Depart Queue
QSZBP	Number of zerotime busy periods	Updated by Queue statement
QSZIP	Number of zerotime idle periods	Updated by Depart Queue statement
QMNBPL	Minimum non-zero busy period length	Updated by Depart Queue statement
QMXBPL	Maximum busy period length	Updated by Depart Queue statement
QCBPL	Length of current busy period	Updated by Depart Queue statement
QMNIPL	Minimum non-zero idle period length	Updated by Queue statement
QMXIPL	Maximum idle period length	Updated by Queue statement
QCIPL	Length of current idle period	Updated by Queue statement
QMXTT	Maximum transit time per transaction	Updated by Depart Queue statement
QMNTT	Minimum non-zero transit time per transaction	Updated by Depart Queue statement

 $[\]star$ All data are stored in the statistics area associated with the designated facility.

Table C.2.4-3 IPSS Queueing Statistics - Definitions

Statistic Type	Specific Statistic	Definition
Busy Period	Total number of busy periods	QSBP
	Number of zerotime busy periods	QSZBP
	%zero time busy periods	(QSZBP/QSBP)*100
	Total Busy time	QSBPL
	%busy	(QSBPL/CLOCK3)*100
	Mean length of busy period	(QSBPL/QSBP)
	Standard deviation of mean length	((QSBP*QSBPLS~QSBPL**2)/ (QSBP*(QSBP~1)))**0.5
	Mean length of nonzero time busy periods	(QSBPL/(QSBP-QSZBP))
	Standard deviation of nonzero time busy periods	<pre>(((QSBP-QSZBP)*QSBPLS-QSBPL**2) /((QSBP-QSZBP)*(QSBP-QSZBP-1)) **0.5</pre>
	Maximum length of a busy period	QMXBPL
	Minimum non-zero length	QMNBPL
	Length of current busy period	QCBPL
[d]e Period	Total number of idle periods	QSIP
	Number of zerotime idle periods	QSZIP
	%zero time idle periods	(QSZIP/QSIP)*100
	Total idle time	QSIPL
	%idle	(QSIPL/CLOCK3)*100
	Mean length of idle period	(QSIPL/QSIP)
	Standard deviation of mean length	(((QSIP-QSIPLS)-QSIPL**2) /(QSIP*(QSIP-1)))**0.5

Detroit Marie

Total Control

Table C.2.4-3 IPSS Queueing Statistics - Definitions (Continued

Statistic Type	Specific Statistic	Definition
	Mean length of nonzero time idle periods	(QSIPL/(QSIP-QSZIP))
	Standard deviation of nonzero time idle periods	(((QSIP-QSZIP)*QSIPLS-QSIPL**2) /((QSIP-QSZIP)*(QSIP-QSZIP-1)))**0.5
	Maximum length of an idle period	QMXIPL
	Minimum nonzero length	QMNIPL
	Length of current idle period	QCIPL
Queue Length	Mean length	(QSNCE/QSQE)
	Maximum length	QMXNCE
	Current length	QCNCE
	Standard deviation of mean length	((QSQE*QSNCES-QSNCE**2) /(QSQE*(QSQE-1)))**0.5
Queue	Total number of queue entries	QSQE
	Total number of completed queue entries	QSQE - QCNCE = QSCQE
	Number of zerotime queue entries	QSZQE
	%zerotime queue entries	(QSZQE/QSQE)*100
Transit Time	Mean transit time per transaction	(QSTT/(QSQE-QCNCE))
	Mean nonzero time per transaction	(QSTT/(QSCQE-QSZQE))
	Standard deviation of mean transaction time	((QSCQE*QSTTS-QSTT**2) /(QSCQE*(QSCQE-1)))**0.5

Table C.2.4-3 IPSS Queueing Statistics - Definitions (Continued)

Statistic Type	Specific Statistic	Definition
	Standard deviation of mean non-zerotime transactions	<pre>(((QSCQUE-QSZQE)*QSTTS-QSTT**2) /((QSCQE-QSZQE)*(QSCQE-QSZQE-1)))**0.5</pre>
	Maximum transit time	QMXTT
	Minimum nonzero transit time	QMNTT

C.2.5 IPSS Utilization Statistics

An IPSS model generates utilization statistics on user defined modeled facilities. Several extensions to these statistics were proposed and evaluated in the research. Table C.2.5-1 identifies all IPSS utilization statistics, whether current or proposed and their desposition after evaluation. Table C.2.5-2 identifies the means for their generation and Table C.2.5-3 contains their definitions.

Table C.2.5-1 IPSS Utilization Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition
Busy Period	Total number of busy periods Number of zero time busy	X		
	periods Percent zero time busy periods	X X		
	Total busy time	x		
	Percent busy	x		
	Mean length of busy period Mean length of busy period	χ̈́		
	<pre>(excluding zero time periods) Standard deviation for</pre>		X	EU-DP
	nonzero time busy periods Maximum length of a busy		X	EU-DP
	period Minimum (non-zero) length		X	EU-DP
	of a busy period Length of the current busy		X	EU-DP
	period		X	EU-DP
dle Period	Total number of idle periods Number of zero time idle	X		
	periods Percent zero-time idle	X		
	periods	X		
	Total idle time	x		
	Percent idle Mean length of an idle	x		
	period	X		
	<pre>Mean length of an idle period (excluding zero-</pre>			
	time periods)		X	EU-DP
	Standard deviation for nonzero time idle periods		X	EU-DP
	Maximum length of an idle period		X	EU-DP
	Minimum (non-zero) length of an idle period		X	EU-DP
	Length of the current idle period		X	EU-DP
Concurrency	Mean level of concurrency	X		
	Maximum level of concurrency	X		
	Current level of concurrency	X		
	Standard deviation of mean			
	level of concurrency	X		

Table C.2.5-1 IPSS Utilization Statistics (Continued)

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition
Seizure	Total number of seizures	X		
	Total number of completed seizures	X		
	Number of zero-time	^		
	seizures	X		
	Percent zero-time			
	seizures	X		
	Mean seizure time	X		
	Mean seizure time (ex- cluding zero-time			
	seizures)	X		
	Standard deviation of non-			
	zero time seizures	X		
	Maximum length of a			
	seizure		X	EU-DP
	Minimum (non-zero) length			
	of a seizure		X	EU-DP

Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, not implemented but useful in paper model.

Table C.2.5-2 IPSS Utilization Statistics - Required Model Statement

Variable Name*	Description of the Variable	Method of Data Collection
	bescription of the variable	Method of Data Coffection
UCNCS	Current number of Concurrent seizures	Updated by Seize/Release statements
UMXNCS	Maximum number of current concurrent seizures	Updated by Seize statement
USNCS	Sum of number of concurrent seizures	Updated by Seize/Release statement
USNCSS	Sum of (number of concurrent seizures)**2	Updated by Seize/Release statement
USSEZ	Total number of seizures	Updated by Seize statement
USZSEZ	Total number of zero-time seizures	Updated by Release statement
USTT	Sum of time per transaction	Updated by Release statement
USTTS	Sum of (time per trans)**2	Updated by Release statement
UCTI	Cumulative Time Integral	Update by Seize and Release statements
UTSR	Time of last seize or release	Updated by Seize and Release statements
USBP	Total number of busy periods	Update by Seize statement
USBPL	Sum of length of busy periods	Update by Release statement
USBPLS	Sum of (length of busy periods)**2	Updated by Release statement
USIP	Number of idle periods	Updated by Release statement
USIPL	Sum of length of idle periods	Updated by Seize statement
USIPLS	Sum of (length of idle periods)**2	Updated by Seize statement
UTBIP	Time of last Busy or idle period start	Updated by Seize and Release statement

Table C.2.5-2 IPSS Utilization Statistics - Required Model Statement (Continued)

Variable Name*	Description of the Variable	Method of Data Collection
USZBP	Number of zerotime busy periods	Updated by Release statement
USZIP	Number of zerotime idle periods	Updated by Seize statement
UMNBPL	Minimum non-zero busy period length	Updated by Release statement
UMXBPL	Maximum busy period length	Updated by Release statement
UCBPL	Length of current busy period	Updated by Release statement
UMNIPL	Minimum non-zero idle period length	Updated by Seize statement
UMXIPL	Maximum idle period length	Updated by Seize statement
UCIPL	Length of current idle period	Updated by Seize statement
UMXTT	Maximum transit time per transaction	Updated by Release statement
UMNTT	Minimum non-zero transit time per transaction	Updated by release statement

^{*}All data are stored in the statistics area associated with the designated facility.

Table C.2.5-3 IPSS Utilization Statistics - Definitions

Statistic Type	Specific Statistic	Definition
Busy Period	Total number of busy periods	USBP
	Number of zerotime busy periods	USZBP
	%zero time busy periods	(USZBP/USBP)*100
	Total busy time	USBPL
	%busy	(USBPL/CLOCK3)*100
	Mean length of busy period	(USBPL/USBP)
	Standard deviation of mean length	(((USBP-USBPLS)-USBPL**2) /(USBP*(USBP-1)))**0.5
	Mean length of non-zero time busy periods	(USBPL/(USBP-USZBP))
	Standard deviation of non-zero time busy periods	(((USBP-USZBP)*USBPLS-USBPL**2) /((USBP-USZBP)*(USBP- USZBP-1)))**0.5
	Maximum length of a busy period	UMXBPL
	Minimum non-zero length	UMNBPL
	Length of current busy period	UCBPL
Idle Period	Total number of idle periods	USIP
	Number of zero time idle periods	USZIP
	%zero time idle periods	(USZIP/USIP)*100
	Total idle time	USIPL
	%idle	(USIPL/CLOCK3)*100
	Mean length of idle period	(USIPL/USIP)
	Standard deviation of mean length	((USIP*USIPLS-USIPL**2) /(USIP*(USIP-1)))**0.5

Table C.2.5-3 IPSS Utilization Statistics - Definitions (Continued)

Statistic Type	Specific Statistic	Definition
	Mean length of non-zero time idle periods	(USIPL/(USIP-USZIP))
	Standard deviation of non-zero time idle periods	(((USIP-USZIP)*USIPLS-USIPL**2) /((USIP-USZIP)*(USIP- USZIP-1)))**0.5
	Maximum length of an idle period	UMXIPL
	Minimum non-zero length	UMNIPL
	Length of current Idle period	UCIPL
Concurrency	Mean level	(USNCS/USSEZ)
	Maximum level	UMXNCS
	Current level	UCNCS
	Standard deviation of mean level	((USSEZ*USNCSS-USNCS**2) /(USSEZ*(USSEZ-1)))**0.5
Seizures	Total number of seizures	USSEZ
	Total number of completed seizures	USCSEZ = USSEZ - UCNCS
	Number of zero-time seizures	USZSEZ
	%zero-time seizures	(USZSEZ/USSEZ)*100
Transit Time	Mean transit time per trans- action	(USTT)/USCSEZ
	Mean non-zero time per transaction	(USTT/(USCSEZ-USZSEZ))
	Standard deviation of mean transaction time	((USCSEZ*USTTS-USTT**2) /(USCSEZ*(USCSEZ-1)))**0.5
	Standard deviation of mean non- zero-time transactions	(((USCSEZ-USZSEZ)*USTTS-USTT**2) /((USCSEZ-USZSEZ)*(USCSEZ- USZSEZ-1)))**0.5

Table C.2.5-3 IPSS Utilization Statistics - Definitions (Continued

Statistic Type	Specific Statistic	Definition	
	Maximum transit time	UMXTT	
	Minimum non-zero transit time	UMNTT	

C.2.6 IPSS Wait Statistics

These statistics measure the degree to which individual transactions were delayed until specific system resources attained appropriate status. They are similar to types of statistics gathered for Queueing Statistics. IPSS automatically collects these statistics as a by-product of the following IPSS Built-in procedures:

- 1. Wait Facility Status,
- 2. Wait Service Complete,
- 3. Wait I/O Complete,
- 4. Wait Semaphore,
- 5. Wait Access Path, and
- 6. Wait Signal.

The effect of these built-in procedures is to suspend the execution of a service until the identified condition is met. The statistics for each are the same. Only the WAIT FACILITY STATUS definitions are described in this section. Table C2.6-1 identifies all the IPSS Wait Statistics, Table C2.6-2 the means for their operation and Table C2.6-3 their definitions.

Table C.2.6-1 IPSS Wait Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition
Busy Period	Total number of busy periods Number of zero-time busy periods Percent zero-time busy periods Total busy time Percent busy Mean length of a busy period Mean length of a busy period (excluding zero-time periods) Standard deviation for non- zero time busy periods Maximum length of a busy period	X X X X X X	X X X	EU-DP EU-DP EU-DP
	Minimum (non-zero) length of a busy period Length of the current busy period		X X	EU-DP EU-DP
Idle Period	Total number of idle periods Number of zero-time idle periods Percent zero-time idle periods Total idle time Percent idle Mean length of an idle period (excluding zero-time periods) Standard deviation for non- zero time idle periods Maximum length of an idle period Minimum (non-zero) length of an idle period Length of the current idle period		X X X X	EU-DP EU-DP EU-DP EU-DP
Wait Chain Length Characteristics	Mean length Maximum length Current length Standard deviation of mean length	X X X		

Table C.2.6-1 IPSS Wait Statistics (Continued)

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition
Wait Chain Entry	Total number of entries	X		
Characteristics	Total number of completed entries	X		
	Number of zero-time	^		
	entries	X		
	Percent zero-time entries	X		
	Mean time in chain Mean time in chain (ex- cluding zero-time	X		
	entries)	X		
	Standard deviation of mean time non-zero-time			
	entries	X		
	Maximum time in chain Minimum (non-zero) time		X	EU-DP
	in chain		X	EU-DP

Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, not implemented but useful in paper model.

Table C.2.6-2 IPSS Wait Statistics - Required Model Statements

Variable Type*	Description of the Variable	Method of Data Collection
WCNE	Current number of Concurrent wait chain elements	
WMXNCE	Maximum number of current concurrent wait chain elements	
WSNCE	Sum of number of concurrent wait chain elements	
WSNCES	Sum of (number of concurrent wait chain elements)**2	
WSWCE	Total number of wait chain elements	
WSZWCE	Total number of zero-time wait chain elements	A11 -1-1/-1/-
WSTT	Sum of time per transaction	All statistics are maintained by the Wait statements
WSTTS	Sum of (time per trans)**2	
WCTI	Cumulative Time integral	
WTWFSS	Time of last wait start or wait end	
WSBP	Total number of busy periods	
WSBPL	Sum of length of busy periods	
WSBPLS	Sum of (length of husy periods)**2	
WSIP	Number of idle periods	
WSIPL	Sum of length of idle periods	
WSIPLS	Sum of (length of idle periods)**2	
WTBIP	Time of last Busy or Idle period start	

Table C.2.6-2 IPSS Wait Statistics - Required Model Statements (Continued)

Variable Type*	Description of the Variable	Method of Data Collection
WSZBP	Number of zerotime busy periods	
WSZDP	Number of zerotime idle periods	
WMNBPL	Minimum non-zero busy period length	
WMXBPL	Maximum busy period length	
WCBPL	Length of current busy period	All statistics are maintained
WMNIPL	Minimum non-zero idle period length	by the Wait Statement
WMXIPL	Maximum idle period length	
WCIPL	Length of current idle period	
WMXTT	Maximum transit time per transaction	
WMNTT	Minimum non-zero transit time per transaction	

^{*}All are stored in the statistics area associated with the designated facility.

Table C.2.6-3 IPSS Wait Statistics - Definitions

Statistics Type	Specific Statistic	Definitions
Busy Period	Total number of busy periods	WSBP
	Number of zerotime husy periods	WSZBP
	%zero time busy periods	(WSZBP/WSBP)*100
	Total busy time	WSBPL
	%busy	(WSBPL/CLOCK3)*100
	Mean length of busy period	(WSBPL/WSBP)
	Standard deviation of mean length	((WSBP*WSBPLS-WSBPL**2) /(WSBP*(WSBP-1)))**0.5
	Mean length of non-zero time busy periods	(WSBPL/(WSBP-WSZBP)
	Standard deviation of non-zero time busy periods	(((WSBP-WSZBP)*WSBPLS- WSBPL**2)/((WSBP-WSZBP) *(WSBP-WSZBP-1)))**0.5
	Maximum length of a busy period	WMXBPL
	Minimum non-zero length	WMNBPL
	Length of current busy period	WCBPL
Idle Period	Total number of idle periods	WSIP
	Number of zerotime idle periods	WSZIP
	%Zerotime idle periods	(WSZIP/WSIP)*100
	Total idle time	WSIPL
	%Idle	(WSIPL/CLOCK3)*100
	Mean length of idle period	(WSIPL/WSIP)
	Standard deviation of mean length	((WSIP*WSIPLS-WSIPL**2) /(WSIP*(WSIP-1)))**0.5
	Mean length of non-zero time idle periods	(WSIPL/(WSIP-WSZIP))

Table C.2.6-3 IPSS Wait Statistics - Definitions (Continued)

Statistics Type	Specific Statistic	Definitions
	Standard deviation of non-zero time idle periods	(((WSIP-WSZIP)*WSIPLS-WSIPL**2 /((WSIP-WSZIP)*(WSIP-WSZIP-1))**0.5
	Maximum length of an idle period	WMXIPL
	Minimum non-zero length	WXNIPL
	Length of current Idle period	WCIPL
Chain Length	Mean length	(WSNCE/WSWCE)
	Maximum length	WMXNCE
	Current length	WCNCE
	Standard deviation of mean length	((WSWCE*WSNCES-WSNCE**2) /(WSWCE*(WSWCE-1)))**0.5
Chain	Total number of wait chain entries	WSWCE
	Total number of completed chain entries	WSCCE = WSWCE - WCNCE
	Number of zero-time entries	WSZWCE
	%zero-time chain entries	(WSZWCE/WSWCE)*100
Transit Time	Mean transit time per trans- action	(WSTT/WSCCE)*100
	Mean non-zero time per transaction	(WSTT/(WSCCE-WSZWCF))
	Standard deviation of mean transaction time	((WSCCE*WSTTS-WSTT**2) /(WSCCE*(WSCCE-1)))**0.5
	Standard deviation of mean non- zero-time transactions	(((WSCCE-WSZWCE)*WSTTS- WSTT**2)/((WSCCE-WSZWCE) *(WSCCE-WSZWCE-1)))**0.5
	Maximum transit time	QMXTT
	Minimum non-zero transit time	WMNTT

C.2.7 IPSS Service Related Statistics

IPSS Exogenous and Endogenous services are modeler-defined model components representing information system processes. These processes could be manual or automated. IPSS provides the modeler with over 100 statements in addition to the use of all of ANSI Fortran to define these services. In the current IPSS, services are treated the same as other acquirable resources for statistic gathering purposes.

In the proposal a number of additional statistics specific to services were proposed. The purpose of these new statistics is to characterize service performance in a more meaningful manner. This is accomplished in the following two ways:

A. Modeler Explicit Statistics

Three new sets of IPSS statements permit the modeler to divide overall service time into three segments;

- Queue Time time which a request for service spends in a service maintained queue waiting for its service to commence.
- Acquire Time time spent by the service acquiring the resources needed to service a request, and
- Service Time time spent by the service servicing the request.

These statistics are generated only when explicitly requested by the modeler through the employment respectively of the following IPSS statement pairs:

- 1. MARK QUEUE START MARK QUEUE END
- 2. MARK ACQUIRE START MARK ACQUIRE END
- 3. MARK SERVICE START MARK SERVICE END

B. Modeler Implicit Statistics

The purpose of these statistics is to allocate a service's time to subordinate-services. The IPSS translator will identify

all Endogenous services (modeler defined or IPSS Built-in) requested by a modeler defined service. Statistics on subservice performance will be gathered and displayed with the calling service. These statistics show total service time for the invoked service with respect to the invocations by the requestor service and total time. The requestor service was delayed wait for the requested service to complete processing. These statistics are automatically generated by the REQUEST SERVICE - WAIT SERVICE statements.

Table C.2.7-1 identifies the specific statistics generated. Tables C.2.7-2 and C.2.7-3 specify the IPSS mechanisms used in their generation.

Table C.2.7-1 IPSS Service Related Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition*
Execution	Total number of executions		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean execution time		X	EU-DP
	Mean non-zero execution time		X	EU-DP
Queue Time	Total number of periods		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean period length		X	EU-DP
	Mean non-zero period length Mean percent of execution		X	EU-DP
	time		X	EU-DP
	Mean non-zero percent of execution time		X	EU-DP
Acquire Time	Total number of periods		X	EU-DP
Acquire Time	Percent zero-time		x	EU-DP
	Mean period length		â	EU-DP
	Mean non-zero period		^	LU-UF
	length		X	EU-DP
	Mean percent of execution			20 0.
	time		X	EU-DP
	Mean non-zero percent of			
	execution time		X	EU-DP
Service Time	Total number of periods		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean period length Mean non-zero period		X	EU-DP
	length Mean percent of execution		X	EU-DP
	time Mean non-zero percent of		X	EU-DP
	execution time		X	EU-DP
Service by Service	Total number of invocations		X	EU-DP
(For each	Percent zero-time invocations	3	X	EU-DP
Service invoked	Mean time per invocation		X	EU-DP
by a parent	Mean time per invocation			
service)	<pre>(excluding zero-time invocations)</pre>		X	EU-DP
	Total number of Wait			
	Service Completes		X	EU-DP

Table C.2.7-1 IPSS Service Related Statistics (Continued)

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition*
	Demont Thus time Waits		Y	EU-DP
	Percent zero-time Waits Mean time per Wait		x	EU-DP
	Mean time per Wait (excluding zero-time Waits)		X	EU-DP

*Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, but not implemented

Table C.2.7-2 IPSS Service Related Statistics - Required Model Statistics

Variable Type	Description of the Variable	Method of Data Collection
The following area associa	ng variables identify storage locat ated with each modeler defined Exog	ions in the fixed statistics enous and Endogenous Service
SNT	Total number of executions	Updated by REQUEST SERVICE Statement
SNZT	Total number of zero-time executions	Updated by WAIT RETURN* statement
SCUMT	Cumulative execution time	Update by TERMINATE* statement
SNQ	Number of queue periods	Updated by MARK QUEUE START statement
SNZQ	Number of zero-time queue periods	Updated by MARK QUEUE END,
SCUMQ	Cumulative queue period time	MARK ACQUIRE START, MARK SERVICE START or TERMINATE statements
SNAC	Number of acquire periods	Updated by MARK ACQUIRE START statement
SNZAC	Number of zero time acquire periods	Updated by MARK ACQUIRE END, MARK SERVICE START or TERMINATE
SCUMAC	Cumulative acquire period time	statements
SNS	Number of service periods	Updated by MARK SERVICE START statement
SNZS	Number of zero time periods	Updated by MARK SERVICE END
SCUMS	Cumulative service period time	or TERMINATE statements
the statist	ng variables identify storage locat ics area. These are replaced for e y a service. (Note: intermediate s created for a requested service u)	ach subordinate service statistics are kept in the

SSNIA

Total number of invocations of Service A

Updated by REQUEST SERVICE statement

Table C.2.7-2 IPSS Service Related Statistics - Required Model Statistics (Continued)

Variable Type	Description of the Variable	Method of Data Collection
SSNZIA	Total number of zero-time invocations of Service A	Update by TERMINATE statement
SSTMIA	Total execution time	Update by TERMINATE statement
SSNWA	Total number of Wait Service	Update by WAIT SERVICE Complete statement
SSNZWA	Total number of zerotime	Updated by WAIT RETURN statement
SSTWSA	Total Wait service time	Updated by WAIT RETURN statement

^{*}These statements are internal to IPSS and not available to the modeler.

Table C.2.7-3 IPSS Service Related Statistics - Definitions

Statistics Type	Specific Statistic	Definitions
Execution	Total number of executions	SNT
	Percent zero-time	(SNZT/SNT)*100
	Mean execution time	SAVT = (SCUMT/SNT)
	Mean non-zero time execution time	SAVTX = (SCUMT/(SNT - SNZT))
Queue Time	Total number of queue periods	SNQ
	Percent zero-time	(SNZQ/SNQ)*100
	Mean period length	SAVQ = (SCUMQ/SNQ)
	Mean non-zero time period length	SAVZQ = (SCUMQ/(SNQ - SNZQ))
	Mean percent of execution time	SAVQ/SAVT
	Mean non-zero % of non-zero execution time	SAVQZ/SAVTX
Acquire Time	Total number	SNAC
	Percent zero-time	(SNZAC/SNAC)*100
	Mean length	SAVAC = (SCNMAC/SNAC)
	Mean non-zero length	SAVACX = (SCUMAC/(SNAC - SNZAC))
	Mean % of execution time	(SAVAC/SAVT)*100
	Mean non-zero % of non-zero execution time	(SAVACX/SAVTX)*100
Service Time	Total number	SNS
	Percent zero-time	(SNZS/SNS)*100
	Mean length	SAVS = (SCUMS/SNS)

Table C.2.7-3 IPSS Service Related Statistics - Definitions (Continued)

Statistics Type	Specific Statistic	Definitions
	Mean non-zero length	SAVSX = (SCUMS/(SNS - SNZS))
	Mean % execution time	(SAVS/SAVT)*100
	Mean non-zero % of non-zero execution time	(SAVSX/SAVTX)*100
Subservice Behavior	Total number of invocations	SSNIA
for Requestor Service	Percent zero-time	(SSNZI _A /SSNI _A)*100
	Mean time per invocation	(SSTMI _A /SSNI _A)
	Mean time per non-zero time execution	(SSTMIA/(SSNIA - SSNZIA))
	Total number of WAIT SERVICE COMPLETES	SSNWA
	Percent zero time Waits	(SSNZWA/SSNWA)*100
	Mean time per Wait	(SSTW _A /SSNW _A)
	Mean time per non-zero time Wait	(SSTWA/(SSNWA - SSNZWA))

C.2.8 IPSS Task and Activity Statistics

These statistics provide insight into the service activity subordinate to a specific task, where a Task is a modeler defined facility activated as part of the execution of a Exogenous and Endogenous service. A Task is assumed to be comprised of a linear combination of one or more modeler defined activities which are identified as part of the Task declaration. A modeler defines an activity by associating an activity name with an endogenous service. Statistics are gathered for a Task by gathering statistics for its activities. This occurs in the following two step process:

- 1. The Task is explicitly activated via the INITIATE TASK statement.
- Statistics are automatically gathered for each Endogenous service which
 - a. is subordinate to the inservice containing the INITIATE TASK statement, and
 - has been identified as an activity associated with the Task.

The statistics gathering continues until the TERMINATE TASK statement is executed.

Table C.2.8-1 identifies the statistics to be gathered. However, the actual collection mechanism to be used within an IPSS model has not yet been defined.

Table C.2.8-1 IPSS Task and Activity Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition*
By Task General	Total number of tasks initiated		X	EU-IM
	Total number of tasks completed		X	EU-IM
	Total number of tasks still in execution		X	EU-IM
	Mean time per task initiation		X	EU-IM
By Activity Within Task	Total number of invocations by task		X	EU-IM
Task Initiation Statistics	Mean total activity processing Mean total queue time per		X	EU-IM
	task initiation		X	EU-DP
	Percent queue time of processing time		X	EU-DP
	Mean total activity acquire time per task initiation Percent acquire time of		X	EU-DP
	processing time Mean total activity service		X	EU-DP
	time per task initiation Percent service time of		X	EU-DP
A-+::+	processing time		X	EU-DP
Activity Invoke Statistics			X	EU-IM
	Mean activity queue time per invocation by task		X	EU-DP
	Percent queue time of pro- cessing time per invocation		X	EU-DP
	Mean activity acquire time per invocation by task		X	EU-DP
	Percent acquire time of pro- cessing time per invocation		X	EU-DP
	Mean activity service time per invocation by task		X	EU-DP
	Percent service time of pro- cessing time by task		X	EU-DP

^{*}Disposition codes used in this table have the following meanings:

EU-IM:

EU-DP:

C3. IPSS EXTENSIONS TO CHARACTERIZE COMPUTER NETWORKS

The standard IPSS and the IPSS/DBS have been designed to operate as independent simulators. However, an integrated model capability is required in order to characterize a complete system. In the IPSS this capability was achieved by permitting the concurrent execution of two asynchronous simulation processes which have the ability to communicate. Once this cability was defined, it was possible to extend the mechanism to accommodate any number of asynchronous processes.

The purpose of this section is to describe the mechanism to be used which will permit the modeling of information processing system networks. For this discussion a mode is assumed to be an individual model containing IPSS components:

Component	Number	Permitted
Request Streat	0	or 1
System Resources	0	or 1*
Storage Structure	0	or 1**
Data Base Access		or 1*
Data Structure	0	or 1***

- *Either the System Resources or Data Base Access Component must be present. Both can also be present at each mode.
- **Storage Structure Component can be present only when the System Resources Component is also present.
- ***Data Structure Component can be present only when the Data Base Access Component is also present.

A network is comprised of multiple nodes which communicate by initiating requests for services to another node. These services requests are treated the same as exogenous requests generated by a node's request stream. To establish a network the modeler defines a multiple facility pool - which is described in the following subsection.

C3.1 Multiple Facility Pools

Currently, in order to simulate a computer network, all nodes of the network have to be combined together to form a single facility pool consisting of at most one of each of the following components: system resource component, storage structure component, request stream component and model director component. Since the above approach does not facilitate independent testing of each node or evaluating the interaction between nodes, a modeling structure which allows multiple facility pools is needed. The basic structure of the IPSS nucleus which directs model execution is given in Figure C3.1-la. In order to consider multiple facility pools, this structure will have to be modified to that shown in Figure C3.1-lb. In the multiple facility pool configuration, the nucleus consists of three drivers:

- 1. Model Driver
 - The function of the model driver is the same as in the case of a single facility pool. It's basic function is to start and terminate the simulation drivers.
- 2. Global Simulation Driver The global simulation driver controls the FEQ (Future Events Queue) and the simulation clock. All future events generated by the different facility pools are combined together into a single FEQ so as to provide correct simulated time synchronization between the facility pools. The global simulation driver has a priority list that is used to determine which local simulation driver has control.
- 3. Local Simulation Driver Each facility pool has a local simulation driver which controls its corresponding CEQ (Current Event Queue). Once a local simulation driver is given control, it retains control until either (i) a transaction has to be placed in a different CEQ or (ii) a transaction has to be placed in the FEQ. In these cases, the local driver

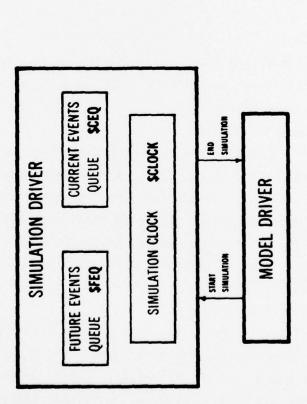


Figure C3.1-la Basic Structure of IPSS Simulation Nucleus

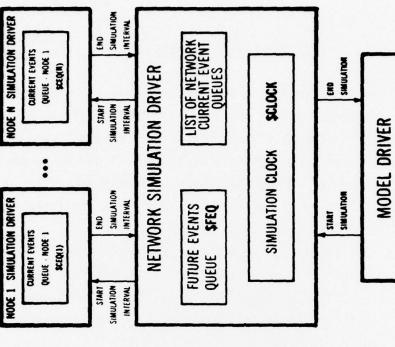


Figure C3.1-1b Structure of Model Nucleus to Support Networking

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returns control to the global driver which has the responsibility for placing the transaction into the FEQ or calling the appropriate local nucleus to insert a transaction into its CEQ. If placement is in a CEQ with higher priority than the currently active one, the interrupt flag is set. The currently active nucleus again resumes processing its active transaction until its processing is suspended. If the interrupt flag is set, control is given to the local driver with the higher priority via the global nucleus, otherwide the current driver continues with its next transaction. The process continues until all the CEQ's are empty at which time the clock is advanced to the next most imminent departure time in the FEQ. One or more transactions with this time are then moved to their corresponding CEQ and control is given to the nucleus with highest priority and the process begins a new.

C3.2 Changes to IPSS Simulation Nucleus

In order to accommodate the multiple facility pool structure, a number of changes had to be made to the IPSS nucleus and system subroutines.

- 1. The IPSS definitional tables (e.g., \$TOC, \$LBL, \$DIR, \$DEF) of the individual facility pools have to be combined together to form a single set of tables.
- To avoid naming conflicts between facility pools, a global interface routine is needed to invoke the currently implemented local interface routines.
- The system subroutines (e.g., EQUATE, INVOKE) that handle intrafacility pool communications have to be modified to handle inter-facility pool communications.

Detailed programming changes to IPSS are given below.

Changes to the IPSS Fortran System Subroutines to Handle Multiple Facility Pools

(A) The following procedures are to be modified:

1. \$FDEF

This procedure has to be modified so that it will accept base pointers to the definitional tables as parameters. A procedure for searching the System Resources and Data Base Access \$DEF tables through \$TOC, \$LBL and \$DIR has to be added to this procedure.

2. \$EQUAT

This procedure has to be modifed so that it will accept definitional table addresses and bases for the equated components as input parameters.

3. \$CREAT

This procedure has to be modified so that both the Facility's Pool ID and Driver ID are stored in the created transaction.

- \$GENER, \$GADR, \$SYSEX2, \$SEEK, \$EXINT
 All calls to the interface routine are to be modified.
- \$LINK, \$UNLINK, \$INSRT, \$REMOV
 All these procedures have to be modified to accept multiple CEQ's.

6. \$DRIVR

This procedure has to be modified to handle multiple CEQ's. The driver should automatically interchange base for the definitional table when a new CEQ is selected. All calls to interface routines should be modified. Also, all calls to initialization routines have to be modified so that all Facility Pool/Driver-relative components are initialized.

7. \$DVINT, \$ABORT

Parameter lists are to be deleted from these procedures and appropriate COMMON's should be added.

- 8. All statistical routines are to be modified so that separate statistics will be printed for each facility pool.
- All procedures that use the Create Data Set (\$CDS) arrays have to be changed.
- 10. All procedures that use absolute addresses for accessing the definitional table have to be changed so that they use relative addresses in the base table.
- (B) The following procedures to be added.
 - 1. \$BLOC

When the Facility Pool ID and Driver ID are given, this procedure will search through the base table to identify the bases to be used for the local simulation driver.

2. \$SUPEQ

This procedure handles all inter-facility pool equates and makes use of \$EQUAT to resolve all intra-facility pool equates.

3. \$INIT

This procedure initializes all endo and exo services with the facility pool and driver ID's. The endo and exo service definitions are found by searching through \$TOC, \$LBL and \$DIR.

4. \$INVEX

This procedure handles the invoking of an Exo service in another facility pool. Basically, this procedure facilitates interfacility pool communications.

- (C) The following Fortran tables and statements are to be generated by the language parsers.
 - 1. CEQ Table

This table will be used by the flobal simulation driver to determine the priority of the CEQ's.

\$PRTY	\$FPID	\$DRIVR	\$CBEG	\$CEND
			•	
•	•			•
	•	•	•	

\$PRTY: Priority of the CEQ

\$DPIF: Facility Pool ID
\$DRIVR: type of driver

1. if SR, RS or SS component driver

2. if DBA or DT component driver

\$CBEG: beginning addresses of CEQ's (initialized to zero's)

\$CEND: ending addresses of CEQ's (initialized to zero's)

2. BASE Table

This table stores the bases used for addressing the definitional tables.

\$FPID	\$CMPNT	\$\$TOC	\$\$LBL	\$\$DIR	\$\$DEF	\$CDS	\$NEVNT
			•	•	•	•	
	•		•				

\$FPID: Facility Pool ID

\$CMPNT: Component ID

type = 1 if SR component
type = 2 if RS component
type = 3 if SS component
type = 4 if DBA component
type = 5 if DT component

\$\$TOC: list of bases for \$TØC
\$\$ LBL: list of bases for \$LBL
\$\$DIR: list of bases for \$dir
\$\$DEF: list of bases for \$DEF

\$CDS: list of bases for create data set array

\$NEVNT: number of events that has occurred in the associated facility pool.

3. The following interface routine will be generated for each (facility pool, driver) pair.
SUBROUTINE \$xxxxx (TYPE, SYSIND, INDX1, INDX2,*)
INTEGER TYPE, SYSIND, INDX1, INDX2
GO TO (100, 200, 300, 400), TYPE

50 RETURN 1

100 CALL \$CFPR (SYSIND, &50) RETURN

200 CALL SYSER (INDX2, INDX1, &50)
RETURN

300 CALL \$DBPR (SYSIND, &50)
RETURN

400 CALL \$EXPR (SYSIND, &50)
RETURN
END

Where xxxxx is a unique name generated for each (facility pool, driver) pair.

4. The following global interface routine is generated for the model pool. SUBROUTINE \$SUPER (FPID, DRIV, TYPE, SYSIND, INDX1, XINDX2,*) INTEGER FPID, DRIV, TYPE, SYSIND INTEGER INDX1, INDX2, I
GO TO 9999

CALL \$xxxxx (TYPE, SYSIND INDX1, INDX2 \$10000)
RETURN

9999 I = FPID *2 + DRV-2 GO TO (1, 2, ... n), I 10000 RETURN 1 END

Where n is equal to the number of (facility pool, driver) pairs.